

Spectral Representations of the Transition Probability Matrices for Continuous Time Finite

Markov Chains

Author(s): Nan Fu Peng

Source: Journal of Applied Probability, Vol. 33, No. 1 (Mar., 1996), pp. 28-33

Published by: <u>Applied Probability Trust</u> Stable URL: http://www.jstor.org/stable/3215261

Accessed: 28/04/2014 10:47

Your use of the JSTOR archive indicates your acceptance of the Terms & Conditions of Use, available at http://www.jstor.org/page/info/about/policies/terms.jsp

JSTOR is a not-for-profit service that helps scholars, researchers, and students discover, use, and build upon a wide range of content in a trusted digital archive. We use information technology and tools to increase productivity and facilitate new forms of scholarship. For more information about JSTOR, please contact support@jstor.org.



Applied Probability Trust is collaborating with JSTOR to digitize, preserve and extend access to Journal of Applied Probability.

http://www.jstor.org

SPECTRAL REPRESENTATIONS OF THE TRANSITION PROBABILITY MATRICES FOR CONTINUOUS TIME FINITE MARKOV CHAINS

NAN FU PENG.* National Chiao Tung University

Abstract

Using an easy linear-algebraic method, we obtain spectral representations, without the need for eigenvector determination, of the transition probability matrices for completely general continuous time Markov chains with finite state space. Comparing the proof presented here with that of Brown (1991), who provided a similar result for a special class of finite Markov chains, we observe that ours is more concise.

MARKOV CHAINS: TRANSITION PROBABILITY MATRICES; SPECTRAL REPRESENTATIONS

AMS 1991 SUBJECT CLASSIFICATION: PRIMARY 60J35
SECONDARY 60J27

1. Introduction

It is undoubtedly important to calculate numerically the time-dependent transition probabilities of continuous time Markov chains. We focus our attention on those with a finite state space. Keilson developed in his book [5] the methods of spectral decomposition and the uniformization technique. Ross [10] found the external uniformization; this was followed by related work such as [7] and [12]. Some results on finite queues can be found in [1], [8], [9] and [11]. Brown [2] gave spectral representations, without eigenvectors, of the transition probability matrices of finite continuous time Markov chains with diagonalizable infinitesimal matrices (see also theorem 5 of [3]). Here we present an easy linear-algebraic technique which enables us to extend the result of [2] to completely general continuous time Markov chains with finite state space. The method used in this paper is also more concise and efficient than that of [2].

2. A simple linear-algebraic method

Consider a Markov chain (X(t)) defined on a finite state space $\{0, 1, 2, \dots, N\}$. Denote by $\lambda_0 = 0, \lambda_1, \dots, \lambda_N$ (maybe complex) the eigenvalues of its infinitesimal matrix Q. It is well known [5] that the transition probability matrix P(t) of X(t) is

Received 15 July 1994; revision received 19 December 1994.

^{*} Postal address: Institute of Statistics, National Chiao Tung University, 1001 TA Hsueh Road, Hsinchu, Taiwan.

(1)
$$P(t) = e^{Qt} = \sum_{n=0}^{\infty} \frac{Q^n t^n}{n!}.$$

Obviously, (1) implies the following:

(2)
$$P(0) = 1 \quad \text{and} \quad \frac{d^n P(t)}{dt^n} \bigg|_{t=0} = \left(\frac{d^n P_{ij}(t)}{dt^n}\right) = Q^n, \quad \forall n \ge 1.$$

If P(t) is a transition function or, more generally, sufficiently smooth, then (2) implies (1); hence we obtain the equivalence of (1) and (2). The linear algebra used below can be found in many textbooks, e.g. [4].

Lemma 1. Let A and B be two complex $n \times n$ matrices and $\{\bar{\alpha}_1, \dots, \bar{\alpha}_n\}$ be any basis of C^n . Then $A\bar{\alpha}_i = B\bar{\alpha}_i$ for all i implies A = B.

Although Theorem 1 is a special case of Theorem 3 below, it is worth listing the proof here for comparison with that of Theorem 3 and that of [2].

Theorem 1. If the λ_i are all distinct, then

(3)
$$P(t) = \prod_{i=1}^{N} (I - Q/\lambda_i) + \sum_{m=1}^{N} (Q/\lambda_m) \prod_{i \neq m, 0} [(I - Q/\lambda_i)/(1 - \lambda_m/\lambda_i)] \exp(\lambda_m t).$$

Proof. Call the right-hand side of (3) $\tilde{P}(t)$. It is easy to see that, for $m = 0, 1, \dots, N$,

$$\frac{d^n \tilde{P}(t)}{dt^n} \bigg|_{t=0} \bar{x}_m = \lambda_m^n \bar{x}_m = Q^n \bar{x}_m, \qquad n = 0, 1, 2, \dots$$

where \bar{x}_m is an eigenvector associated with the eigenvalue λ_m . Since the λ_m are all distinct, the \bar{x}_m form a basis of C^{N+1} . The $\tilde{P}(t)$ is obviously smooth, hence we obtain (3) from the fact that (2) implies (1) and Lemma 1.

The above proof gives us a natural extension of Theorem 1 to Theorem 2 below. We allow repeated eigenvalues here, and relabel them $\lambda_0 = 0, \lambda_1, \dots, \lambda_M$ as the distinct values.

Theorem 2. If the minimal polynomial of Q is of the form

$$g(x) = x \prod_{i=1}^{M} (x - \lambda_i), \qquad M \leq N,$$

with distinct $\lambda_0 = 0, \lambda_1, \dots, \lambda_M$, then P(t) is of the form (3) with N replaced by M.

The next corollary also appeared in [2].

Corollary 1. If (X(t)) is a finite birth and death process, then P(t) is of the form (3).

30 NAN FU PENG

Proof. The infinitesimal matrix Q of (X(t)) is tridiagonal and it is shown in [2] that its eigenvalues are real and distinct.

The following example makes Theorem 1 more plausible.

Example 1. Consider a continuous time Markov chain having state space $\{0, 1, 2, 3\}$ and starting from state 0 with infinitesimal matrix

$$Q = \begin{bmatrix} 0 & 1 & 2 & 3 \\ -\lambda & \lambda & 0 & 0 \\ 1 & 0 & -\lambda & \lambda & 0 \\ 0 & 0 & -\lambda & \lambda \\ \lambda & 0 & 0 & -\lambda \end{bmatrix}.$$

A simple argument shows that

(4)
$$P_{03}(t) = \sum_{n=1}^{\infty} P(T=4n-1) = e^{-\lambda t} \sum_{n=1}^{\infty} \frac{(\lambda t)^{4n-1}}{(4n-1)!},$$

where T is a random variable distributed as Poisson (λt). In a similar fashion, we have

(5)
$$P_{12}(t) = e^{-\lambda t} \sum_{n=1}^{\infty} \frac{(\lambda t)^{4n-3}}{(4n-3)!}.$$

Alternatively, observing that $0, -2\lambda, -\lambda + i\lambda$ and $-\lambda - i\lambda$ are the eigenvalues of Q, we obtain from (3) that

(6)
$$P_{03}(t) = e^{-\lambda t} \left[\frac{1}{4} e^{\lambda t} - \frac{1}{4} e^{-\lambda t} - \frac{1}{2} \sin(\lambda t) \right]$$

and

(7)
$$P_{12}(t) = e^{-\lambda t} \left[\frac{1}{4} e^{\lambda t} - \frac{1}{4} e^{-\lambda t} + \frac{1}{2} \sin(\lambda t) \right].$$

By introducing the Taylor expansions of the terms in the brackets of the right-hand sides of (6) and (7), we obtain the respective equivalence of (6) and (7) to (4) and (5).

3. The general result

A matrix Q is defined to be lower semitriangular if $Q_{ij} = 0$ for j > i + 1. It was claimed in Theorem 1.2 of [6] that, if Q is lower semitriangular with $Q_{i,i+1} \neq 0$ for all i, then its eigenvalues are distinct but may be complex. This statement is incorrect as the next simple counterexample shows.

Example 2. Let the matrix Q be

$$Q = \begin{bmatrix} 0 & 1 & 2 \\ -1 & 1 & 0 \\ 1 & -2 & 1 \\ 2 & 1 & 0 & -1 \end{bmatrix}.$$

The eigenvalues of Q are 0, -2 and -2. Neither Theorem 1 nor Theorem 2 can be applied to this case because the null space of Q+2I is of dimension 1. Theorem 3 below

deals with general Q and provides us with a way to settle the problem. Several lemmas are needed in order to prove that theorem.

Lemma 2.

$$\frac{d^{n}(t^{k}e^{\lambda t})}{dt^{n}}\bigg|_{t=0} = \begin{cases} 0 & n < k \\ k! & n = k \\ \binom{n}{k}k! \lambda^{n-k} & n > k. \end{cases}$$

Proof. By the product rule of derivatives, it is easy to see that if f(t) and g(t) are continuously differentiable functions,

(8)
$$(fg)^{(n)} = \sum_{i=0}^{n} \binom{n}{i} f^{(i)}g^{(n-i)}.$$

We immediately obtain the lemma by letting $f(t) = t^k$ and $g(t) = e^{\lambda t}$.

Lemma 3. For given $M \ge 1$, let

$$f(t) = \left[\prod_{m=1}^{M-1} \left(\frac{t}{a_m} + 1 \right)^{d_m} \right] \left(1 + \sum_{i=1}^{k} c_i t^i \right)$$

where the d_m are non-negative integers and $a_m \neq 0$ for $m = 1, \dots, M-1$. Then $f^{(n)}(0) = 0$, $n = 1, 2, \dots, K$, if and only if the c_n satisfy

(9)
$$-c_n = \sum_{\substack{i_m \le d_m \\ 0 < i_1 + \dots + i_{M-1} \le n}} \left(\prod_{m=1}^{M-1} \frac{\binom{d_m}{i_m}}{a_m^{i_m}} \right) c_{n-i_1 - \dots - i_{M-1}}, \qquad n = 1, 2, \dots, K,$$

with the conventions that $c_0 = 1$ and the right-hand side of (9) is zero when M = 1.

Proof. A quick application of (8) shows, for $M=2, 3, \cdots$ and any f_1, \cdots, f_M ,

$$\frac{d^{n}\prod_{i=1}^{M}f_{i}(t)}{dt^{n}} = \sum_{0 \leq i_{1}+\cdots+i_{M-1} \leq n} {n \choose i_{1}i_{2}\cdots i_{M}} \left(\prod_{m=1}^{M-1}f_{m}^{(i_{m})}(t)\right) f_{M}^{(n-i_{1}-\cdots-i_{M-1})}(t),$$

with $i_M = n - i_1 - \dots - i_{M-1}$ here. Hence for $n = 1, 2, \dots, K$,

32 NAN FU PENG

$$\begin{split} \frac{d^{n}f(t)}{dt^{n}} \bigg|_{t=0} &= \sum_{\substack{i_{m} \leq d_{m} \\ 0 \leq i_{1} + \dots + i_{M-1} \leq n}} \left\{ \frac{n!}{i_{1}! \cdots i_{M-1}! (n-i_{1} - \dots - i_{M-1})!} \right. \\ & \times \left(\prod_{m=1}^{M-1} \frac{d_{m}!}{(d_{m} - i_{m})! a_{m}^{i_{m}}} \right) (n-i_{1} - \dots - i_{M-1})! c_{n-i_{1} - \dots - i_{M-1}} \right\} \\ &= n! \sum_{\substack{i_{m} \leq d_{m} \\ 0 \leq i_{1} + \dots + i_{M-1} \leq n}} \left(\prod_{m=1}^{M-1} \frac{d_{m}!}{a_{m}^{i_{m}}} \right) c_{n-i_{1} - \dots - i_{M-1}} = 0 \end{split}$$

if and only if (9) holds.

Theorem 3. Let the minimal polynomial of Q be of the form $f(x) = \prod_{i=0}^{M} (x - \lambda_i)^{d_i}$ where the λ_i are distinct and $d_i \ge 1$. Then

(10)
$$P(t) = \sum_{i=0}^{M} \left(\sum_{j=0}^{d_i-1} \frac{R_{(i,j)}}{j!} (Q - \lambda_i I)^j t^j \right) e^{\lambda_i t}$$

where

(11)
$$R_{(i,j)} = \left(\prod_{m \neq i} \frac{(Q - \lambda_m I)^{d_m}}{(\lambda_i - \lambda_m)^{d_m}} \right) \left(I + \sum_{n=1}^{d_{i-j-1}} c_{i,n} (Q - \lambda_i I)^n \right)$$

and

$$-c_{i,n} = \sum_{k_m \leq m \neq i d_m, \ 0 < \Sigma_{m \neq i} \ k_m \leq n} \left(\prod_{m \neq i} \frac{\binom{d_m}{k_m}}{(\lambda_i - \lambda_m)^{k_m}} \right) c_{i,n - \Sigma_{m \neq i} \ k_m}, \qquad 1 \leq n \leq d_i - 1,$$

with $c_{i,0} = 1$.

Remark. It is easy to check that Theorem 3 reduces to Theorem 2 when $d_i = 1$ for $i = 0, \dots, M$.

Proof. Call the right-hand side of (14) $\tilde{P}(t)$. Due to the fact $(Q - \lambda_m I) = (Q - \lambda_i I) + (\lambda_i - \lambda_m)I$ and Lemma 3, $R_{(i,j)}(Q - \lambda_i I)^j$ for $0 \le j < d_i$ can be written as

(12)
$$R_{(i,j)}(Q - \lambda_i I)^j = w_{\beta}(Q - \lambda_i I)^{\beta} + \dots + w_{d_i}(Q - \lambda_i I)^{d_i} + (Q - \lambda_i I)^j$$

where the w are complex scalars depending on i and $\beta = \sum d_m - 1$.

With some algebra, Lemma 2 together with (10), (11) and (12) yield the following: $\tilde{P}(0)\bar{x}_i = I\bar{x}_i$ and

$$\frac{d^n \tilde{P}(t)}{dt^n} \bigg|_{t=0} \tilde{x}_i = \sum_{m=0}^n \binom{n}{m} (Q - \lambda_i I)^m (\lambda_i I)^{n-m} \tilde{x}_i$$
$$= (Q - \lambda_i I + \lambda_i I)^n \tilde{x}_i = Q^n \tilde{x}_i$$

where \bar{x}_i belongs to the null space of $(Q - \lambda_i I)^{d_i}$. Note that $(Q - \lambda_i I)^m \bar{x}_i = \bar{0}$ if $m \ge d_i$. Since these \bar{x}_i form a basis for C^{N+1} and $\tilde{P}(t)$ is sufficiently smooth, Lemma 1 and the implication of (2) to (1) yield the desired result.

Remark. Supposing the minimal polynomial is difficult to obtain, Theorem 3 still holds if we replace it with the characteristic polynomial.

Corollary 2. If (X(t)) is ergodic, then $\bar{\pi}' = (1/(N+1))\bar{1}' \prod_{i=1}^{M} (I - Q/\lambda_i)^{d_i}$ is the unique stationary vector of P(t), where $\bar{1}$ is the vector with all entries equal to 1.

Proof. Since $0 \le P_{ij}(t) \le 1$, the real part of each λ_k $(k \ne 0)$ is strictly negative and $d_0 = 1$. Hence $P(t) \to \prod_{i=1}^{M} (I - Q/\lambda_i)^{d_i}$ as $t \to \infty$. Since (X(t)) is ergodic, each row of $\prod_{i=1}^{M} (I - Q/\lambda_i)^{d_i}$ is the unique stationary vector $\bar{\pi}'$.

Note that irreducibility of (X(t)) implies ergodicity of (X(t)) [5].

Example 2 (Continued). The probability transition matrix P(t) corresponding to the infinitesimal matrix Q is

$$P(t) = \begin{pmatrix} 1/2 & 1/4 & 1/4 \\ 1/2 & 1/4 & 1/4 \\ 1/2 & 1/4 & 1/4 \end{pmatrix} + \begin{pmatrix} 1/2 & -1/4 & -1/4 \\ -1/2 & 3/4 & -1/4 \\ -1/2 & -1/4 & 3/4 \end{pmatrix} e^{-2t} + \begin{pmatrix} 0 & 1/2 & -1/2 \\ 0 & -1/2 & 1/2 \\ 0 & -1/2 & 1/2 \end{pmatrix} t e^{-2t}.$$

Acknowledgment

I am very thankful to the referee for many helpful comments.

References

- [1] ABATE, J., KIJIMA, M. AND WARD, W. (1991) Decompositions of the *M/M/I* transition function. *Queueing Systems* 9, 323–336.
- [2] Brown, M. (1991) Spectral analysis, without eigenvectors, for Markov chains. *Prob. Eng. Inf. Sci.* 5, 131-144.
- [3] FILL, J. A. (1992) Strong stationary duality for continuous-time Markov chains. Part I: Theory. J. Theor. Prob. 5, 45–70.
 - [4] HOFFMAN, K. AND KUNZE, R. (1971) Linear Algebra. Prentice Hall, New York.
 - [5] Keilson, J. (1979) Markov Chain Models Rarity and Exponentiality. Springer, Berlin.
- [6] KIJIMA, M. (1987) Spectral structure of the first-passage-time densities for classes of Markov chains. J. Appl. Prob. 24, 631-643.
- [7] KIJIMA, M. (1992) A note on external uniformization for finite Markov chains in continuous time. *Prob. Eng. Inf. Sci.* 6, 127–131.
- [8] Kijima, M. (1992) The transient solution to a class of Markovian queues. *Comput. Math. Appl.* 24, 17–24.
- [9] Parthasarathy, P. R. and Sharafali, M. (1989) Transient solution to the many server Poisson queue. A simple approach. *J. Appl. Prob.* **26**, 584–594.
- [10] Ross, S. M. (1987) Approximating transition probabilities and mean occupation times in continuous-time Markov chains. *Prob. Eng. Inf. Sci.* 1, 251–264.
- [11] SHARMA, O. P. AND DASS, S. (1988) Multiserver Markovian queues with finite waiting space. Sankhyā B 50, 428–431.
- [12] YOON, B. S. AND SHANTHIKUMAR, J. G. (1989) Bounds and approximations for the transient behavior of continuous-time Markov chains. *Prob. Eng. Inf. Sci* 3, 175–198.