# Performance Analysis of a Truncated Closed-Loop Power-Control Scheme for DS/CDMA Cellular Systems

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*Abstract—***This paper analyzes the system performance of a truncated closed-loop power-control (TCPC) scheme for uplinks in direct-sequence/code-division multiple-access cellular systems over frequency-selective fading channels. In this TCPC scheme, a mobile station (MS) suspends its transmission when the** *short-term fading* **is less than a preset cutoff threshold; otherwise, the MS transmits with power adapted to compensate for the** *short-term fading* **so that the received signal power level remains constant. Closed-form formulas are successfully derived for performance measures, such as system capacity, average system transmission rate, MS average transmission rate, MS power consumption, and MS suspension delay. Numerical results show that the analysis provides reasonable accuracy and that the TCPC scheme can substantially improve the system capacity, the average system transmission rate, and power saving over conventional closed-loop power-control schemes. Moreover, the TCPC scheme under realistic consideration of power-control error due to power-control step size, power-control period, power-control command loop delay, and MS velocity is further investigated. A closed-form formula is obtained to accurately approximate the system capacity of the realistic TCPC scheme.**

*Index Terms—***Power consumption, short-term fading, suspension delay, system capacity, transmission rate, truncated closed-loop power control (TCPC).**

#### I. INTRODUCTION

**I** N A direct-sequence/code-division multiple-access<br>(DS/CDMA) system, many users can transmit messages si-<br>multiprecently such the come redio obtained each using a gracific multaneously over the same radio channel, each using a specific spread-spectrum pseudonoise (PN) code [\[1](#page-9-0)]. Within a cell, the code channels in downlinks can be considered as mutually orthogonal because downlinks may exhibit synchronous CDMA transmission. However, these code channels in uplinks cannot be exactly mutually orthogonal for a set of asynchronous users; thus, mutual interference occurs among the uplinks. In such a case, a strong signal increases communication quality and a weak signal may suffer from strong interference. This problem is referred to as the near–far effect and limits the CDMA system capacity [[2\]](#page-9-0). Hence, power control is an essential issue in a DS/CDMA system.

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*Open-loop* power control, that is, the average power control, is applied to compensate for the long-term channel fading such that the average received signal power level is constant and the near–far problem is solved [\[3](#page-9-0)]. *Closed-loop* power control, however, is typically used to mitigate the short-term channel fading so that an acceptable received signal quality can be attained for the uplink communication. Several *closed-loop/open-loop* power-control schemes have been investigated, such as: 1) the well-known perfect power control, within which MS transmission power is adjusted to the exact inverse of the short-term fading and, thus, the received signal power level remains constant (such a method is also referred to as the "channel-inversion scheme") [[4\]](#page-9-0), [[5\]](#page-9-0); 2) combined power/rate control proposed in [\[6](#page-9-0)], which is the same as the perfect power control except in that MS holds its transmission power at  $Q_0/X_0$  and adapts its transmission rate to  $S(t) \cdot R_0/X_0$  when  $S(t) < X_0$ , where  $Q_0$  is the desired received power level,  $X_0$  is a preset cutoff threshold,  $S(t)$  is the short-term fading at time t, and  $R_0$  is the data symbol rate; 3) truncated average power control (TAPC) proposed in [\[7](#page-9-0)], which applies a truncated channel-inversion scheme to conventional average power control. This truncated channel-inversion scheme suspends transmission when the *long-term channel fading* falls below a cutoff threshold; otherwise, it adaptively controls power according to the channel-inversion scheme. By suspending transmission in this way, an improvement of system capacity was reported.

In this paper, we propose a truncated closed-loop power-control (TCPC) scheme, which extends the TAPC scheme by considering the fast closed-loop power control, since mitigating the rapid variations of fading by power-control mechanism had been adopted in third-generation (3G) systems such as cdma2000 and WCDMA. In the TCPC scheme, when the short-term fading satisfies  $S(t) < X_0$ , MS suspends transmission; otherwise, MS adaptively transmits power to compensate for the *short-term fading*, so that the received signal power level is constant. This paper analyzes the performance of the TCPC scheme for uplinks in DS/CDMA cellular systems over frequency-selective fading channels. It first analyzes an ideal TCPC scheme in which the transmission power is continuously and instantaneously adjusted. Based on the signal-to-interference ratio (SIR) formula over a medium-term period, closed-form formulas are successfully derived for system performance including system capacity, average system transmission rate, MS average transmission rate, MS power consumption, and MS suspension delay. The upper limit of the average system transmission rate of the TCPC scheme

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achieved in a multicell system is also obtained and it is found to be the same as that of the perfect power-control scheme in a single-cell system. Numerical results show that the analysis is quite accurate; the TCPC scheme is more effective than the conventional closed-loop power-control schemes such as perfect power control [\[4](#page-9-0)], [[5\]](#page-9-0), combined power/rate control [[6\]](#page-9-0), and TAPC [\[7](#page-9-0)]. This paper further analyzes the realistic TCPC scheme that considers power-control error due to power-control step size, power-control command loop delay, and MS velocity. A closed-form formula is also derived to accurately approximate the system capacity of the realistic TCPC scheme. Numerical results indicate that the power-control error has a significant impact on the system capacities of the TCPC and conventional closed-loop power-control schemes.

The rest of this paper is structured as follows. Section II introduces the system model and derives a formula of the received SIR per bit. Section III derives the performance measures of the TCPC scheme in both ideal and realistic cases. Section IV presents numerical results as well as simulation results for validation. Comparisons between TCPC and other conventional schemes are also presented. Finally, Section V is the conclusion.

# II. SYSTEM MODEL

# *A. Channel Model*

Consider  $N_B$  cells in a CDMA cellular system of which the central cell is surrounded by other cells in a hexagonal-grid configuration. Each cell has a base station (BS) located at the center and has  $N_M$  MSs uniformly distributed in the system. Generally, the frequency-selective multipath channel of the uplink  $(m, b)$ , denoting the channel from any MS m to any BS  $b$ , is described by a time-variant impulse response , where  $T_c$  represents the chip duration, K is the number of resolvable paths, and the random variables  $\alpha_{mb,p}(t)$  and  $\theta_{mb,p}(t)$  are the power gain and phase of the pth path at time t, respectively [[1\]](#page-9-0).

The power gain of the *th path can be divided into two* parts,  $\alpha_{mb,p}(t) = L_{mb}(t) \cdot S_{mb,p}(t)$  [\[8](#page-9-0)], where  $L_{mb}(t)$ and  $S_{mb,p}(t)$  represent long-term and short-term fading, respectively. Long-term fading is normally modeled as  $L_{mb}(t) = |\mathbf{z}_m(t) - \mathbf{x}_b(t)|^{-\eta} \cdot 10^{x_{mb}(t)/10}$ , where  $\mathbf{z}_m(t)$  and  $x_b(t)$  are complex numbers that represent the locations of MS m and BS b, respectively.  $\eta$  is the propagation exponent that depends on the environment and  $x_{mb}(t)$  is a zero-mean Gaussian-random process with standard deviation  $\sigma_x$ . The short-term fading is widely modeled as Rayleigh fading and  $S_{mb,p}(t)$  is exponentially distributed. For convenience, the multipath intensity profile (MIP) [[1](#page-9-0)] of each uplink is assumed to be identical and every signal path has the same average received power. Without loss of generality, the overall short-term fading effect  $S_{mb}(t) \triangleq \sum_{p=1}^{K} S_{mb,p}(t)$ , given the optimal maximal ratio combing, is normalized to have unit mean. Thus, the probability density function (pdf) of  $S_{mb}(t)$ , denoted by  $f_S(\cdot)$ , is found to have a gamma distribution.

# *B. Transmitter Model*

If binary phase-shift keying (BPSK) modulation is considered, the transmitted low-pass equivalent signal of MS  $m$  is given by  $s_m(t) = \sqrt{P_m(t)} \cdot d_m(t - \tau_m) \cdot c_m(t - \tau_m) \cdot e^{-j\phi_m}$ , where  $P_m(t)$  is MS transmission power at time t,  $d_m(t)$  is the MS bipolar data stream,  $c_m(t)$  is the corresponding PN sequence whose chip waveform is rectangular,  $\tau_m$  is a value between zero and the data symbol duration  $T$  and indicates each MS independent symbol timing due to asynchronous transmission, and  $\phi_m$  is the random carrier phase. Under a transmission power control,  $P_m(t)$  can be conceptually divided into two parts as  $P_m(t) = P_{O,m}(t) \cdot P_{C,m}(t)$ , where  $P_{O,m}(t)$ and  $P_{C,m}(t)$  are the transmission powers controlled by openand closed-loop power-control schemes, respectively. This paper assumes perfect open-loop power control; hence, for each MS m served by its home BS h, we have  $P_{O,m}(t) = 1/L_{mh}(t)$ .

# *C. Receiver Model*

Here, we consider a medium-term period that is sufficiently long to make the short-term fadings be averaged out while the long-term fadings remain almost constant, since MSs move only a little. Also, the set of the long-term fading from any MS  $m$  to any BS *b* is referred to as a scenario and is represented as  $\{L_{mh}\},\$ from which its time index  $t$  is dropped. Based on the models of channel impulse response and the MS transmitted signal, the total signal received at a BS  $H$  can be found to be as shown in (1) at the bottom of the page, where the background noise is ignored since it is much smaller than the other-cell interference. Consider a RAKE receiver with an optimal maximal ratio combiner to take full advantage of the multipath diversity. Here, we assume slow fading such that  $S_{mb}(t)$ ,  $S_{mb,p}(t)$ ,  $P_m(t)$ , and  $\theta_{mb,p}(t)$  can be treated as constants  $S_{mb}[k]$ ,  $S_{mb,p}[k]$ ,  $P_m[k]$ , and  $\theta_{mb,p}[k]$ , respectively, for all possible m, b, and p, during

$$
r_H(t) = \sum_{m=1}^{N_B \cdot N_M} \sum_{p=1}^{K} \sqrt{L_{mH} \cdot S_{mH,p}(t) \cdot P_m(t)} \cdot d_m(t - pT_c - \tau_m) \cdot c_m(t - pT_c - \tau_m) \cdot e^{-j(\theta_{mH,p}(t) + \phi_m)}
$$
(1)

$$
Z_M[k] \triangleq \text{Re}\left\{\sum_{p=1}^K \int_{kT+pT_c+\tau_M}^{(k+1)T+pT_c+\tau_M} r_H(t) \cdot \left(\sqrt{S_{MH,p}(t)} \cdot c_M(t-pT_c-\tau_M) \cdot e^{-j(\theta_{MH,p}(t)+\phi_M)}\right)^* dt\right\}
$$
  
=  $r_M[k] + I_A[k] + I_S[k].$  (2)

the kth data symbol period. The decision statistics  $Z_M[k]$  corresponding to an MS  $M$  at cell  $H$  can be derived and divided into three parts, as shown in (2) at the bottom of the previous page.

The first term  $r_M[k]$  is the desired signal part given by  $r_M[k] = T \cdot d_M[k] \cdot S_{MH}[k] \cdot \sqrt{L_{MH} \cdot P_M[k]}$ . The second term  $I_A[k]$  represents the multiple access and multipath interference from the other MSs and can be obtained as shown by (3) at the bottom of the page, where [see (4) at the bottom of the page]. Over the medium-term period, the long-term fading, which is the local mean of the channel gain, is treated as constant and the short-term fading is typically modeled as a stationary random process. Hence, the channel gain, which is the product of the long- and short-term fading, becomes a stationary random process. The transmission power under strength-based power control, which generally is a function of channel gain, will also turn out to be stationary. If both MS data streams and PN sequences are assumed to be stationary, then the interference signal  $I_A[k]$  in (4) will resemble a noise-like stationary random process and its distribution can be approximated by a Gaussian one according to the central limit theorem. Notably,  $I_A[k]$  has a mean of zero and a variance, which is the mean of the interference signal power at the output of the RAKE receiver, given by

$$
\mathbf{E}\left(I_A^2[k]\right) = \frac{G \cdot T_c^2}{3} \cdot S_{MH}[k] \cdot S
$$

where  $G \triangleq T/T_c$  is the processing gain [[1\]](#page-9-0) and  $\mathbf{E}(\cdot)$  is the expectation function. The third term  $I_S[k]$  represents the multipath interference itself, since  $N_B \cdot N_M \gg 1$ ,  $I_S[k]$  would be much smaller than  $I_A[k]$  and the effect of  $I_S[k]$  is ignored in the following.

The SIR per bit during the  $k$ th data symbol period of MS  $M$  at the output of RAKE receiver, defined in [[1,](#page-9-0) p. 244 ] and denoted by  $\Gamma_M[k]$ , is given by

$$
\Gamma_M[k] = \frac{r_M^2[k]}{2\mathbf{E}(I_A^2[k])} = \frac{3}{2} \cdot \frac{G \cdot Q_M[k]}{I_I + I_O}
$$
(6)

where  $Q_M[k]$  is the instantaneously received signal power and  $I_I$  and  $I_O$  represent the mean power of the intracell interference signals and the mean power of the other-cell interference signals, respectively. In general, the term  $Q_M[k]$  can be obtained by

$$
Q_M[k] = P_M[k] \cdot S_{MH}[k] \cdot L_{MH}
$$
  
=  $(P_{C,M}[k] \cdot S_{MH}[k]) \cdot (P_{O,M}[k] \cdot L_{MH}).$ 

Assuming that  $\{P_{C,M}[k], \forall m\}$  are independently and identically distributed (i.i.d.) with a mean of  $\overline{P_C} \triangleq \mathbf{E}(P_{C,m}[k]),$  ${Q_m[k], \forall m}$  are also i.i.d. random variables with a mean of  $\overline{Q} \triangleq {\bf E}(Q_m[k])$ . From (5) and (6),  $I_I$  and  $I_O$  can be obtained by

$$
I_I = (N_M - 1) \cdot \overline{Q} \tag{7}
$$

$$
I_O = \overline{P_C} \cdot \sum_{h \neq H} \sum_{m \in MS^{(h)}} \frac{L_{mH}}{L_{mh}}
$$
(8)

where the  $MS^{(h)}$  in (8) refers to the set of the MSs that are served by BS  $h$ .

# III. PERFORMANCE ANALYSIS

In the TCPC scheme, MS adjusts its transmission power to compensate for the *short-term fading* when the it is above a preset cutoff threshold  $X_0$  and suspends its transmission otherwise. Accordingly, considering an uplink  $(M, H)$ 

$$
P_M[k] = \begin{cases} \frac{Q_0}{L_{MH}[k] \cdot S_{MH}[k]}, & \text{as } S_{MH}[k] \ge X_0\\ 0, & \text{as } S_{MH}[k] < X_0 \end{cases} \tag{9}
$$

where  $Q_0$  denotes the preset desired received power level. If the transmission power can be adjusted *continuously* and *immediately*, as shown in (9), the scheme is referred to as an *ideal* TCPC.

# *A. Ideal TCPC*

Conditional with  ${L_{mb}}$ , the SIR per bit in (6) given  $P_M[k] >$  $\theta$  is time constant over the medium-term period. The unconditioned SIR per bit is then obtained by considering all possible  ${L_{mb}}$ . The formula of the unconditioned SIR per bit is the same as (6), but its denominator is now a random variable due to the randomness of  ${L_{mb}}$ . Again, applying the central limit theorem, the normalized  $I_0$ , represented as  $I_O^* \triangleq I_O / \overline{P_C}$ , can be approximated by a Gaussian random variable  $\xi$  and expressed as

$$
I_O^* \approx \xi(N_M \mu_O, N_M \sigma_O^2) \tag{10}
$$

$$
I_A[k] = \sum_{p=1}^{K} \sqrt{S_{MH,p}[k]} \cdot \sum_{m=1, m \neq M}^{N_B \cdot N_M} \sum_{p'=1}^{K} \sqrt{L_{mH} \cdot S_{mH,p'}[k] \cdot P_m[k]} \cdot I_{M,p}^{(m,p')}[k]
$$
(3)

# $I_{M,p}^{(m,p')}[k] = \int_{k\tau + n\tau_{-} + \tau_{M}}^{(k+1)T + pT_{c} + \tau_{M}} d_{m}(t-p'T_{c}-\tau_{m}) \cdot c_{m}(t-p'T_{c}-\tau_{m}) \cdot c_{M}(t-pT_{c}-\tau_{M}) \cdot \cos(\theta_{mH,p'}[k] + \phi_{m} - \theta_{MH,p}[k] - \phi_{M}) dt.$  (4)

where  $\mu_O$  and  $\sigma_O^2$  are the mean and variance of  $I_O^*$ , respectively, when  $N_M = 1$ . The terms  $\mu_O$  and  $\sigma_O^2$  can be obtained by

$$
\begin{cases}\n\mu_O = \mathbf{E} \left( \sum_{h \neq H} I_h \right) = \sum_{h \neq H} \mathbf{E} (I_h) \\
\sigma_O^2 = \text{var} \left( \sum_{h \neq H} I_h \right) = \sum_{h \neq H} \text{var} (I_h)\n\end{cases} (11)
$$

where  $I_h \triangleq (L_{mH} L_{mh}^{-1} | L_{mb} \leq L_{mh}, \forall b)$ . The  $I_h$  was analyzed by Zorzi [[9\]](#page-10-0) and its mean and variance can be numerically computed as

$$
\mathbf{E}(I_h) = \lim_{n \to \infty} \left( n - \int_0^n \Pr(I_h \le s) \cdot ds \right) = 1 - \int_0^1 \Pr(I_h \le s) \cdot ds \tag{12}
$$

$$
\text{var}(I_h) = \left(1 - 2 \int_0^1 s \Pr(I_h \le s) \cdot ds\right) - \mathbf{E}^2(I_h) \tag{13}
$$

where

$$
\Pr(I_h \le s) = \frac{\int \Pr(L_{mH}(\mathbf{z}) \le sL_{mh}(\mathbf{z}), L_{mb}(\mathbf{z}) \le L_{mh}(\mathbf{z}), \forall b) \cdot f_{\mathbf{z}_m}(\mathbf{z}) d\mathbf{z}}{\int_{\Pr(L_{mb}(\mathbf{z}) \le L_{mh}(\mathbf{z}), \forall b) \cdot f_{\mathbf{z}_m}(\mathbf{z}) d\mathbf{z}}}
$$
\n
$$
= \frac{\varphi(s)}{\varphi(1)} \tag{14}
$$

 $f_X(\cdot)$  denotes the pdf of a random variable X, and

$$
\varphi(s) = \Pr(L_{mH}(\mathbf{z}) \le sL_{mh}(\mathbf{z}), L_{mb}(\mathbf{z}) \le L_{mh}(\mathbf{z}), \forall b)
$$
  
= 
$$
\int f_{L_{mh}(\mathbf{z})}(y) \Pr(L_{mH}(\mathbf{z}) \le sy) \prod_{b \ne h, H} \Pr(L_{mb}(\mathbf{z}) \le y) dy.
$$
 (15)

Considering an MS  $M$  in central cell  $H$ , the outage probability of the uplink  $(M, H)$ , denoted by  $\Lambda_M$ , is defined as

$$
\Lambda_M \triangleq \Pr(\Gamma_M[k] < \Gamma_0 | P_M[k] > 0) \\
= \Pr\left(\frac{1.5G \cdot Q[k]}{(N_M - 1) \cdot \overline{Q} + \overline{P_C} \cdot I_O^*} < \Gamma_0\right) \tag{16}
$$

where  $\Gamma_0$  is the minimum SIR per bit required to achieve a desired bit-error rate at the output of the RAKE receiver. The outage probability given  $P_M[k] = 0$  is not considered, since no data is transmitted in that case. Notably, only the statistics relating to the central cell  $H$  are taken into account to avoid the corner effect. Using (6) and (10),  $\Lambda_M$  can be approximated by

$$
\Lambda_M \approx \Pr \left( \frac{1.5G \cdot Q_0}{(N_M - 1) \cdot \overline{Q} + \overline{P_C} \cdot \xi \left( N_M \mu_O, N_M \sigma_O^2 \right)} < \Gamma_0 \right)
$$
\n
$$
= 1 - F_{\xi} \left( \frac{1.5G \cdot Q_0 - \Gamma_0 \cdot (N_M - 1) \cdot \overline{Q}}{\Gamma_0 \cdot \overline{P_C} \cdot \sqrt{N_M}} - \frac{N_M \mu_O}{\sqrt{N_M}}, 0, \sigma_O^2 \right) \tag{17}
$$

where  $F_{\xi}(\cdot,\mu_O,\sigma_O^2)$  indicates the cumulative distribution function (cdf) of a Gaussian variable  $\xi$  with mean  $\mu_0$  and variance  $\sigma_O^2$ . The average transmission power  $\overline{P_C}$ , controlled by the closed-loop power control, can be obtained by

$$
\overline{P_C} = \int_{X_0}^{\infty} \frac{Q_0}{y} f_S(y) dy
$$
  
= 
$$
\begin{cases} Q_0 \cdot \int_{X_0}^{\infty} \frac{1}{y \cdot e^y} dy, & K = 1 \\ Q_0 \cdot \frac{K}{K-1} \cdot (1 - F_{\gamma(K-1,K)}(X_0)), & K \ge 2 \end{cases}
$$
(18)

where  $F_{\gamma(p_1,p_2)}(\cdot)$  denotes the cdf of a two-parameter gammarandom variable  $\gamma$  with parameters  $p_1$  and  $p_2$ . The average received power  $\overline{Q} = \mathbf{E}(Q_M[k])$  can be obtained by

$$
\overline{Q} = \int_{X_0}^{\infty} Q_0 \cdot f_S(y) dy = Q_0 \cdot \left(1 - F_{\gamma(K,K)}(X_0)\right). \tag{19}
$$

Five performance measures are investigated. The first measure is the system capacity  $C$ , which is defined as the maximum number of users per cell that the system can support under the constraint that the outage probability of an arbitrary MS in central cell H is less than a preset outage threshold  $\Lambda_0$ . That is

$$
C \triangleq \left[ \max \left\{ \underset{N_M \in \mathfrak{R}}{\arg} \left[ \left( \Lambda_M < \Lambda_0 \right) \middle|_{M \in MS(H)} \right] \right\} \right] \tag{20}
$$

where the operator  $\lvert \cdot \rvert$  denotes the maximum integer below the argument. Substituting (17) into (20) yields  $C = \lfloor C^* \rfloor$  where  $C^*$  satisfies

$$
F_{\xi} \left( \frac{1.5G \cdot Q_0 - \Gamma_0 \cdot (C^* - 1) \cdot \overline{Q}}{\Gamma_0 \cdot \overline{P_C} \cdot \sqrt{C^*}} - \frac{C^* \cdot \mu_O}{\sqrt{C^*}}, 0, \sigma_O^2 \right) = 1 - \Lambda_0 \quad (21)
$$

which is rewritten in the form of  $a \cdot C^* + b \cdot \sqrt{C^*} - c = 0$  where

$$
\begin{cases}\na = \Gamma_0 \cdot (\overline{Q} + \overline{P_C} \cdot \mu_O) \\
b = \Gamma_0 \cdot \overline{P_C} \cdot F_{\xi}^{-1} (1 - \Lambda_0, 0, \sigma_O^2) \\
c = 1.5G \cdot Q_0 + \Gamma_0 \cdot \overline{Q}.\n\end{cases}
$$
\n(22)

Accordingly, the capacity  $C$  under the TCPC scheme is given by

$$
C = \left[ \left( \frac{-b + \sqrt{b^2 + 4ac}}{2a} \right)^2 \right] = \left[ \left( -\frac{b}{2a} + \sqrt{\left( \frac{b}{2a} \right)^2 + \frac{c}{a}} \right)^2 \right] \tag{23}
$$

Notably, when  $X_0 = 0$ , the TCPC is reduced to the perfect power control and the above analyzes remain applicable. On the other hand, when  $X_0 \rightarrow \infty$ ,  $b/a \rightarrow 0$ ,  $c/a \rightarrow \infty$ , and the capacity becomes infinite.

The second measure is the MS average transmission rate  $R$ , which is

$$
\overline{R} \triangleq \int_{X_0}^{\infty} R_0 \cdot f_S(y) dy = R_0 \cdot \left(1 - F_{\gamma(K,K)}(X_0)\right) \quad (24)
$$

where  $R_0 = 1/T$  is the symbol rate.

The third measure is the average system transmission rate  $\Re$ , which is defined as the average transmission rate multiplied by the system capacity. Thus,  $\overline{\Re}$  is given by

$$
\overline{\Re} \triangleq \overline{R} \cdot C \approx \left( -\frac{b \cdot \sqrt{\overline{R}}}{2a} + \sqrt{\left( \frac{b \cdot \sqrt{\overline{R}}}{2a} \right)^2 + \frac{c \cdot \overline{R}}{a}} \right)^2 \tag{25}
$$



Fig. 1. Functional blocks of the realistic TCPC scheme.

such that  $\overline{\Re}$  can be approximated as a monotonously increasing function of  $X_0$ . When  $X_0 = 0$ , (25) corresponds to the average system transmission rate of the perfect power-control scheme and is the lower limit of  $\overline{\Re}$ . While  $X_0 \rightarrow \infty$ , the factor  $\sqrt{b} \cdot \overline{R}^{1/2}/a$  in (25) approaches zero and the upper limit of the average system transmission rate in a multicell system under TCPC is

$$
\lim_{X_0 \to \infty} \overline{\Re} = \lim_{X_0 \to \infty} \frac{c \cdot \overline{R}}{a}
$$
\n
$$
= \lim_{X_0 \to \infty} \frac{R_0 \cdot (1.5G \cdot Q_0 + \Gamma_0 \cdot \overline{Q})}{Q_0 \cdot \Gamma_0 \cdot \left(1 + \frac{\overline{P_C}}{\overline{Q}} \cdot \mu_0\right)}
$$
\n
$$
= \frac{1.5G \cdot R_0}{\Gamma_0} \tag{26}
$$

which is just the upper limit of the average system transmission rate in a single-cell system under perfect power control.

The fourth measure is the MS average transmission energy per bit  $\overline{\mathfrak{E}}$ , which is defined as the average ratio of MS transmission power to its transmission rate when data is being transmitted.  $\overline{\mathfrak{E}}$  is given by

$$
\overline{\mathfrak{E}} \triangleq \mathbf{E}\left(\left.\frac{P_M[k]}{R_M[k]}\right| P_M[k] > 0\right) = \mathbf{E}\left(L_{MH}^{-1}\right) \cdot \frac{\overline{P_C}}{\overline{R}}.\tag{27}
$$

The last measure is the MS average suspension delay  $\overline{D}$ . It is, in fact, the average fade duration defined as [[10,](#page-10-0) p. 189]

$$
\overline{D}\left(r = \sqrt{X_0}\right) \triangleq \frac{\Pr\left(r \leq \sqrt{X_0}\right)}{\operatorname{lcr}\left(r = \sqrt{X_0}\right)}\tag{28}
$$

where  $r \triangleq \sqrt{S_{MH}(t)}$  and  $lcr(r = \sqrt{X_0})$  is the level-crossing rate at a given strength level  $\sqrt{X_0}$ . Envelope r follows a Nakagami distribution, since the overall short-term fading  $S_{MH}(t)$ follows a gamma distribution. Consequently, as in [[11](#page-10-0)],  $lcr(\cdot)$ and  $D(\cdot)$  are

$$
lcr\left(r = \sqrt{X_0}\right) = f_m \cdot f_{\gamma(K,K)}(X_0) \sqrt{\frac{2\pi X_0}{K}} \tag{29}
$$

and

$$
\overline{D}\left(r=\sqrt{X_0}\right) = \frac{F_{\gamma(K,K)}(X_0)}{f_m \cdot f_{\gamma(K,K)}(X_0)} \sqrt{\frac{K}{2\pi X_0}}
$$
(30)

where  $f_m$  is the maximal Doppler frequency and  $f_{\gamma(p_1,p_2)}(\cdot)$ denotes the pdf of a two-parameter gamma-random variable  $\gamma$ with parameters  $p_1$  and  $p_2$ .

# *B. Realistic TCPC*

Considering that MS transmission power is adjusted discretely according to finite-bit power-control commands and factors of power-control error (PCE) due to step size for power adaptation  $\Delta_p$ , power-control command loop delay  $T_d$ , MS velocity V, and power-control period  $T_p$ , the TCPC scheme is referred to as a realistic TCPC scheme.

Fig. 1 presents the functional blocks of the realistic TCPC scheme. The transmitted signal  $s_M(t)$  with power  $P_M(t)$  from MS M will pass through the channel  $h_{MH}(\tau;t)$  and suffer interference from other MS transmissions. The received signal  $r_H(t)$ at the BS  $H$  is then fed into the optimal RAKE receiver to extract the averaged signal power  $Q_M[n]$  and to estimate the average short-term channel fading  $\overline{S_{MH}}[n]$  over the *n*th power-control period. The BS  $H$  has a power-control command unit (PCCU) to generate the power-control command *CMD* periodically

$$
CMD = \begin{cases} +1, & \text{if } \overline{Q_M}[n] < Q_0, \overline{S_{MH}}[n] \ge X_0 \\ -1, & \text{if } \overline{Q_M}[n] \ge Q_0, \overline{S_{MH}}[n] \ge X_0 \\ \text{"suspend,"} & \text{if } \overline{S_{MH}}[n] < X_0 \,. \end{cases}
$$
(31)

The  $CMD$  is then transmitted through the downlink channel to the destination MS, which has a transmission-control unit (TCU) that adjusts  $P_M(t)$  according to the received CMD. It suspends the transmission temporarily if the received  $CMD$  is "*suspend*" and adapts its transmission power by an amount of  $CMD \cdot \Delta_p$  dB otherwise.

In the realistic TCPC scheme, the received power is no longer a time constant. Conditional with  $Q_M[k] > 0$ , the PCE is defined as  $\psi = Q_M[k]/Q_0$ . Thus, the outage probability  $\Lambda_M$  turns out to be the average of (16) over all possible PCE values. Using (10),  $\Lambda_M$  can be approximated by (32), shown at the bottom of the page, where  $f_{\xi}(x; \mu_1, \sigma_1^2) = dF_{\xi}(x; \mu_1, \sigma_1^2)/dx$  is the Gaussian pdf function and

$$
\mu_1 = \frac{\Gamma_0 (N_M - 1) \cdot Q + \Gamma_0 P_C N_M \mu_O}{1.5 G \cdot Q_0}
$$

$$
\sigma_1^2 = \left(\frac{\Gamma_0 \overline{P_C}}{1.5 G \cdot Q_0}\right)^2 \cdot N_M \sigma_O^2. \tag{33}
$$

As shown in (32),  $\Lambda_M$  equals an arithmetic average of  $F_{\psi}(x)$ with weighting function of  $f_{\xi}(x; \mu_1, \sigma_1^2)$ . Since  $f_{\xi}(x; \mu_1, \sigma_1^2)$ is a bell-shaped-like curve with peak at  $x = \mu_1$ ,  $\Lambda_M$  can be accurately approximated by

$$
\Lambda_M \approx \int_0^\infty f_\xi(x; \mu_1, 0) F_\psi(x) dx = F_\psi(\mu_1). \tag{34}
$$

Based on the system-capacity definition given in (20), the system capacity can then be approximated by

$$
C \approx \left[ \frac{\frac{1.5 \cdot F_{\psi}^{-1}(\Lambda_0) \cdot G}{\Gamma_0} + \frac{\overline{Q}}{Q_0}}{\frac{\overline{P_C} \mu_0}{Q_0} + \frac{\overline{Q}}{Q_0}} \right].
$$
 (35)

From (35), it is found that the factor  $G/\Gamma_0$  is usually a number larger than other factors, whereby the PCE  $\psi$  plays an important role in affecting the system capacity  $C$ . The higher the dispersion degree of PCE is, the lower the factor  $F_{\psi}^{-1}(\Lambda_0)$  and the system capacity will be. Note that merely the factor  $F_{\psi}^{-1}(\Lambda_0)$  is enough to estimate the system capacity and that the knowledge of the whole distribution of PCE  $\psi$  is not necessary. On the other hand, the factor  $\mu_{\text{O}}$  has nothing to do with closed-loop power control and is the same as that given in (11). The  $P_C$  and Q in the realistic TCPC scheme depend on several factors including cutoff threshold  $X_0$ , step size  $\Delta_p$ , power-control command loop delay  $T_d$ , number K of resolvable paths, and MS velocity V. Hence, it is difficult to derive general formulas for  $P_C$  and  $Q$ . The simulation shows that the  $\overline{P_C}$  and  $\overline{Q}$  in the realistic TCPC scheme can be adequately approximated by the analytical results given in (18) and (19), respectively. Based on the above approximations and by assuming that the factor  $F_{\psi}^{-1}(\Lambda_0)$  is obtainable, the system capacity can be calculated.

Additionally, we assume that the average short-term channel fading  $\overline{S_{MH}}[n]$  is still Rayleigh distributed. Therefore, the performance measures  $\overline{R}$ ,  $\overline{\mathfrak{E}}$ , and  $\overline{D}$  have the same formulas given in (24), (27), and (30), respectively. Finally,  $\Re$  is computed according to its definition. Notably, as  $X_0 = 0$ , the realistic TCPC becomes a realistic perfect power control and the above analyzes remain applicable.

The realistic TCPC has an issue with how the transmitter can understand when  $\overline{S_{MH}}[n]$  turns out to be good after a bad period. Using the real TCPC applied to a WCDMA system [\[12](#page-10-0)] as an example, the WCDMA system has two types of uplink dedicated physical channels (DPCHs): the uplink dedicated physical data channel (DPDCH) and the uplink dedicated physical control channel (DPCCH). Each connection is allocated a DPCH including one DPCCH and several DPDCHs. The DPDCH is used to send user data and its transmission is suspended when the received CMD is "suspend." The DPCCH

is used to carry control information, including pilot symbols, power-control command for downlink power control, and rate information of current uplink transmission, required by the BS and the DPCCH is basically always on, even when the DPDCH is suspended. Thus, whenever the transmission of DPDCH is suspended, the receiver can somehow monitor the DPCCH channel condition and generate the power-control command according to (31).

Another issue is related to the MS velocity. By (30), for a mobile speed lower than 320 km/h, the average fade duration can be found greater than or equal to one power-control period and the TCPC scheme is still workable. As the mobile speed increases, the appearance probability of "suspend" command decreases and the TCPC scheme degrades toward the conventional closed-loop power-control scheme.

# IV. NUMERICAL RESULTS

In the following examples, the DS/CDMA system has  $N_B =$ 19 cells, the propagation exponent  $\eta = 3.5$ , the shadowing effect  $\sigma_x = 8$ , the resolvable path number  $K = 2$ , the processing gain  $G = 128$ , the outage threshold  $\Lambda_0 = 1\%$ , and  $\Gamma_0 = 7$  dB to yield BER  $= 10^{-3}$ [\[14](#page-10-0)]. The Rayleigh-fading random processes are generated according to Jake's model. Also, MS velocity 80 km/h; the step size of power adaptation  $\Delta_p = 1$  dB, the power-control period  $T_p = 2/3$  ms, and the power-control command loop delay  $T_d = 1 \cdot T_p$  are assumed for a realistic case. The antenna heights used with MS and BS are simply set to 0.003 and 0.01 times the BS radius, respectively, to prevent the occurrence of "zero distance" between an MS and a BS. The following figures show the numerical results of the performance of the ideal and realistic TCPC schemes, together with perfect power control [[4\]](#page-9-0), [[5\]](#page-9-0), combined power/rate control [[6\]](#page-9-0), and TAPC [[7\]](#page-9-0). Simulation results are also provided to validate the theoretical analysis of the TCPC scheme. Notably, the performance of the perfect power-control scheme is directly obtained from TCPC by setting  $X_0 = 0$ ; the cutoff thresholds for TAPC are applied such that the corresponding truncation probabilities are 10%.

For each scenario, the  ${L_{mb}}$  is randomly generated according to the channel model and the MS locations are uniformly distributed.  $I_O^* = I_O / \overline{P_C}$  can, thus, be calculated according to (8). Since TCPC is a type of strength-based power control, the distributions of  $Q_M[k]$  and PCE  $\psi$ , the average received power  $\overline{Q}$ , and the average MS transmission power  $\overline{P_C}$  are obtained by a bit-level single-link simulation with 10 000 power-control cycles. Based on the above factors, the outage probability can then be calculated. For each  $N_M$ , the

(32)

$$
\Lambda_M \approx \int_0^\infty \left( 1 - F_\xi \left( \frac{1.5G \cdot Q_0 \cdot y - \Gamma_0 \cdot (N_M - 1) \cdot \overline{Q}}{\Gamma_0 \cdot \overline{P_C} \cdot \sqrt{N_M}} - \frac{N_M \cdot \mu_O}{\sqrt{N_M}}, 0, \sigma_O^2 \right) \right) \cdot f_\psi(y) dy
$$
  
= 
$$
\int_0^\infty \int_y^\infty f_\xi \left( x; \mu_1, \sigma_1^2 \right) dx \cdot f_\psi(y) dy
$$
  
= 
$$
\int_0^\infty f_\xi \left( x; \mu_1, \sigma_1^2 \right) F_\psi(x) dx
$$



Fig. 2. System capacity C versus cutoff threshold  $X_0$ .



Fig. 3. System capacity C of various realistic power-control schemes versus cutoff threshold  $X_0$ .

overall outage probability is obtained by averaging the outage probability over 10 000 scenarios. Equivalently, the effect of MS movement on the outage probability is taken into account by considering many scenarios. The system capacity is then obtained. The other performance measures, including  $\overline{R}$ ,  $\overline{\mathfrak{E}}$ , and  $\overline{D}$ , are also obtained by a bit-level single-link simulation. Finally,  $\overline{\Re}$  is calculated.

Fig. 2 plots the system capacity  $C$  versus the cutoff threshold  $X_0$ . It can be seen that the differences between analytical results and simulation results are small, which indicates that the theoretical analyzes of the ideal and realistic TCPC schemes are quite accurate. Besides, it can be found that the ideal TCPC achieves the greatest system capacity over all the cutoff thresholds. This is mainly because TCPC suspends transmission as soon as the channel becomes bad, effectively reducing mutual interference. As  $X_0$  increases, the system capacity of the TCPC scheme rises at the cost of average transmission rate reduction and longer suspension delay, shown subsequently. TAPC has a low capacity because it suffers the short-term fading, although it employs a transmission-suspension mechanism to reduce mutual interference. Considering the case of  $X_0 = 0.7$  for instance, the ideal TCPC can accommodate 37 MSs, while the combined power/rate control, perfect power control, and TAPC at the truncation probability of 10% can only accept 20, 12, and 3 MSs,



Fig. 4. Average system transmission rate  $\overline{\Re}$  versus cutoff threshold  $X_0$ .



Fig. 5. MS average transmission rate  $\overline{R}$  versus cutoff threshold  $X_0$ .

respectively. Also, the realistic TCPC scheme yields a system capacity by  $78\% \sim 83\%$  lower than that of the ideal TCPC scheme, which reveals that the PCE indeed has a great impact on the system capacity of the TCPC scheme.

Fig. 3 further plots the system capacity  $C$  of various realistic power-control schemes versus the cutoff threshold  $X_0$ . These simulation results show that the system capacity of each realistic power-control scheme is tremendously reduced by the PCE, as compared with the results of the ideal power-control schemes shown in Fig. 2. Also, the TCPC again attains the greatest system capacity in the realistic situations.

Fig. 4 plots the average system transmission rate  $\Re$  versus  $X_0$ , where  $\overline{\Re}$  is normalized by the default rate  $R_0$ . The ideal TCPC scheme is found to have the highest  $\overline{\Re}$  value. In the example with  $X_0 = 0.7$ , the ideal TCPC  $\overline{\Re}$  is 133% of that of the combined power/rate control scheme, 183% of that of the perfect power-control scheme, and 815% of the TAPC scheme at the truncation probability of 10%. A larger  $X_0$  results in a higher  $\overline{\Re}$ . Notably,  $\overline{\Re}$  is a zigzag increasing function of  $X_0$ , not directly proportional to  $X_0$ . The stepwise constant capacity and declining MS average transmission rate cause the zigzag curves of the TCPC and combined power/rate control schemes in the figure.

Fig. 5 shows the MS average transmission rate  $\overline{R}$  versus  $X_0$ , where  $\overline{R}$  is also normalized by  $R_0$ . Of these power-control schemes, TCPC has the lowest  $\overline{R}$  due to the transmission



Fig. 6. MS average transmission energy per bit  $\overline{\mathfrak{C}}$  versus cutoff threshold  $X_0$ .



Fig. 7. MS average suspension delay  $\overline{D}$  versus cutoff threshold  $X_0$ .

suspension, while TAPC  $\overline{R}/R_0$  is 0.9, since its suspension probability is considered to be 10%. The low  $\overline{R}$  value of the TCPC scheme is a side effect. However, TCPC maintains the greatest average system transmission rate, as shown in Fig. 4. The side effect can be mitigated by either increasing the default rate  $R_0$  or by assigning more than one code channel to an MS. Notably, the analytical curve of the realistic TCPC scheme coincides with that of the ideal TCPC scheme due to the assumption made in Section III-B.

Fig. 6 shows the MS average transmission energy per bit  $\overline{\mathfrak{E}}$ versus  $X_0$ , where  $\mathfrak E$  is normalized by that of the perfect average power-control scheme. It can be found that TAPC saves the power most, followed by TCPC. The reason is intuitive: both schemes apply the transmission-suspension strategy. MS power consumption under perfect power control is highest, since the MS uses high power to compensate for deep fading. One extreme case is  $K = 1$ : the  $\overline{\mathfrak{E}}$  of the perfect power-control scheme is theoretically infinite, which can be obtained by letting  $X_0 \rightarrow 0$  in (18), causing infinite interference and zero capacity. Also, the higher the  $X_0$  is, the more power the TCPC scheme can save. For  $X_0 \geq 0.4$ , it can be seen that the simulation results can be well approximated by the analytical results, which supports the using (27) as the analytical formula of  $\mathfrak E$  for the realistic TCPC scheme in Section III-B.

Fig. 7 shows the MS average suspension delay  $D$  of the realistic TCPC scheme versus  $X_0$  for  $V = \{3, 20, 80, 120\}$  km/h.

<span id="page-9-0"></span>

A larger  $X_0$  induces a longer suspension delay, which is undesirable for real-time service. Such a suspension delay is another side effect of the TCPC scheme. However, once the delay requirement is stated, the upper limit of  $X_0$  for the best performance can be obtained. As shown in the figure,  $X_0$  should be less than 0.75 to guarantee that the suspension delay is less than 80 ms, as the velocity is greater than 3 km/h. The TAPC suspension delay would be much longer than that of the TCPC scheme, since the variation of long-term fading is much slower than that of short-term fading [8].

Fig. 8 plots the system capacity of the realistic TCPC scheme versus the MS velocity V at cutoff threshold  $X_0 = 0.7$  and step size  $\Delta_p = 1$  dB for power-control command loop delay  $T_d = 0$ ,  $1 \cdot T_p$ ,  $2 \cdot T_p$ . It is found that the factor of MS velocity greatly affects system performance. For the case of  $T_d = 1 \cdot T_p$ , the system capacity is 27 for  $V = 20$  km/h, which is 75% of the system capacity in the ideal case. The capacity drops rapidly and becomes only 5 for  $V = 120 \text{ km/h}$ . The performance degradation is even worse for  $T_d = 2 \cdot T_p$ . Although the case of  $T_d = 0$  is practically impossible, its corresponding performance curve serves as an upper bound for the realistic TCPC scheme. The space between the performance curves for  $T_d = 0$  and  $T_d = 1 \cdot T_p$  means the performance gain that a power-control scheme equipped with channel gain predictor to compensate the effect of power-control command loop delay could have.

# V. CONCLUSION

This paper analyzes the system performance of a TCPC scheme for uplinks in DS/CDMA cellular systems over frequency-selective fading channels. Based on the SIR formula over a medium-term period, closed-form formulas for five performance measures are successfully obtained, including system capacity, average system transmission rate, MS average transmission rate, MS average transmission energy per bit, and MS average suspension delay. This TCPC scheme is further analyzed under realistic conditions of PCE due to power-control command loop delay, step size for power adaptation, and MS velocity. Simulation results are also provided to validate the theoretical analysis. Results show that the analysis is accurate and that the TCPC scheme can greatly improve the system performance. This analytical approach can also apply to the conventional perfect power control by setting  $X_0 = 0$ . Notably, based on the SIR formula over the medium-term period, the analytical method presented herein can be easily extended for the study of another type of strength-based closed-loop power-control scheme in either circuit-switched or packet-switched CDMA systems.

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