

Flip-Flop Selection for Mixed Scan and Reset Design Based on Test Generation and Structure of Sequential Circuits*

HSING-CHUNG LIANG AND CHUNG LEN LEE⁺

*Department of Electronic Engineering
Chang Gung University
Taoyuan, Taiwan 333, R.O.C.*

E-mail: hcliang@mail.cgu.edu.tw

⁺*Department of Electronics Engineering*

*National Chiao Tung University
Hsinchu, Taiwan 300, R.O.C.*

E-mail: cllee@cc.nctu.edu.tw

In this paper, a novel mixed selection methodology using flip-flops for scan and reset design is proposed. The method runs test generation for a sequential circuit to obtain reachable states of flip-flops and required states for hard-to-detect faults. The circuit is also explored so as to acquire the structural connection relationship among the flip-flops. By analyzing these three sets of information, the flip-flops can be arranged in an appropriate order for mixed partial scan and reset selection. Instead of selecting the best flip-flop to revise the circuit for the next test generation, we give first priority to independent flip-flops each time in order to reduce the number of iterations. Experimental results show that this method can achieve higher testability with fewer scan/reset flip-flops than can either the scan only or the previous mixed scan/reset methods.

Keywords: partial scan, partial reset, reachable states, test generation, design for testability

1. INTRODUCTION

After decades of research, it is now generally recognized that to solve the testing problem, design for testability (DFT) must be incorporated into VLSI. For sequential circuits, scan design is the most widely used DFT method [1-32]. The full scan [3] can achieve very high testability, but it suffers from the disadvantages of excessive hardware overhead and lengthy test application time. Partial scan [4-32] and partial reset [33-36] have been considered to be more practical approaches to lessen the above problems while still maintaining a certain level of testability.

For partial scan, many methods have been proposed to select the minimal number of flip-flops so as to increase the maximal possible testability. Those methods can basically be classified as follows: (1) testability-measure-oriented [4-14]; (2) structure-analysis-oriented [7, 10, 13, 15-27]; and (3) test-generation-oriented [6, 10, 14, 28-31]. Each category of methods has its advantages and disadvantages. For example, calculating the testability may

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be very fast, but the measure may not be exact enough. Structural analysis is also fast, but hard-to-detect faults are not completely related to the way of circuit's connection. Test generation methods consider fault behavior directly but they take a long time or require too many iterations. Consequently, combinations of the above methods have also been considered [6, 7, 10, 13, 14, 32] in the hope of finding the perfect flip-flops so as to achieve a high level of testability.

Partial reset is an appealing technique in DFT because it increases the controllability of circuits with negligible modification of selected flip-flops and fewer additional inputs compared to scan design. Fig. 1 shows an example of partial reset design. The reset flip-flops can be forced into definite states at any time during the test pattern application. Obviously, the number of valid states of the circuit can be, therefore, increased, and so can the circuit's testability. Compared to partial scan, partial reset allows the circuit-under-test (CUT) to be tested at high speed with less test application time. The authors in [33] selected the reset flip-flops and their reset values by analyzing the state table of the circuit, those in [34] did so by breaking cycles and calculating the controllability of the circuit, and those in [35] did so by utilizing a cost function to conduct sensitivity analysis.

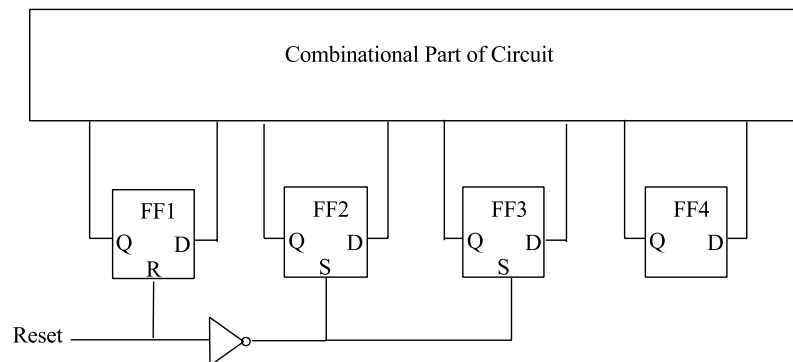


Fig. 1. An example of partial reset design. As the signal 'Reset' becomes 0, the outputs of flip-flops FF1, FF2, and FF3 are forced to be 0, 1, and 1, respectively.

Recently, it has been proposed that partial scan and partial reset be employed together in the CUT to enjoy the advantages of both approaches [36, 37]. Through using this approach, one can improve the level of controllability while achieving the needed level of observability of the CUT by using much less hardware as compared to scan only design. These methods select reset flip-flops first and then scan flip-flops in order to improve the circuit's testability with as few flip-flops as possible.

In this paper, a new method for selecting flip-flops for partial-scan and partial-reset design is proposed. Instead of selecting flip-flops first for reset and then for scan, we combine the steps to achieve better results. The method utilizes the following information to select the optimal flip-flops for partial scan and/or reset:

- (1) the reachable states obtained from test pattern simulation,
- (2) the fault-free states needed for excitation and propagation for hard-to-detect faults, and
- (3) the structural connection relationships among flip-flops.

By analyzing the above information, we can construct the appropriate selection sequence for flip-flops that are to be scanned or reset. The flip-flops which are predicted to improve testability more are given first priority for selection. Because choosing only one flip-flop each time for DFT requires many iterations, especially for large circuits, we simultaneously select those flip-flops which are *independent* in structure so as to reduce the required run time. Experiments on ISCAS89 [38] benchmark circuits show that the proposed method needs fewer flip-flops for scan or reset but achieves higher fault coverage than do the previous mixed scan/reset or scan only methods.

Section 2 describes how the three sets of information are collected. We explain how they are analyzed to guide flip-flop selection in section 3. Experimental results and conclusions are then given in sections 4 and 5, respectively.

2. INFORMATION FOR FLIP-FLOP SELECTION

The overall process involved in this mixed scan/reset flip-flop selection approach is shown in Fig. 2. As mentioned above, information on reachable states, the fault-free states needed for excitation and propagation of hard-to-detect faults, and the structural connection relationships among the flip-flops, is collected first. Consequently, at the beginning, a deterministic test generator is applied to the CUT to generate test patterns for the easy-to-detect faults in the circuit while using little computation time as possible. During this test generation step, the faults that the test generator identifies as being untestable and those for which it cannot find test patterns are called hard-to-detect faults. The states required to excite and propagate these faults to primary outputs are collected during the test generation step. The test patterns obtained by the test generator are then used to simulate the circuit in order to obtain the reachable states. At the same time, the circuit is analyzed in order to construct its *dependence graph* [39], which describes the structural relationships among the flip-flops. After analyzing the data to guide selection of the best flip-flop(s) for scan or

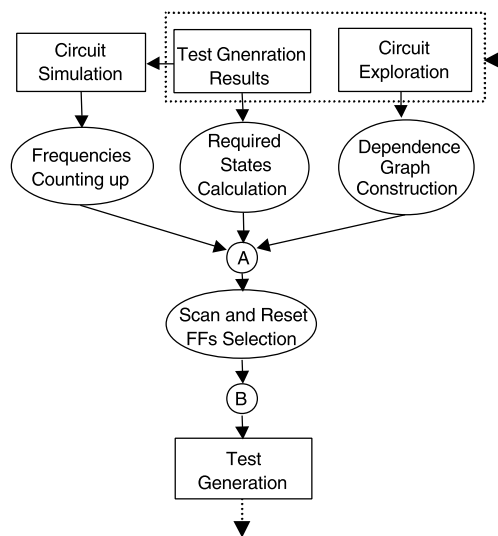


Fig. 2. The overall selection process for scan and reset flip-flops.

reset, we revise the circuit and run the test generation procedure for the subsequent sets of data. The process can be executed repeatedly until the obtained fault coverage satisfies the requirement. The details of the above data are described, respectively, in the following subsections.

2.1 Simulation for Reachable States

The reachable states are those states that can be reached from any state of a sequential machine. They represent the states that are ready for excitation and propagation of target faults for test generation. Obviously, the circuit can be much more testable if it has many valid states. Searching for all the valid states of a circuit is impractical, so we use reachable states that have been simulated based on test patterns generated using an implemented BACK[40,41]-like test generator. From these states, which are recorded in the set RS , we count up the 1 and 0 frequencies present on each flip-flop. These frequency counts indicate the difficulty involved in setting the flip-flops to be 0 and 1. The smaller the frequency, the more difficult it is to set the corresponding value. These count values, therefore, are a great aid in selecting flip-flops for partial scan/reset.

Example 1: Fig. 3 shows an example where the input patterns (010,110,101,001,111, ...) are used to simulate a circuit with three flip-flops. The reachable states obtained after this simulation is conducted are $RS(FF_1, FF_2, \text{ and } FF_3) = (X01,011,111,101, \dots)$. Then the frequencies of occurrence of 1 and 0 on these reachable state cubes are calculated. In Fig. 3, the frequencies of occurrence of 0 for FF_1 , FF_2 , and FF_3 are 1, 2, and 0, respectively. This indicates that the value 0 does not occur on flip-flop FF_3 until now, which implies that it is difficult for the circuit to set FF_3 to be 0 during test generation.

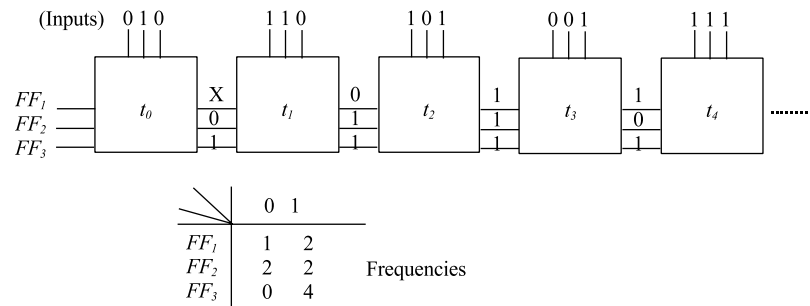


Fig. 3. An example showing the reachable states obtained through logic simulation. On the flip-flops, 'X' means don't care.

2.2 States Required for Excitation and Propagation for Fault Detection

To generate a test for a fault, it is necessary to excite and propagate the fault, perhaps through several time frames, to a primary output. During excitation and propagation, if it is difficult or impossible for the circuit to have the required states, then the target fault is likely to be aborted or judged untestable by the test generator. Accordingly, we will collect these needed states for hard-to-detect faults during the test generation phase in the set NS for further analysis.

Example 2: In Fig. 4, to excite the fault f , flip-flops FF_1 , FF_2 , and FF_3 must be set to (001). It takes two time frames to propagate the fault to an output. In the two time frames, FF_1 , FF_2 , and FF_3 are required to take values of (011) and (001), respectively. From the counting of value 1 and 0 in the states set, if f is a hard-to-detect fault, it is natural to select FF_1 and FF_3 as flip-flops to be reset to 0 and 1, respectively. In addition, as the fault passes through flip-flop FF_2 twice, FF_2 is better for being scanned.

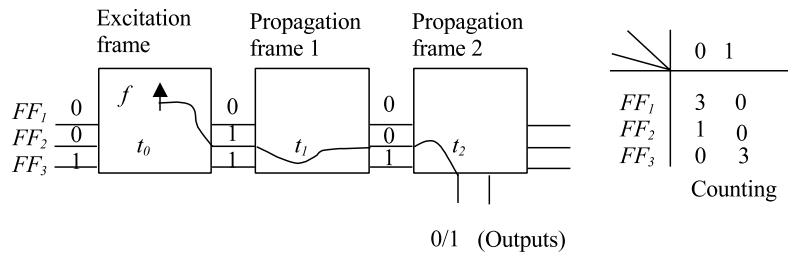


Fig. 4. An example showing the excitation and propagation states needed for a hard-to-detect fault. The frequencies of occurrence of 1 and 0 on flip-flops reveal which flip-flops are good/bad candidates for selection for partial scan or reset.

From the example, it is obvious that we need to record the required states (NS) for excitation and propagation of hard-to-detect faults and also the frequencies (FP) of the flip-flops for the faults to pass through. These data are also good indices for selecting candidate flip-flops for partial scan/reset.

2.3 Structural Relationships Among Flip-Flops

The structural relationships among flip-flops in a circuit play an important role in determining the testability of the circuit. Here a *dependence graph* [39] is used to describe the structural relationships among flip-flops. Fig. 5 shows the *dependence graph* for the ISCAS benchmark circuit s400 [38]. In the graph, each node represents a flip-flop or a group of flip-flops, which themselves are in a strongly connected graph, and each branch represents a signal path from a source node to a destination node. In this graph, each node is leveled such that flip-flops in lower levels determine the values of flip-flops in higher levels. The graph provides information about how the signal on one flip-flop may be affected by those on some other flip-flops, and about which flip-flops the signal will be able to go through. This information will also be helpful in selecting reset and scan flip-flops.

3. FLIP-FLOP SELECTION FOR SCAN AND RESET

After the above information has been collected, it is then analyzed to help select flip-flops for partial scan or partial reset. The steps are shown in detail in Fig. 6, which shows a sub-process of the process shown in Fig. 2. From the states in the reachable states set RS , we accumulate the occurrence frequencies for 1 and 0 on flip-flop FF_n and record them as $RS_1(n)$ and $RS_0(n)$, respectively. The smaller the values of $RS_1(n)$ or $RS_0(n)$ are, the more difficult it will be to set the flip-flop to “1” or “0”, respectively. If $RS_1(n)$ is 0 (or very small) and $RS_0(n)$ is very large, FF_n is a good candidate for being reset to 1 rather than being

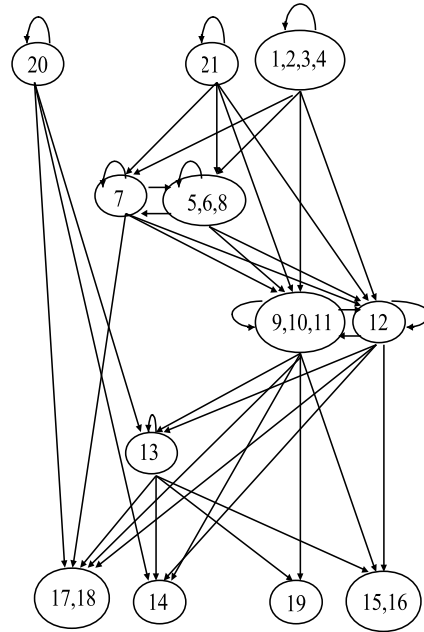


Fig. 5. The dependence graph of circuit s400.

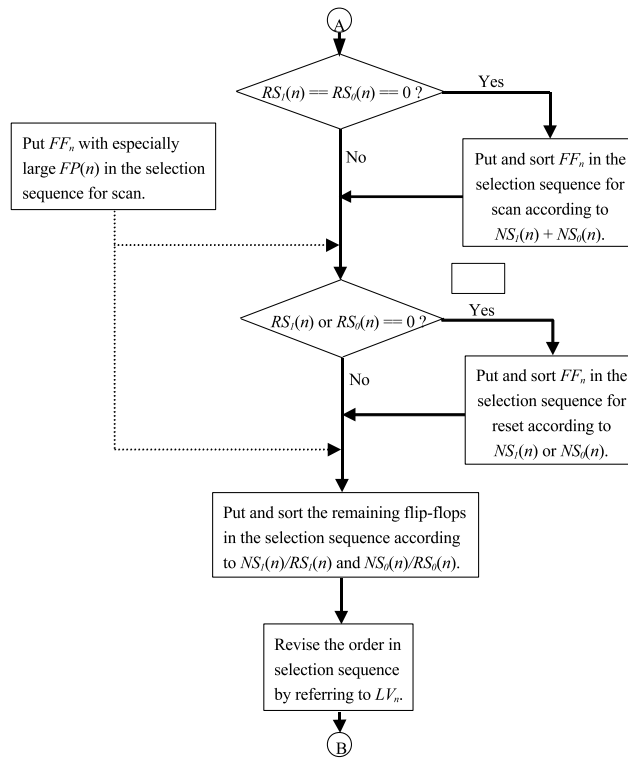


Fig. 6. The sub-process for selecting flip-flops for partial scan/reset.

scanned. If both $RS_1(n)$ and $RS_0(n)$ are equal to 0 (or very small), then FF_n is a good candidate for scanning. When there is more than one flip-flop for which both $RS_1(n)$ and $RS_0(n)$ are equal to 0, we need more information to help distinguish the importance of these flip-flops in order to improve the level of testability.

As mentioned in section 2.2, we collect the needed states for excitation and propagation of hard-to-detect faults in the states set NS . Let $NS_1(n)$ and $NS_0(n)$ be the frequency of occurrence of 1 and 0, respectively, on flip-flop FF_n in NS . The larger the values of $NS_1(n)$ and $NS_0(n)$, the more often the values 1 and 0 are required on flip-flop FF_n , respectively, during test generation. For those flip-flops for which both $RS_1(n)$ and $RS_0(n)$ are equal to 0, we add the values of $NS_1(n)$ and $NS_0(n)$ for each flip-flop and sort these flip-flops according to the summation in descending order. These flip-flops are first selected for partial scan as they are difficult to be set values, i.e., either 1 or 0, but are urgently required for test generation.

Example 3: Assume a circuit with six flip-flops has the data $RS_1(n)$, $RS_0(n)$, $NS_1(n)$, and $NS_0(n)$ as shown in the following:

FF_n	$RS_1(n)$	$RS_0(n)$	$NS_1(n)$	$NS_0(n)$
1	0	0	17	22
2	16	34	20	42
3	0	0	48	31
4	0	50	29	19
5	50	0	10	42
6	28	22	15	24

Because flip-flops FF_1 and FF_3 have RS_1 and RS_0 values both equal to 0, we calculate the summation of NS_1 and NS_0 for these flip-flops and obtain the values $17 + 22 = 39$ and $48 + 31 = 79$, respectively. Consequently, we put FF_3 first in order and FF_1 second for scanning.

For the next consideration, the flip-flops for which only $SS_1(n)$ or $SS_0(n)$ are equal to 0 are put into the selection sequence. They are considered for resetting and sorted in order according to their $NS_1(n)$ or $NS_0(n)$ values. The larger the $NS_1(n)$ (or $NS_0(n)$) is, the more urgent it is for FF_n to be reset to 1 (or 0).

Example 4: In *Example 3*, $RS_1(4)$ and $RS_0(5)$ are equal to 0 but $RS_0(4)$ and $RS_1(5)$ are not. Therefore, they will be selected for resetting rather than scanning. Because $NS_1(4) = 29 < NS_0(5) = 42$, flip-flop FF_5 , which is to be reset to 0, will be selected before FF_4 , which is to be reset to 1.

In section 2.2, we explained that besides the required states set NS , we also collect the frequencies of passing through the flip-flops for the hard-to-detect faults during propagation. Let $FP(n)$ be the frequency of faults passing through flip-flop FF_n . A flip-flop with a large $FP(n)$ is also a good selection for scanning because it can let many hard-to-detect faults become easily observed. In our experiment, we put the flip-flops whose $FP(n)$ s were at least twice the values of the others next in the selection sequence. It is noteworthy that as scan design requires much more hardware overhead and test application time than reset does, selecting flip-flops with respect to $FP(n)$ is usually considered behind the previous selection for reset. However, the value of $NS_1(n)$ or $NS_0(n)$ for reset selection may sometimes be very small; i.e., selecting the corresponding flip-flop for reset will not increase the circuit's testability very much, so the flip-flops with large $FP(n)$ values will be considered first to achieve better testability.

Example 5: Assume that the circuit in *Example 3* has the $FP(n)$ values for the six flip-flops shown in the following table. Obviously, FF_3 will be considered for scanning as it has a larger FP value compared to the others.

FF_n	$FP(n)$
1	14
2	2
3	57
4	19
5	5
6	11

For the remaining flip-flops, neither $RS_1(n)$ nor $RS_0(n)$ being equal to 0, which implies that they may not be suitable for resetting because the resulting circuit may only increase a few new valid states. More analysis of reachable states and required states is needed to find the possibility of increasing valid states if resetting these flip-flops. For simplicity, in our process, they are chosen for scanning according to the values of $NS_1(n)/RS_1(n)$ and $NS_0(n)/RS_0(n)$, which are accounted for because the flip-flop with larger $NS_1(n)$ (or $NS_0(n)$) and smaller $RS_1(n)$ (or $RS_0(n)$) is more demanding than those with the inverse condition.

Example 6: For the data given in *Example 3*, flip-flops FF_2 and FF_6 have RS_1 and RS_0 values which are not equal to 0. For flip-flop FF_2 , $NS_1(2)/RS_1(2) = 20/16 = 1.25$ and $NS_0(2)/RS_0(2) = 42/34 = 1.24$. For flip-flop FF_6 , $NS_1(6)/RS_1(6) = 15/28 = 0.54$ and $NS_0(6)/RS_0(6) = 24/22 = 1.09$. Having larger values, FF_2 will therefore be placed in the selection sequence before FF_6 . Both flip-flops are for scan only when being chosen to aid the design for testability.

To find the fewest flip-flops for DFT to improve the highest testability, it is theoretical to choose the best flip-flop each time and to then obtain new data for the revised circuit for the next round of analysis. The iterations are continued until the required testability for the circuit is achieved. Analyzing the data is fast, but running test generation takes much time. To quickly obtain an enough testable circuit, we apply a strategy of selecting *independent* flip-flops each time. Here *independent* means that these flip-flops do not affect each other in the dependence graph; i.e., there is no directed path from one flip-flop to another in the dependence graph. For two flip-flops in a directed path in the graph, selecting one flip-flop for scanning may affect the testability of the other one; therefore, they are not suitable for simultaneous selection for DFT purposes. As for independent flip-flops, if they are adjacent in the flip-flop selection sequence, they can be selected at the same time for DFT in order to reduce number of iterations required for test generation.

Example 7: Assume a circuit with six flip-flops has the dependence graph shown in Fig. 7. Flip-flops FF_1 and FF_2 are independent because there are no directed paths between them. If after analyzing the data for scanning and resetting, we find that they are placed in adjacent positions in the selection sequence, then they can be selected together for DFT purposes in the same iteration.

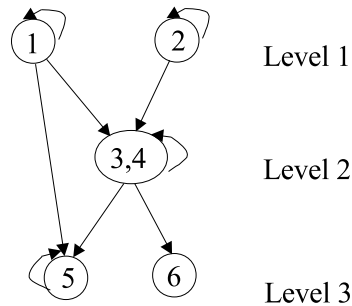


Fig. 7. A dependence graph of a circuit with six flip-flops.

In Fig. 7, the dependence graph has three levels. The values of the flip-flops in the lower level implicitly determine those of the flip-flops in higher levels. Consequently, the flip-flops in the lower level are likely to be selected first for resetting or scanning because they may affect the testability of the flip-flops in higher levels. Let LV_n be the level of FF_n in the dependence graph. For the flip-flops with higher LV_n 's, they affect less number of flip-flops if being reset and behave more likely to be only new primary outputs if being scanned. Accordingly, when two flip-flops are adjacent in the selection sequence and have comparable values for resetting or scanning, their levels in the dependence graph are used to slightly modify the selection order.

Example 8: In Fig. 7, assume that we obtain a selection sequence in which flip-flops FF_6 and FF_3 have similar values and FF_6 is prior to FF_3 . Their positions will be exchanged according to the consideration for levels in the dependence graph.

Example 9: We will use circuit s400 [38], whose dependence graph is shown in Fig. 5, to explain the entire process of selecting flip-flops. After performing test generation, we found that none of the $RS_1(n)$ and $RS_0(n)$ values of the flip-flops were equal to 0 based on the reachable states of test patterns. There was also no flip-flop with an especially large $FP(n)$ value. Therefore, we compared the flip-flops with $NS_1(n)/RS_1(n)$ and $NS_0(n)/RS_0(n)$. From the comparison, we found that flip-flop FF_{12} had the largest values and, hence, was a good selection for scanning. The scanned circuit was then subjected to test generation again. None of the new $RS_1(n)$ and $RS_0(n)$ of all the flip-flops were equal to 0, either. There was still no flip-flop with an especially large $FP(n)$. According to the $NS_1(n)/RS_1(n)$ and $NS_0(n)/RS_0(n)$ of the remaining flip-flops, we found the next flip-flop, i.e., FF_8 , for being selected for scan. The process can be continued until the required fault coverage is achieved. It will be shown in the next section that, for circuit s400, we only selected two flip-flops for scanning but obtained fault coverage comparable to that of the other two scanning methods, which required three and five flip-flops, respectively. The *fault efficiency*, i.e., the sum of the counts of detectable and identified untestable faults divided by the total number of faults for this circuit appears to be 100% for the test generation of all the methods. This means that the reported levels of fault coverage are the highest achievable results for the three selection methods. Therefore, our method requires fewer scan flip-flops but achieves the same level of testability as do the other methods for circuit s400.

4. EXPERIMENTAL RESULTS

The above procedure, called RESCAN, was implemented, on a PC with a Pentium II 300Mhz CPU and 256M RAM. A BACK[40,41]-like test generator was implemented and used for test generation in the procedure. Table 1 shows the results obtained using RESCAN compared to those of three other recently reported methods for partial scan [14, 31, 32]. In

Table 1. Comparison of RESCAN with the other three methods of scan selection. In the table, ‘-’ means no data given, and ‘*’ means that the original circuits are at least different in terms of the number of flip-flops.

Circuit	BELLONA[31]			Opscan[32]		IDROPS[14]		RESCAN			
	#FF	#SFF	FC(%)	#SFF	FC(%)	#SFF	FC(%)	#SFF	#RFF	FC(%)	#Iter
s298	14	1	94.8	1	94.8	-	-	1	0	98.05	1
s344	15	1	96.2	3	98.8	-	-	1	0	98.54	1
s349	15	2	98.0	3	98.3	-	-	1	0	98.00	1
s382	21	3	97.2	5	97.5	-	-	3	0	97.49	3
s386	6	-	-	2	92.2	-	-	2	0	94.79	2
s400	21	3	95.8	5	95.8	-	-	2	0	95.75	2
s420	16	-	-	-	-	3	20.9	0	3	50.23	3
s444	21	2	94.5	5	94.9	3	93.2	2	0	93.25	1
								3	0	94.73	2
s526	21	2	91.4	7	98.7	3	87.2	2	0	93.69	2
s526n	21	-	-	8	99.1	-	-	6	0	99.10	3
s641	19	1	95.7	5	94.2	-	-	0	1	99.36	1
s713	19	1	87.4	5	88.1	-	-	0	1	92.97	1
s820	5	1	98.9	2	100.0	-	-	1	0	100.00	1
								0	1	99.53	1
s832	5	1	97.7	2	98.4	-	-	1	0	98.39	1
s953	29	-	-	3	100.0	-	-	3	0	100.00	2
s1423	74	34	97.6	-	-	15	95.8	26	2	95.27	7
s1488	6	1	99.1	2	100.0	-	-	1	0	99.87	1
s1494	6	1	98.3	3	99.2	-	-	1	0	99.00	1
s5378	179	27	93.8	80	97.5	21	94.7	15	8	94.91	10
s9234	228	-	-	-	-	*27	*79.9	41	13	73.74	9
s13207	669	-	-	-	-	*78	*76.5	82	4	77.15	7
s15850	597	-	-	-	-	*66	*65.2	69	7	65.58	4
s35932	1728	-	-	150	89.8	-	-	0	8	89.72	1

the table, column #FF lists the number of flip-flops in each circuit, and column #SFF gives those selected for scan. Column #RFF lists the number of flip-flops selected for reset design. Column FC(%) shows the final fault coverage of the circuits. For RESCAN, column #Iter gives the number of iterations required to select flip-flops. As shown in the table, BELLONA[31] and OpSCAN[32] provided most of the circuits except s5378 with 100% *fault efficiency*. In other words, the fault coverage of most of the circuits listed in the table is the highest achievable using the two methods. However, our method can achieve even higher fault coverage for many circuits by selecting fewer flip-flops for DFT. The authors of IDROPS[14] used a simulation-based test generator[42] which could not identify untestable faults, so that they could not achieve the highest fault coverage for the circuits. The authors also did not provide results for many circuits. In addition, they did not explain why the number of flip-flops for circuits s9234, s13207, and s15850 was 211, 638, and 534, respectively, which results are different from those provided by ISCAS89[38]. From the table, RESCAN exhibits better results with fewer scanned flip-flops but higher fault coverage for most of the circuits compared to the other methods. For some circuits, RESCAN can even use flip-flops for resetting only yet can achieve higher fault coverage than the other three methods using scanning. Some circuits, e.g. s420 listed, in the table can be further run to achieve 100% fault efficiency or coverage. For the larger circuits, iterations of test generation take much time, which is the common disadvantage of flip-flop selection methods that use test generation results. Due to the limitations of our test generator, the real fault coverage for each of the large circuits can be even higher than shown. As the run time for the other methods is not available for comparison, we only show the time required by our test generator in Table 2 for some of the circuits. For the circuits that are not shown in the table, their test generation jobs were usually interrupted for various reasons, which prevented us from conducting complete runs and recording the time used. We estimated that our test generator took more than two weeks to run iterations for each circuit. If available, more efficient test generators may be used to speed up collection of needed data. Another suggested way is to first find test patterns for the easily testable faults only. The remaining faults may all be considered as difficult to detect, and their required states for fault excitation and propagation can then be quickly found by means of a simplified test generation process. The test generator need not take time to ensure whether there are test patterns or not. The process of collecting data for analysis can then be speeded up, but the obtained data may not be accurate enough for selecting the correct flip-flops.

Table 3 lists the results obtained using RESCAN compared with those of two other methods that incorporated both partial reset and partial scan [36, 37]. The two methods both applied reset selection first and then scan selection for flip-flops, which is different from our mixed selection approach. They need some criteria to decide how many reset flip-flops are sufficient. In our approach, we choose the best flip-flop(s) for scanning or resetting according to whether they can make most of the hard-to-detect faults testable. The information about the reachable states, the states required for fault excitation and propagation, the number of hard-to-detect faults passing through each flip-flop, and the structural dependence graph enable our approach to select flip-flops more accurately. As for MIXPART [36], in addition to partial scan and reset, it employs a partial observation technique that makes some selected flip-flops observable. However, RESCAN out-performs it significantly on many circuits. Comparing to [37], RESCAN still exhibits better results for many circuits. As for circuits larger than s5378, because our test generator cannot achieve 100% fault efficiency, we believe that the real fault coverage can be even higher if better test generator is used.

Table 2. Run time of RESCAN on some benchmark circuits.

Circuit	#FF	#SFF	#RFF	FC(%)	#Iter.	Time(sec.)
s298	14	1	0	98.05	1	1.6
s344	15	1	0	98.54	1	189.1
s349	15	1	0	98.00	1	45.2
s382	21	2	0	95.24	3	53.6
s386	6	2	0	94.79	2	0.8
s400	21	2	0	95.75	2	120.8
s420	16	0	3	50.23	3	2.3
s444	21	3	0	94.73	2	453.1
s526	21	2	0	93.69	2	1125.6
s526n	21	6	0	99.10	3	49.6
s641	19	0	1	99.36	1	1.2
s713	19	0	1	92.97	1	1.4
s820	5	1	0	100.00	1	1243
s832	5	1	0	98.39	1	5947
s953	29	3	0	100.00	2	299
s1488	6	1	0	99.87	1	50815
s1494	6	1	0	99.00	1	67961

5. CONCLUSIONS

In this work, a novel methodology for flip-flop selection for partial resetting and partial scanning on sequential circuits designed to improve testability has been proposed. Previous methods which combined scanning and resetting selected flip-flops first for resetting and then for scanning. The proposed method is not limited to this selection approach but instead finds the flip-flops that most improve the testability improvement if selected for scanning or resetting. To find the best flip-flops, the procedure collects information about the reachable states of test patterns and states required for hard-to-detect faults through test generation and simulation, and explores for structural inter-connection information of flip-flops. The obtained data are analyzed to produce a selection sequence for all the flip-flops being reset or scanned. With respect to the highest priority flip-flop(s), the process revises the circuit with scan or reset and subjects the resulting circuit to test generation again. This can be repeated until the required fault coverage is achieved. Experimental results on ISCAS benchmark circuits show that the methodology requires fewer flip-flops for scanning and resetting to achieve a higher level of testability than can be achieved by the other scan only and mixed scan/reset methods. This proves the better accuracy achieved by our method in selecting flip-flops. In order to reduce the time needed, superior test generators are needed to speed up the process of searching for test patterns and states required for selecting flip-flops.

Table 3. Comparison of RESCAN with two other methods that combine reset and scan.

Circuit	MIXPART[36]				[37]			RESCAN		
	#FF	#RFF	#SFF	FC(%)	#RFF	#SFF	FC(%)	#RFF	#SFF	FC(%)
s298	14	2	1	95.37	2	1	95.54	0	1	98.05
s344	15	4	2	98.89	3	1	99.14	0	1	98.54
s349	15	4	3	98.37	3	1	98.60	0	1	98.00
s382	21	6	3	98.58	2	3	97.28	0	3	97.49
s386	6	2	2	98.21	1	2	98.96	0	2	94.79
s400	21	6	2	93.29	2	1	94.42	0	2	95.75
s420	16	4	5	63.17	3	1	72.15	3	0	50.23
s444	21	6	5	95.98	4	5	96.27	0	3	94.73
s526	21	6	6	97.75	2	5	98.40	0	2	93.69
s526n	21	–	–	–	2	5	98.57	0	6	99.10
s820	5	2	2	99.77	1	1	100.00	0	1	100.00
s832	5	2	2	98.18	1	1	98.39	0	1	98.39
s953	29	6	2	98.29	2	1	98.53	0	3	100.00
s1423	74	16	25	93.65	4	25	94.89	2	26	95.27
s1488	6	2	2	99.40	2	2	99.66	0	1	99.87
s1494	6	2	3	98.88	2	2	98.88	0	1	99.00
s5378	179	36	15	94.67	8	15	92.34	8	15	94.91
s9234	228	94	70	89.99	51	70	90.20	13	41	73.74
s13207	669	–	–	–	–	–	–	4	82	77.15
s15850	597	–	–	–	–	–	–	7	69	65.68
s35932	1728	16	200	90.19	16	200	89.80	8	0	89.72

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Hsing-Chung Liang (梁新聰) received the B.S., M.S., and Ph.D. degrees in electronics engineering from National Chiao Tung University, Taiwan, R.O.C., in 1989, 1991, and 1997, respectively. From 1997 to 1999, he was an Assistant Professor in the Department of Electronic Engineering, Van Nung Institute of Technology. Since August 1999, he joined the faculty of the Department of Electronic Engineering, Chang Gung University. His research interests include VLSI testing and design for testability.



Chung Len Lee (李崇仁) obtained his B.S. from National Taiwan University in 1968 and M.S. and Ph.D. from Carnegie Mellon University in 1971 and 1975 respectively, all in Electrical Engineering. He has been with Department of Electronics Engineering, National Chiao Tung University, since 1975, engaging in teaching and research in the fields of semiconductor devices, integrated circuits, VLSI, computer aided design and testing. He has supervised over 120 M.S. and Ph.D. students to complete their thesis and has published over 250 papers in the above areas. He has been involved in various technical activities in the above areas in Taiwan as well as in Asia. He is on the editorial board of JETTA.