

Half-band filter design with odd-polynomial Chebyshev approximation

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Abstract: An efficient method for the design of the half-band filter is described. Both the characteristics of the time response and the frequency response of the half-band filter are fully utilised to reduce the computation time and to increase the computation accuracy. The algorithm is based on the Chebyshev approximation criterion and uses only odd-symmetry polynomials for frequency interpolation to meet the required frequency response of the half-band filter.

1 Introduction

The use of digital technique in signal processing has become very popular in many areas. In several applications, a system may even be required to handle data of different type with various sampling rates [1]. This raises the need for a rate conversion system. A tree structure composed of sections each capable of coping with a sampling rate change of two is currently the most widely used. The design of half-band filter for decimation and interpolation has received considerable attention in the past decade [4-8].

Mintzer [2] had shown that the half-band filter can be designed by the Parks and McClellan [12] procedure with symmetry constrain on the passband and the stopband cutoff frequencies as well as with equal weighting on the passband and the stopband error. The filter thus designed is not optimum because of the quantisation error. Grenez [3] proposed a linear programming method with constrained Chebyshev approximation to design the half-band filter. The dense grid of points replacing the frequency axis is made symmetric with respect to $\pi/2$. The specifications of the half-band filter was exactly met with the Remez exchange algorithm [20]. The symmetry property of the extremal points is not fully utilised in the general Parks and McClellan procedure. The search for the new external points in the Remez algorithm is thus not optimised. Vaidyanathan and Nguyen [19] proposed another method which exploits knowledge about the coefficient of the impulse response of the half-band filter to reduce the design time.

An efficient design procedure based on the odd-polynomial interpolation technique for the half-band

filter is presented. The frequency response is always guaranteed to be symmetric with respect to $w = \pi/2$. The computation time is about 3/8 of that of the P-M method in each iteration because of the use of only odd-coefficients. The resultant filter is guaranteed to be a half-band filter. A filter obtained using this procedure has the following properties:

- (a) The even terms of the impulse response are zero
- (b) The searching time for the new external points can be reduced by one-half compared with the Parks-McClellan method
- (c) The external points obtained will be symmetrical with respect to the quarter-Nyquist frequency
- (d) The overall computation time can be reduced by at least 3/8 compared with that of the Parks-McClellan method and one quarter is found in the implementation.

2 Characteristics of half-band filter

Multirate digital signal processing concepts have recently become popular. The basic concept of multirate digital signal processing is that the sampling rate of signals to be processed in the same system can be different. The signal is decomposed in the frequency domain into different channels to reduce the bandwidth. The signal in each channel can be decimated to lower rates because of the reduction in bandwidth. The original signal can be reconstructed without error if these subsignals are properly recombined through interpolation.

If the factor of decimation and interpolation is two, a 2-band QMF bank will result. The system function of this filter bank is

$$X(Z) = \frac{1}{2}[H_0(Z)G_0(Z) + H_1(Z)G_1(Z)] X(Z) + \frac{1}{2}[H_0(-Z)G_0(Z) + H_1(-Z)G_1(Z)] X(-Z) \quad (1)$$

The term containing $X(-Z)$ is caused by the aliasing effect.

Smith and Barnwell [7, 8] proposed a method for the choice of the analysis and synthesis filters to reach the perfect reconstruction condition. The choices are

$$\begin{aligned} H_1(Z) &= Z^{-(N-1)}H_0(-Z^{-1}) \\ G_1(Z) &= Z^{-(N-1)}H_1(Z^{-1}) \\ G_0(Z) &= Z^{-(N-1)}H_0(Z^{-1}) \end{aligned} \quad (2)$$

Using this, the transfer function can be written as

$$T(Z) = F(Z)Z^{-(N-1)} + F(-Z)Z^{-(N-1)} \quad (3)$$

where $F(Z) = H_0(Z)H_0(Z^{-1})$. For perfect reconstruction, the transfer function is required to be $Z^{-(N-1)}$. From eqn.

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3, it is necessary that

$$F(Z) + F(-Z) = 1 \quad (4)$$

$F(Z)$ is a linear phase FIR filter over the interval $[-(N-1), N-1]$ and is called the half-band filter. From eqn. 4, it can be shown that the even terms of $f(n)$ are zero except for $f(0) = 0.5$. If expressed in spectral domain

$$F(\omega) + F(\pi - \omega) = 1 \quad (5)$$

and

$$F(\omega) = |H_0(\omega)|^2 \quad (6)$$

From eqn. 5, it is clear that $F(\pi/2) = 0.5$. By this, eqn. 5 can be written as

$$F(\omega) - F\left(\frac{\pi}{2}\right) = -\left[F(\pi - \omega) - F\left(\frac{\pi}{2}\right)\right] \quad (7)$$

This means that $F(\omega)$ is odd-symmetry with respect to $\omega = \pi/2$. The design of this odd-symmetry filter is the main topic of this paper.

The P-M [12] algorithm is a popular method for linear phase FIR filter design. This method uses Chebyshev approximation for frequency interpolation. It uses the Remes exchange algorithm to search external points during the iteration. The half-band filter is known to be an extra-ripple filter and is odd-symmetry with respect to $\omega = \pi/2$, i.e. even terms must be zero. The P-M method can be used to design such a filter with equal weighting on both the pass band and the stop band and with symmetrical frequency response with respect to $\pi/2$. The even terms are not usually zero because of the quantisation error in the calculation. The filter obtained is not the half-band filter required. If the even term is deliberately set to zero to meet the requirement of the half-band filter, the resultant error ripple is not greater than that of the original [2]. This phenomenon indicates that a filter designed by the conventional P-M method is not optimal. The results from both the P-M method and the proposed odd-polynomial based method for the half-band filter design are to be compared.

The frequency response of the half-band filter is odd-symmetrical, so the search time for extreme points in the Remes exchange algorithm is reduced by half in the proposed method. The even terms of the time response are assumed to be exactly zero, so the size of the interpolation basis in the formulation is reduced by half. The proposed method is discussed in detail and the design formula is derived from the point of view of linear algebra.

3 Chebyshev approximation based on odd-polynomial

An efficient technique to design the halfband filter with maximum stop band attenuation is presented. The filter spectral response is interpolated by a set of odd-polynomials. The Remes exchange algorithm is used to search the new extreme points in the Chebyshev approximation.

3.1 Chebyshev approximation

Chebyshev approximation [9] has long been used for optimisation designs such as network synthesis [10]. The application of this in linear phase FIR filter design was proposed by Parks and McClellan [12]. Many techniques [11, 12, 17, 18] exist that can be used to solve this kind of optimisation problem. Linear programming and

the second Remes algorithm (the so-called Remes exchange algorithm) are the two most popular methods for solving the Chebyshev sensed optimisation problems.

The Remes exchange algorithm has fast convergence speed but is difficult to use or implement. Linear programming is easier and more flexible in use, but it is much slower in convergence than the Remes algorithm. The Remes exchange algorithm for optimisation in Chebyshev approximation is discussed.

3.1.1 Problem modelling

Let $\{h(n)\}$ be the impulse response of the product filter over the interval $[-(N-1), (N-1)]$. The frequency response of the product filter can be described as

$$H(\omega) = h(0) + \sum_{n=1}^{N-1} 2h(n) \cos(n\omega) \quad 0 \leq \omega \leq \pi \quad (8)$$

The Chebyshev approximation for the FIR filter design could be modelled as

$$\min \left(\max_{\omega \in A} W(\omega) |H(\omega) - D(\omega)| \right) \quad (9)$$

where $W(\omega)$ is the weighting function of the approximation error, $D(\omega)$ is the desired frequency response and A is a compact subset within $[0, \pi]$.

If the weighted error function of the approximation is written as

$$E(\omega) = W(\omega)[H(\omega) - D(\omega)] \quad \omega \in A \quad (10)$$

then the Chebyshev approximation can be rewritten in terms of $E(\omega)$, as

$$\min \left(\max_{\omega \in A} |E(\omega)| \right) \quad (11)$$

3.1.2 Remes exchange algorithm

The basic concepts of Chebyshev approximation for filter design is as mentioned above. An iterative algorithm to solve this optimisation problem is discussed — the well known Remes exchange algorithm. This algorithm is simple and efficient, and is known [13, 14] to converge uniformly according to $|P^k - P^*| < C\theta^k$ and all components of the approximation must be continuous in A . P^k is the k th approximation polynomial, P^* is the best approximation polynomial, C is a constant and $0 < \theta < 1$.

Several conditions and properties are first stated.

Haar condition

Let $\{g_1, g_2, \dots, g_n\}$ be a prescribed function set defined on a compact metric space A . The system of functions $\{g_i\}$ is said to satisfy the Haar condition iff every $g_i(x)$ is continuous in A for $i = 1, \dots, n$, and for every set of n -vectors $\{\bar{x}_i\}$ with the form $\bar{x}_i = [g_1(x_i), \dots, g_n(x_i)]$, $i = 1, \dots, n$, is linear independent for all distinct $x_i \in A$. That is, $\{g_1, g_2, \dots, g_n\}$ forms a basis of the n -space over A .

The linear combination $\sum_i c_i g_i$ is defined as the generalised polynomial. From the Haar condition, several properties can be obtained.

Property 3.1.1

If the set of functions $\{g_1, g_2, \dots, g_n\}$ satisfy the Haar condition. There exists a square matrix M , such that

$$M(x_1, x_2, \dots, x_n) = \begin{pmatrix} g_1(x_1) & g_2(x_1) & \cdots & g_n(x_1) \\ g_1(x_2) & g_2(x_2) & \cdots & g_n(x_2) \\ \vdots & \vdots & \ddots & \vdots \\ g_1(x_n) & g_2(x_n) & \cdots & g_n(x_n) \end{pmatrix} \quad (12)$$

$M(x_1, x_2, \dots, x_n)$ is clearly nonsingular, if the samples $\{x_i\}, i = 1, \dots, n$, are all distinct over A .

Property 3.1.2

Let $\{g_1, \dots, g_n\}$ satisfy the Haar condition over A and span the n -space V^n , i.e., $\forall g \in V^n, \exists! \{a_i\}, i = 1, \dots, n, \ni g = \sum_i a_i g_i$. If there is another function g_{n+1} which is continuous on A . Thus $\{g_1, \dots, g_n, g_{n+1}\}$ would also satisfy the Haar condition if $g_{n+1} \notin V^n$.

Alternation theorem

Let $\{g_1, g_2, \dots, g_n\}$ be a system of functions on A satisfying the Haar condition. For a given function $f(x), x \in A$, there is one generalised polynomial $P = \sum_i c_i g_i$ that would be the best approximation of f on A iff the error function $r = f - P$ exhibits at least $n + 1$ 'alternations' on A . In other words, $r(x_i) = -r(x_{i+1}) = \pm \rho, x_0 < x_1 < \dots < x_n, x \in A$.

A means the interested regions of x and from this theorem, A must be chosen to have $n + 1$ alternations or more, and $\rho = \max_{x \in A} |r(x)|$.

Iterative procedure

Let the functions $\{g_i\}$ satisfy the Haar condition, and $r(x) = W(x)|f(x) - \sum_i c_i g_i|$ is the weighted error function satisfying the alternation theorem on A . $f(x)$ and $W(x)$ are continuous on A . Let $\{x_0^k, x_1^k, \dots, x_n^k\}$ be the k th iteration extreme points, $x_0^k < x_1^k < \dots < x_n^k, x_i^k \in A$, and ρ^k the deviation of the k th iteration according to the k th extreme points.

The solution of the Chebyshev approximation problems can then be obtained by the following iterative process:

- (a) Initial guessing of the extreme points $\{x_i^0\}$
- (b) From alternation theorem, suppose in the k th iteration $\{x_i^k\}$ and ρ^k satisfy

$$W(x_i^k) \left[f(x_i^k) - \sum_{j=1}^n c_j^k g_j(x_i^k) \right] = (-1)^i \rho^k \quad i = 0, \dots, n \quad (13)$$

in matrix form

$$\begin{pmatrix} g_1(x_0^k) & \dots & g_n(x_0^k) & \frac{-1}{W(x_0^k)} \\ g_1(x_1^k) & \dots & g_n(x_1^k) & \frac{1}{W(x_1^k)} \\ \vdots & \ddots & \vdots & \vdots \\ g_1(x_n^k) & \dots & g_n(x_n^k) & \frac{(-1)^{n+1}}{W(x_n^k)} \end{pmatrix} \times \begin{pmatrix} c_1^k \\ c_2^k \\ \vdots \\ c_n^k \\ \rho^k \end{pmatrix} = \begin{pmatrix} f(x_0^k) \\ f(x_1^k) \\ \vdots \\ f(x_n^k) \end{pmatrix} \quad (14)$$

Assume $g_{n+1}(x_i) = (-1)^{i+1}/W(x_i), i = 0, \dots, n$. If eqn. 14 is to be invertible, then from property 3.1.1 and 3.1.2, $\{g_1, \dots, g_n, g_{n+1}\}$ has to satisfy the Haar condition. That is, g_{n+1} is not in the n -space spanned by $\{g_1, \dots, g_n\}$. If $W(x)$ is chosen to satisfy the condition above, then the above system of equations can be solved to find ρ^k and the k th approximation polynomial as $P^k = \sum_{i=1}^n c_i^k g_i(x)$

- (c) From the polynomial P^k to find new extreme points $\{x_i^{k+1}\}$

(d) Check if $\{x_i^{k+1}\}$ is exactly equal to $\{x_i^k\}$, or equivalently, if ρ^{k+1} is equal to ρ^k . If not, go back to step b

(e) P^k is the best approximation polynomial to $f(x)$, and ρ^k is the ultimate equal ripple deviation with respect to the weighting function $W(x)$ on A . Denote this best approximation polynomial and deviation as $P^*(x)$ and ρ^* , respectively.

3.2 Polynomial interpolation [15, 16]

The interpolation technique is one of the simplest ways to determine the original function in many applications. The spline curve is commonly used as the interpolation basis in curve-fitting applications. In many applications the problem is how to find a well defined interpolation basis that is suitable for the application. A method, based on the linear functional theory of linear algebra, is introduced to find the interpolation basis required for the application in question.

3.2.1 Linear functional theory [15]

In the linear transformation domain, there exists a special case that transforms a n -dimensional vector space V^n into a field F (i.e., one-dimensional vector space). If some specifications are given to the n -dimensional space basis of V^n , for example, let the basis be $\{1, x, \dots, x^{n-1}\}$, this transformation that transforms from V^n to F then acts like a linear function. A formal definition of this theory is presented in the following:

Definition

Let V^n be a n -dimensional vector space over the field F ; l is a linear transformation from V^n to F (denote $l: V^n \rightarrow F$) and l satisfies

$$l(c\alpha + \beta) = cl(\alpha) + l(\beta) \quad \forall c \in F \text{ and } \alpha, \beta \in V^n \quad (15)$$

Thus, l is defined as the linear functional on V^n .

Dual space

Let $L(V^n, F) \triangleq V^*$ represent the set of all the linear functionals from V^n to F and $\{l_1, l_2, \dots, l_n\}$ be a set of linear functionals that satisfy $l_i(g_j) = \delta_{ij}$. Here $\{g_1, g_2, \dots, g_n\}$ is the basis of V^n and $\delta_{ij} = \delta(i - j)$. It can be shown that $\{l_1, l_2, \dots, l_n\}$ is linearly independent and forms a basis of V^* , that is, $\dim V^* = \dim V^n$. V^* is defined as the dual space of V^n or linear functional space.

Theorem 3.2.1

Let V^n be an n -dimensional vector space over $F, B = \{g_1, g_2, \dots, g_n\}$ be an ordered basis for V^n , and $B^* = \{l_1, l_2, \dots, l_n\}$ be an ordered basis for V^* . Then B^* is the unique basis, such that

$$l_i(g_j) = \delta_{ij} \quad i, j = 1, 2, \dots, n \quad (16)$$

and B^* is called the dual basis of B . For every g in V^n and l in V^* , $g = a_1 g_1 + a_2 g_2 + \dots + a_n g_n$ and $l = b_1 l_1 + b_2 l_2 + \dots + b_n l_n$. It can be easily shown that

$$\begin{aligned} a_i &= l_i(g) \\ b_i &= l(g) \quad i = 1, \dots, n \end{aligned} \quad (17)$$

Theorem 3.2.1 is an important property that will provide the theoretical foundation for finding the interpolation basis for some special applications.

3.2.2 Odd-polynomials

In many approximation applications, it is desirable to formulate the problems with a set of polynomial functions or a linear combination of some special functions, e.g., sine or cosine functions. For the present situation, if the complexity of analysis is considered, it is obvious that polynomial representation is the most suitable.

The desired filter spectral response has a symmetrical property. The ordinary technique, such as Lagrange interpolation, may not be able to take this characteristic into consideration, and will suffer from unnecessary computational complexity. The requirement for the half-

band filter necessitates the even terms of its impulse response to be zero. This characteristic indicates that the spectral response consists of only odd power polynomials. There is no need to consider the polynomial of even power.

If this is included in the formulation, $f(x) = a_0 + a_1x + a_2x^2 + \dots + a_{2n}x^{2n}$. Since $f(x)$ has odd-symmetry property, i.e., $f(x)$ is an odd-function, then $f(x) + f(-x) = 0$ and

$$2(a_0 + a_2x^2 + a_4x^4 + \dots + a_{2n}x^{2n}) = 0$$

It is clear that $a_{2i} = 0, i = 0, \dots, n$, and

$$f(x) = a_1x + a_3x^3 + \dots + a_{2n-1}x^{2n-1} \quad (18)$$

By this, the function that satisfies eqn. 18 is defined as the odd-polynomial function. In the same way, the function that contains only even power terms, is called even-polynomial function.

3.2.3 Interpolation basis

The classical Lagrange interpolation formula is a general interpolation technique for arbitrarily spaced sample points. The Lagrange interpolation may be unsuitable in this case. A new interpolation formula that will satisfy the requirements must be designed.

Question modelling

Let $B = \{x, x^3, \dots, x^{2n-1}\}$ be an ordered basis on V^n , and $B' = \{g_1, \dots, g_n\}$ be another ordered basis on V^n with

$$g_i = \sum_{j=1}^n b_{ij}x^{2j-1} \quad (19)$$

Let $\{f(x_1), f(x_2), \dots, f(x_n)\}$ be a set of n -distinct samples. Then, the question for finding an interpolation basis can be stated as follows:

A basis set $B' = \{g_1, \dots, g_n\}$, is desired that satisfies

$$f(x) = \sum_{i=1}^n f(x_i)g_i(x) \quad (20)$$

It is known that the dual space basis $B^* = \{l_1, \dots, l_n\}$ and

$$l_i(f) = f(x_i) \quad (21)$$

Before solving this question, it is necessary to ensure that B^* is indeed a basis of the dual space V^* . This can be proved if $\{l_1, \dots, l_n\}$ are linear independent.

Proof: Suppose $\{l_i, i = 1, \dots, n\}$ are linearly dependent and $\{a_1, \dots, a_n\}$ are not all zero and $a_i \in F, i = 1, \dots, n$. Such that

$$l = a_1l_1 + a_2l_2 + \dots + a_nl_n = 0$$

Then from theorem 3.2.1

$$l(f) = \sum_{i=1}^n a_i f(x_i) = 0 \quad \forall x_i \in F \quad i = 1, \dots, n$$

It is clear that $a_i = 0$, for $i = 1, \dots, n$ and this contradicts the assumption. So $\{l_1, l_2, \dots, l_n\}$ is linearly independent and forms a basis of V^* .

Solution of odd-polynomial basis

B^* is a basis of the dual space, so from theorem 3.2.1 and eqn. 21

$$g_j(x_i) = \delta_{ij} \quad i, j = 1, 2, \dots, n \quad (22)$$

Consider $j = 1$ and eqn. 19, eqn. 22 means that

$$b_{11}x_i + b_{12}x_i^3 + \dots + b_{1n}x_i^{2n-1} = \begin{cases} 1 & i = 1 \\ 0 & \text{otherwise} \end{cases}$$

in matrix form

$$A \cdot B = C \quad (23)$$

where

$$A = \begin{pmatrix} x_1 & x_1^3 & \dots & x_1^{2n-1} \\ x_2 & x_2^3 & \dots & x_2^{2n-1} \\ \vdots & \vdots & \ddots & \vdots \\ x_n & x_n^3 & \dots & x_n^{2n-1} \end{pmatrix}$$

and

$$B = (b_{11} \quad b_{12} \quad \dots \quad b_{1n})^t$$

$$C = (1 \quad 0 \quad \dots \quad 0)^t$$

It can be seen that A is similar to the Vandermonde matrix and will be nonsingular, that is

$$\det A = \prod_{i=1}^n x_i \prod_{1 \leq i < j \leq n} (x_j^2 - x_i^2) \quad (24)$$

and will be nonzero for distinct $\{x_i\}$. So, eqn. 23 is invertible and

$$b_{1k} = \frac{\det A_k}{\det A} \quad (25)$$

where A_k is obtained by replacing the k th column of A and C . From eqns. 19 and 25

$$g_1(x) = \sum_{k=1}^n \det A_k \cdot x^{2k-1} / \det A$$

$$= \det \begin{vmatrix} x & x^3 & \dots & x^{2n-1} \\ x_2 & x_2^3 & \dots & x_2^{2n-1} \\ \vdots & \vdots & \ddots & \vdots \\ x_n & x_n^3 & \dots & x_n^{2n-1} \end{vmatrix} / \det A$$

$$= \frac{\det A|_{x_1=x}}{\det A}$$

Similarly, for arbitrary j , say $j = k$, g_k can be described as

$$g_k(x) = \frac{\det A|_{x_k=x}}{\det A}$$

$$= \frac{x \prod_{j=1, j \neq k}^n (x^2 - x_j^2)}{x_k \prod_{j=1, j \neq k}^n (x_k^2 - x_j^2)} \quad k = 1, 2, \dots, n \quad (26)$$

This set of functions, $\{g_1, \dots, g_n\}$, is the interpolation basis required for the odd-polynomial function. The standard odd-polynomial interpolation formula can be expressed as

$$f(x) = \sum_{i=1}^n f(x_i) \frac{x \prod_{j=1, j \neq i}^n (x^2 - x_j^2)}{x_i \prod_{j=1, j \neq i}^n (x_i^2 - x_j^2)} \quad (27)$$

Uniqueness theorem

Let $B = \{g_i\}$ form an interpolation basis of V^n . If the two sample sets, $\{f(x_i)\}$ and $\{f'(x_i)\}$, simultaneously satisfy

$$f(x) = \sum_{i=1}^n f(x_i)g_i(x) = \sum_{i=1}^n f'(x_i)g_i(x)$$

then

$$f(x_i) = f'(x_i) \quad i = 1, 2, \dots, n$$

Proof: Let $f(x) = \sum_{i=1}^n f(x_i)g_i(x)$, if there exists another set of samples $\{f'(x_i)\}$ which also satisfies $f(x) = \sum_{i=1}^n f'(x_i)g_i(x)$. Then

$$0 = \sum_{i=1}^n (f(x_i) - f'(x_i)) \cdot g_i(x)$$

because $\{g_i\}$ is a basis of V^n , it is clear that $f(x_i) = f'(x_i)$, for $i = 1, \dots, n$. Then it can be concluded that for any $f(x)$ there exists a unique expression with $\{g_i\}$.

3.2.4 Barycentric form

The interpolation method is very useful for reconstructing the original function from arbitrary spaced sample points. Unfortunately, it will take a lot of computations when the order of the interpolation is high. Hamming [16] introduced a special barycentric form for Lagrange interpolation which has about $n/2$ fewer computations than that of the standard Lagrange interpolation formula. This form will make the interpolation technique more useful.

In the odd-polynomial case, if the desired function $f(x)$ is known for a given x , that is for a distinct sample set $\{x_i\}$, $i = 1, \dots, n$, then $f(x_i) = x_i$.

From eqn. 27

$$\sum_{i=1}^n \left(\frac{\prod_{j=1, j \neq i}^n (x^2 - x_j^2)}{\prod_{j=1, j \neq i}^n (x_i^2 - x_j^2)} \right) = 1 \quad (28)$$

because of the uniqueness property. If eqn. 27 is divided by eqn. 28

$$f(x) = x \frac{\sum_{k=1}^n [\alpha_k / (x^2 - x_k^2)] \cdot f(x_k)}{\sum_{k=1}^n [\alpha_k / (x^2 - x_k^2)] \cdot x_k} \quad (29)$$

where

$$\alpha_k = \frac{1}{x_k} \prod_{j=1, j \neq k}^n \frac{1}{(x_j^2 - x_k^2)} \quad (30)$$

Eqn. 29 is the barycentric form for the odd-polynomial interpolation formula, and it has about $(n+1)/3$ fewer computations than that of the standard form.

3.3 Design method and procedure

There are many methods [11, 12, 19, 20] that can be used to design the half-band filter. Suppose $\{f(n)\}$, for $n = -(N-1), \dots, (N-1)$ is the impulse response of the half-band filter and

$$F(Z) = \sum_{n=-(N-1)}^{N-1} f(n)Z^{-n} \quad (31)$$

because $F(Z) + F(-Z) = 1$, by applying eqn. 31

$$2 \left[f(0) + \sum_{\substack{n=-N+1 \\ n=\text{even}}}^{N-1} f(n)Z^{-n} \right] = 1 \quad (32)$$

Clearly, eqn. 32 means that

$$f(n) = \begin{cases} 0.5 & n = 0 \\ 0.0 & n = \text{even} \end{cases} \quad (33)$$

Since $f(0)$ is always a value of 0.5, it is more convenient to set $f(0)$ to zero during the design process. The perfect reconstruction condition then becomes

$$F'(Z) + F'(-Z) = 0 \quad \text{and} \quad f'(2m) = 0 \quad m \in Z \quad (34)$$

The design for the half-band filter is to first design a linear phase filter $f'(n)$ according to eqn. 34, and then let $f(n) = f'(n)$ and $f(0) = 0.5$.

3.3.1 Filter response represented in terms of odd-polynomials

The condition for designing the perfect reconstructed half-band filter in eqn. 34 has been presented. Since the phase and $f(0)$ are temporarily neglected, only the amplitude response is of concern

$$F(\omega) = 2 \sum_{n=0}^{N/2-1} f'(2n+1) \cos(2n+1)\omega \quad (35)$$

where N is even for the QMF design. It is clear that $f'(N-1) = f'(-N+1) = 0$ if N is odd. So, only the case when N is even is of concern.

Property 3.3

For a cosine function $\cos(n\omega)$, if $n = 2k+1$, there is a set $\{a_i\}$, $i = 0, \dots, k$, such that

$$\cos(2k+1)\omega = \sum_{i=0}^k a_i (\cos \omega)^{2i+1} \quad (36)$$

If $n = 2k$, then there is another set $\{b_i\}$, $i = 0, \dots, k$, such that

$$\cos(2k)\omega = \sum_{i=0}^k b_i (\cos \omega)^{2i} \quad (37)$$

eqns. 36 and 37 can be combined as

$$\cos(n\omega) = \sum_{i=0}^n t_{ni} (\cos \omega)^i \quad (38)$$

where

$$t_{ni} = \begin{cases} a_k & i = 2k+1 \\ 0 & \text{otherwise} \end{cases} \quad n = 2m+1$$

and

$$t_{ni} = \begin{cases} b_k & i = 2k \\ 0 & \text{otherwise} \end{cases} \quad n = 2m$$

$\{t_{ni}\}$ $i = 1, 2, \dots, n$, is the coefficient of the n th Chebyshev polynomial [9, 16] $T_n(x)$ and eqn. 38 becomes

$$\cos(n\omega) = T_n(\cos \omega) \quad (39)$$

From property 3.3.1, eqn. 35 could be rewritten as

$$F(\omega) = 2 \sum_{n=0}^{N/2-1} f''(n) (\cos \omega)^{2n+1} \quad (40)$$

where

$$f''(n) = \sum_{k=n}^{N/2} f'(2k+1) \cdot t_{(2k+1)(2n+1)} \quad (41)$$

and $t_{(2k+1)(2n+1)}$ is the $(2n+1)$ th coefficient of the $(2k+1)$ th Chebyshev polynomial. The half-band filter spectral response can be considered as an odd-polynomial function, and will be useful when applying the Remes exchange algorithm for finding the extreme points.

3.2.2 Iterative procedure

Since $F(\omega) = -F'(\pi - \omega)$, only the response during the interval $[0, \pi/2]$ has to be considered. Let A denote the region $[0, \omega_p]$ and ω_p the passband edge. It has been shown that $\{\cos(2i+1)\omega\}$ and $\{(\cos \omega)^{2i+1}, i = 0, 1, \dots, N/2 - 1\}$, satisfy the Haar condition in $[0, \pi/2]$. From the alternation theorem, there is a set of extreme points $\{\omega_i\}$, $\omega_i \in A$, $i = 0, \dots, N/2$, and $\omega_0 = 0$ and $\omega_{N/2} = \omega_p$. Since the half-band filter is an extra-ripple filter, the Chebyshev approximation for the half-band filter design then becomes

$$\min_{\omega \in [0, \omega_p]} \left(\max W(\omega) |F(\omega) - D(\omega)| \right)$$

Apply the Remes iterative procedure as stated earlier, the following $N/2 + 1$ equations are obtained

$$F(\omega_i) - \frac{(-1)^i \rho}{W(\omega_i)} = D(\omega_i) \quad i = 0, \dots, N/2 \quad (42)$$

From eqn. 14

$$\begin{pmatrix} \cos \omega_0 & \cos 3\omega_0 & \cdots & \cos(N-1)\omega_0 & \frac{-1}{W(\omega_0)} \\ \cos \omega_1 & \cos 3\omega_1 & \cdots & \cos(N-1)\omega_1 & \frac{1}{W(\omega_1)} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \cos \omega_{N/2} & \cos 3\omega_{N/2} & \cdots & \cos(N-1)\omega_{N/2} & \frac{(-1)^{N/2}}{W(\omega_{N/2})} \end{pmatrix} \times \begin{pmatrix} f'(1) \\ f'(3) \\ \vdots \\ f'(N-1) \\ \rho \end{pmatrix} = \begin{pmatrix} D(\omega_0) \\ D(\omega_1) \\ \vdots \\ D(\omega_{N/2}) \end{pmatrix} \quad (43)$$

and

$$\begin{pmatrix} x_0 & x_0^3 & \cdots & x_0^{(N-1)} & \frac{-1}{W(\omega_0)} \\ x_1 & x_1^3 & \cdots & x_1^{(N-1)} & \frac{1}{W(\omega_1)} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ x_{N/2} & x_{N/2}^3 & \cdots & x_{N/2}^{(N-1)} & \frac{(-1)^{N/2}}{W(\omega_{N/2})} \end{pmatrix} \times \begin{pmatrix} f''(0) \\ f''(1) \\ \vdots \\ f''(N/2-1) \\ \rho \end{pmatrix} = \begin{pmatrix} D(\omega_0) \\ D(\omega_1) \\ \vdots \\ D(\omega_{N/2}) \end{pmatrix} \quad (44)$$

with respect to the two different bases. Here $x_i = \cos \omega_i$. The weighting function $W(\omega)$ is generally set to be unity because of band separability, and $D(\omega)$ is set to be 0.5 so that the ultimate $|F(\omega)| \leq \rho^*$, for $\omega \in [\omega_s, \pi]$.

If eqn. 43 or 44 is solved directly and eqn. 35 or 40 is used to find the new set of extreme points, it will be time-consuming. In the following, a simple method, using the odd-polynomial interpolation basis discussed earlier, is presented and eqn. 27 can be restarted as

$$F(\omega) = \sum_{i=0}^{N/2-1} F'(\omega_i) g_{i+1}(x_i) \quad (45)$$

where $\{g_k\}$, $k = 1, \dots, N/2$, are the interpolation basis and satisfies eqn. 26. Thus eqn. 14 can be rewritten as

$$\begin{pmatrix} & & & \frac{-1}{W(\omega_0)} \\ & & & \vdots \\ & & & \frac{-(-1)^{N/2-1}}{W(\omega_{N/2-1})} \\ I & & & \frac{-(-1)^{N/2}}{W(\omega_{N/2})} \\ g_1(x_{N/2}) & \cdots & g_{N/1}(x_{N/2}) & \end{pmatrix} \times \begin{pmatrix} F'(\omega_0) \\ F'(\omega_1) \\ \vdots \\ F'(\omega_{N/2-1}) \\ \rho \end{pmatrix} = \begin{pmatrix} D(\omega_0) \\ D(\omega_1) \\ \vdots \\ D(\omega_{N/2}) \end{pmatrix} \quad (46)$$

From eqn. 26

$$g_{k+1}(x_{N/2}) = \frac{x_{N/2}}{x_k} \prod_{\substack{j=0 \\ j \neq k}}^{N/2-1} \left(\frac{x_{N/2}^2 - x_j^2}{x_k^2 - x_j^2} \right) = \frac{a_k}{a_{N/2}} \quad k = 0, \dots, N/2 - 1 \quad (47)$$

and

$$\det \begin{pmatrix} & & & a_1 \\ & & & a_2 \\ & & & \vdots \\ & & & a_n \\ b_1 & b_2 & \cdots & b_n & c \end{pmatrix} = c - (a_1 b_1 + a_2 b_2 + \cdots + a_n b_n) \quad (48)$$

where

$$a_k = \frac{1}{x_k} \prod_{\substack{i=0 \\ i \neq k}}^{N/2} \frac{1}{x_i^2 - x_k^2} \quad k = 0, \dots, N/2 \quad (49)$$

From eqns. 47 and 48, ρ and $F'(\omega_i)$ in eqn. 46 can easily be shown to be

$$\rho = \frac{\sum_{k=0}^{N/2} a_k D(\omega_k)}{\sum_{k=0}^{N/2} (-1)^k a_k / W(\omega_k)} \quad (50)$$

and

$$F'(\omega_i) = D(\omega_i) + (-1)^i \rho / W(\omega_i) \quad i = 0, \dots, N/2 - 1 \quad (51)$$

Since eqn. 46 is invertible and from eqn. 45

$$F'(\omega_{N/2}) = F'(\omega_0) g_1(x_{N/2}) + \cdots + F'(\omega_{N/2-1}) g_{N/2}(x_{N/2})$$

then, the condition

$$F'(\omega_{N/2}) = D(\omega_{N/2}) + (-1)^{N/2} \rho / W(\omega_{N/2})$$

will be automatically satisfied. So, eqn. 51 is exactly the same as eqn. 42. Only the deviation ρ need be solved and $F'(\omega_i)$ will be obtained immediately.

It is more efficient to use the barycentric form in eqn. 29 than that of eqn. 44 and write $F'(\omega)$ as

$$F'(\omega) = x \frac{\sum_{k=1}^{N/2-1} [\alpha_k / (x^2 - x_k^2)] F'(\omega_k)}{\sum_{k=1}^{N/2-1} [\alpha_k / (x^2 - x_k^2)] x_k} \quad (52)$$

where

$$\alpha_k = a_k(x^2 - x_{N/2}^2) \quad k = 0, \dots, N/2 - 1 \quad (53)$$

$$x = \cos \omega \quad \omega \in \left[0, \frac{\pi}{2}\right]$$

After several iterations, this procedure will converge, and the polynomial that best approximates the desired filter response will be obtained. The impulse response could then be calculated by taking the $2N$ points IDFT. This approximation is based on the odd-polynomial technique, so the best approximation $P^*(x)$ is guaranteed to be odd-symmetry with respect to $\omega = \pi/2$, and the inverse transformation would be

$$f'(n) = \begin{cases} \frac{1}{N} \left(2 \sum_{k=1}^{N/2-1} F'(k) \cos \frac{2\pi kn}{2N} + F'(0) \right) & n = \text{odd} \\ 0 & n = \text{even} \end{cases} \quad (54)$$

3.3.3 Normalisation

Since the definition of the half-band filter is $F(Z) = H_0(Z)H_0(Z^{-1})$, the spectrum of $F(Z)$ will be $|H_0(\omega)|^2$. It can be easily seen that $F(\omega)$ must be non-negative for all ω . An adjustment is required to make $F(\omega)$ greater than zero for ω in $[\pi/2, \pi]$. This can be achieved by adding the stop band deviation to the impulse response $f(0)$. If $f(0) = 0.5$ is to be maintained, a scaling to all the impulse responses can be first applied and then a separate bias is added to $f(0)$. The scaling factor and the bias factor were chosen to be $(1 + \rho)^{-1}$ and ρ respectively. The filter before and after adjustment is shown in Figs. 1 and 2, respectively.

4 Numerical results and conclusions

The impulse responses of the half-band filters designed by the proposed odd-polynomial based algorithm and the Parks-McClellan algorithm were compared. The devi-

ation ρ and the filter coefficients of three filters are listed in Tables 1, 2 and 3 with passband cutoff frequencies at

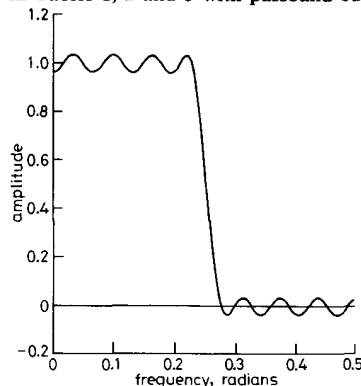


Fig. 1 Half-band filter before impulse response adjustment

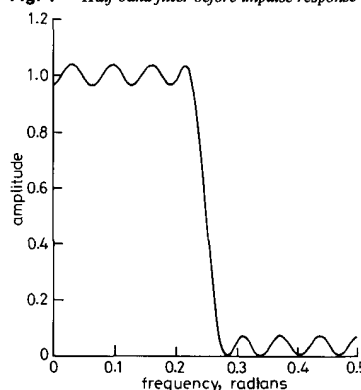


Fig. 2 Half-band filter after impulse response adjustment

Table 1: Coefficients of half-band filter: $\omega_p = 0.2$; 63 taps

$f(n)$	Odd-polynomial method	P-M method	P-M and fix DELF
0	0.5000000	0.49997663	0.50000000
1	0.3168833114447524	0.31688333	0.31688380
2	0.0000000	$0.22997244 \times 10^{-4}$	$0.75975381 \times 10^{-14}$
3	-0.1018922883168539	-0.10189227	-0.10189371
4	0.0000000	$-0.21782890 \times 10^{-4}$	$-0.73153564 \times 10^{-14}$
5	$5.8866083966100550 \times 10^{-2}$	$0.56866088 \times 10^{-1}$	$0.56868309 \times 10^{-1}$
6	0.0000000	$0.19965335 \times 10^{-4}$	$0.68201617 \times 10^{-14}$
7	$-3.6403494753731619 \times 10^{-2}$	$-0.36403461 \times 10^{-1}$	$-0.36406313 \times 10^{-1}$
8	0.0000000	$-0.17579249 \times 10^{-4}$	$-0.62368541 \times 10^{-14}$
9	$2.4419601528327786 \times 10^{-2}$	$0.24419587 \times 10^{-1}$	$0.24422762 \times 10^{-1}$
10	0.0000000	$0.14952601 \times 10^{-4}$	$0.55508948 \times 10^{-14}$
11	$-1.6553439221261703 \times 10^{-2}$	$-0.16553388 \times 10^{-1}$	$-0.16556697 \times 10^{-1}$
12	0.0000000	$-0.12161705 \times 10^{-4}$	$-0.48492956 \times 10^{-14}$
13	$1.1121337371768617 \times 10^{-2}$	$0.11121307 \times 10^{-1}$	$0.11124457 \times 10^{-1}$
14	0.0000000	$0.95174840 \times 10^{-5}$	$0.41351402 \times 10^{-14}$
15	$-7.3105661915249653 \times 10^{-2}$	$-0.73104931 \times 10^{-2}$	$-0.73133593 \times 10^{-2}$
16	0.0000000	$-0.70589279 \times 10^{-5}$	$-0.35042516 \times 10^{-14}$
17	$4.6535299490812190 \times 10^{-3}$	$0.46534820 \times 10^{-2}$	$0.46558763 \times 10^{-2}$
18	0.0000000	$0.50157642 \times 10^{-5}$	$0.29498009 \times 10^{-14}$
19	$-2.8403192665608168 \times 10^{-3}$	$-0.28402353 \times 10^{-2}$	$-0.28421697 \times 10^{-2}$
20	0.0000000	$-0.33100121 \times 10^{-5}$	$-0.25090159 \times 10^{-14}$
21	$1.6441697294677111 \times 10^{-3}$	$0.16441241 \times 10^{-2}$	$0.16455292 \times 10^{-2}$
22	$0.20815364 \times 10^{-5}$	0.20815364×10^5	$0.21314520 \times 10^{-14}$
23	$-8.9026472022364720 \times 10^{-4}$	$-0.89019217 \times 10^{-3}$	$-0.89118005 \times 10^{-3}$
24	0.0000000	$-0.11609320 \times 10^{-5}$	$-0.19186592 \times 10^{-14}$
25	$4.4216740908140399 \times 10^{-4}$	$0.44213563 \times 10^{-3}$	$0.44273468 \times 10^{-3}$
26	0.0000000	$0.61545419 \times 10^{-6}$	$0.17472796 \times 10^{-14}$
27	$-1.9523598087233321 \times 10^{-4}$	$-0.19518295 \times 10^{-3}$	$-0.19554197 \times 10^{-3}$
28	0.0000000	$-0.24318148 \times 10^{-6}$	$-0.16955132 \times 10^{-14}$
29	$7.2256687490477233 \times 10^{-5}$	$0.72250165 \times 10^{-4}$	$0.72396340 \times 10^{-4}$
30	0.0000000	$0.10312851 \times 10^{-6}$	$0.16111451 \times 10^{-14}$
31	$-1.9781406364411270 \times 10^{-6}$	$-0.19752272 \times 10^{-4}$	$-0.19827561 \times 10^{-4}$
Deviation	$5.8635426463977397 \times 10^{-6}$	$5.8710522358252987 \times 10^{-6}$	0.0000059

Table 2: Coefficients of half-band filter: $\omega_p = 0.2$; 31 taps

$f(n)$	Odd-polynomial method	P-M method	P-M and fix DELF
0	0.0000000	0.49999567	0.50000000
1	0.3156775975760746	0.31567876	0.31567777
2	0.0000000	$0.47849695 \times 10^{-5}$	$0.59987857 \times 10^{-16}$
3	$-9.8417417341069860 \times 10^{-2}$	$-0.98416379 \times 10^{-1}$	$-0.98417448 \times 10^{-1}$
4	0.0000000	$-0.41355751 \times 10^{-5}$	$-0.25964893 \times 10^{-16}$
5	$5.1521577152782043 \times 10^{-2}$	$0.51521795 \times 10^{-1}$	$0.51521801 \times 10^{-1}$
6	0.0000000	$0.51690321 \times 10^{-5}$	$0.52825128 \times 10^{-16}$
7	$-2.9870007475233616 \times 10^{-2}$	$-0.29778452 \times 10^{-1}$	$-0.29780828 \times 10^{-1}$
8	0.0000000	$-0.20391174 \times 10^{-5}$	$-0.20592846 \times 10^{-16}$
9	$1.7205509741091583 \times 10^{-2}$	$0.17204932 \times 10^{-1}$	$0.17205631 \times 10^{-1}$
10	0.0000000	$0.28744002 \times 10^{-5}$	$0.47453081 \times 10^{-16}$
11	$-9.4239049264127355 \times 10^{-3}$	$-0.94219386 \times 10^{-2}$	$-0.94244290 \times 10^{-2}$
12	0.0000000	$-0.73721880 \times 10^{-6}$	$-0.13430117 \times 10^{-16}$
13	$4.6470034769669449 \times 10^{-3}$	$0.46465782 \times 10^{-2}$	$0.46479407 \times 10^{-2}$
14	0.0000000	$0.19344144 \times 10^{-5}$	$0.50586775 \times 10^{-16}$
15	$-2.1047008418378838 \times 10^{-3}$	$-0.21042776 \times 10^{-2}$	$-0.21068203 \times 10^{-2}$
Deviation	$1.3486852752779199 \times 10^{-3}$	$1.3493378417375243 \times 10^{-3}$	0.0013528

Table 3: Coefficients of half-band filter: $\omega_p = 0.2$; 15 taps

$f(n)$	Odd-polynomial method	P-M method	P-M and fix DELF
0	0.0000000	0.50001572	0.50000000
1	0.3139355169302561	0.31392599	0.31391401
2	0.0000000	$0.13742684 \times 10^{-4}$	$-0.46259293 \times 10^{-17}$
3	$-9.3435674787155573 \times 10^{-2}$	$-0.93403660 \times 10^{-1}$	$-0.93415821 \times 10^{-1}$
4	0.0000000	$0.15855960 \times 10^{-6}$	0.00000000
5	$4.4115594365228733 \times 10^{-2}$	$0.44110376 \times 10^{-1}$	$0.44101428 \times 10^{-1}$
6	0.0000000	$0.10384331 \times 10^{-4}$	$-0.64763010 \times 10^{-17}$
7	$-2.6477022404434195 \times 10^{-2}$	$-0.26487629 \times 10^{-1}$	$-0.26488660 \times 10^{-1}$
Deviation	$2.3723171792209918 \times 10^{-2}$	$2.3774135388997422 \times 10^{-2}$	0.0237781

0.18, 0.20 and 0.23, respectively. The deviation of the filter designed by the odd-polynomial based method is smaller than that of the P-M method. A better numerical accuracy is thus obtained. The third column of the Table is obtained by setting the DELF of the original Parks-McClellan program to 0.001. This made the dense grid of points replacing the frequency axis symmetric about $\pi/2$. With this setting, the even terms of the filter coefficients by the P-M method were calculated to be zero to within the computer accuracy.

The barycentric interpolation formula of the odd-polynomial based method has 3/4 fewer multiplications than that of the P-M method. The odd-polynomial based method searches only one-half of the extreme points, so, in every iteration, there is a saving of 3/8 on the multiplications. That is, in every iteration, the odd-polynomial based method will be 3/8 times faster than the P-M method.

Another difference is the number of iterations required by the two methods. The P-M method uses only the $N + 1$ extreme points for iteration. It needs an internal exchange among the $N + 2$ extreme points. The odd-

polynomial based method uses all the extreme points for iteration because of its symmetry property. It will thus require fewer iterations than the P-M method. For example, in Table 1, it takes only seven iterations for the odd-polynomial based method and 12 iterations for the P-M method. A comparison between the computing time is shown in Figs. 3, 4 and 5, where an average gain of

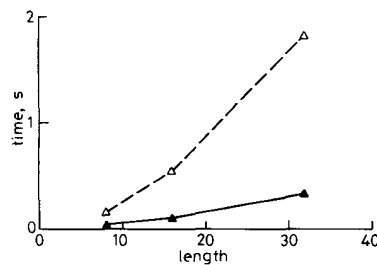


Fig. 3 Computing time against filter length
 $\omega_p = 0.18$
—▲— odd-polynomial method
--△-- P-M method

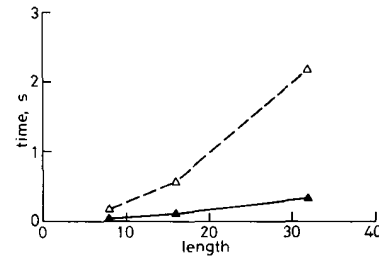


Fig. 4 Computing time against filter length
 $\omega_p = 0.20$
—▲— odd-polynomial method
--△-- P-M method

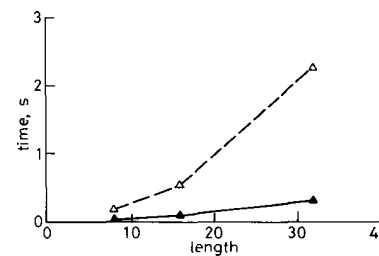


Fig. 5 Computing time against filter length
 $\omega_p = 0.23$
—▲— odd-polynomial method
--△-- P-M method

four for the proposed method over the P-M method is obtained.

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