decision threshold ambiguity width (DTAW) and the phase margin (PM). The DTAW is defined as the minimum data input required to obtain an error rate of less than  $1 \times 10^{-9}$ . The PM is the maximum extent of the phase difference between the clock and the data under 0.8 V peak-to-peak data input and under the above mentioned error rate.3 The measured DTAW and PM against data rate are plotted in Fig. 3.

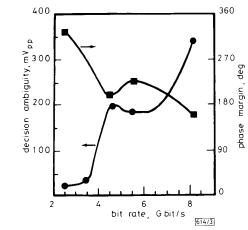
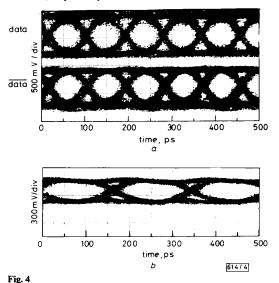


Fig. 3 Decision threshold ambiguity width and phase margin for decision circuit against data rate

Demultiplexing can be achieved with the use of masterslave D-type flip-flops in a series-to-parallel convertor or in a parallel processing arrangement, where two flip-flops demultiplex alternating bits from the input sequence. Demultiplexing operation using the latter architecture was verified by demultiplexing a single channel up to 11.2 Gbit/s. This was accomplished by using data and half the clock frequency as inputs to the IC. A bit error rate of less than  $1 \times 10^{-12}$  was measured, along with a phase margin (phase shift of the f/2 clock relative to the 11.2 Gbit/s data) of 50.5 deg. The 11.2 Gbit/s input eye pattern and the resulting demultiplexed eye are shown in Figs. 4a and b, respectively.



a 11.2 Gbit/s NRZ test signal at demultiplexer input b 5.6 Gbit/s demultiplexed eye

Conclusion: In conclusion, we have designed and implemented a submicron silicon bipolar D-type flip-flop integrated circuit which can be employed either as a decision circuit or a demultiplexer, operating at data rates substantially higher than previously reported silicon bipolar flip-flop integrated circuits, and with performance comparable to or better than, state-ofthe-art GaAs IC technology. The performance of our integrated circuit indicates that silicon bipolar technology, with its relatively mature process and proven reliability as well as its low-cost potential, could play an important role in highspeed lightwave communication systems of up to 10 Gbit/s, or even higher data rates.

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## THROUGHPUT EFFICIENCY OF SOME ARQ STRATEGIES UNDER MARKOV ERROR **MODELS**

Indexing terms: Telecommunications, Data transmission. Errors

In the letter we compute and compare the throughput efficiencies of three continuous ARQ strategies under Markov error models. The three strategies are the classic go-back-N, the Sastry modification, and the Moeneclaey and Bruneel scheme. Results show that Sastry's modification is outperformed by either the classic go-back-N scheme or the Moeneclaey and Bruneel scheme.

Introduction: Automatic repeat request (ARO) strategies are common techniques used to handle transmission errors. Stopand-wait, go-back-N, and selective-repeat are the three most popular ARQ strategies in use and have been extensively studied. The independent error assumption is often adopted to simplify the analysis of ARQ strategies. This assumption, however, is not realistic for channels where transmission errors occur in bursts. Markov error models are usually used to describe the behaviour of such channels.

In this letter we compute and compare the throughput efficiencies of three ARQ strategies under Markov error models. The three ARQ strategies are classic go-back-N, Sastry's modification, and the one proposed by Moeneclaey and Bruneel. By a Markov error model is meant that the probability that a particular data block arrives at the receiver with error depends on whether the immediately preceding transmission is a success or a failure. Generally, a Markov channel can be described as follows. Let  $X_i$  denote the outcome of the transmission in the ith slot so that  $X_i = 0$  means the transmission is a success and  $X_i = 1$  means the transmission is a failure. Then the error model can be described by  $P[X_{i+1} = 0 \mid X_i = 0] = p$  and  $P[X_{i+1} = 1 \mid X_i = 1] = q$ . For convenience, let

$$T = \begin{bmatrix} p & 1-p \\ 1-q & q \end{bmatrix} \tag{1}$$

denote the state transition matrix. It is clear that p + q = 1 corresponds to the independent error model. Moreover, it can be shown<sup>4</sup> that the k-step transition matrix is given by

$$T^{k} = \frac{1}{2 - (p+q)} \begin{bmatrix} 1 - q & 1 - p \\ 1 - q & 1 - p \end{bmatrix} + \frac{(p+q-1)^{k}}{2 - (p+q)} \begin{bmatrix} 1 - p & p-1 \\ q-1 & 1-q \end{bmatrix}$$
(2)

For convenience, we use  $t_{ij}^k$  (i, j = 0, 1) to denote the (i, j)th entry of  $T^k$ .

Throughput efficiency: The throughput efficiency of an ARQ scheme is defined as the inverse of the expected number of transmissions to deliver successfully a data block. For simplicity, we assume that round-trip delay s is fixed and is equal to an integral number of blocks. Notice that round-trip delay is defined as the interval between the end of a transmission and the receipt of its response. In the following, we consider separately the three ARQ strategies of interest under the assumption of a noiseless feedback channel:

Fig. 1 Typical transmission sequence of Sastry scheme for s=3

(a) Classic go-back-N ARQ scheme: Since the classic go-back-N ARQ strategy is well known its operation is not explained here. The throughput efficiency of this scheme under Markov error models has been derived<sup>2</sup> and is given by

$$\eta_c = \frac{(1-q)[1-(p+q-1)^{s+1}]}{(s+1)[2-(p+q)](1-p)+(1-q)[1-(p+q-1)^{s+1}]}$$
(3)

(b) Sastry's modification: The operation of Sastry's modification<sup>2</sup> is similar to the classic go-back-N ARQ scheme. However, whenever a repeat request is received, the same data block is retransmitted continuously until a positive acknowledgment (ACK) is received. It is noted that the transmission sequence of this scheme can be divided into cycles so that each cycle starts with a 1, i.e. an unsuccessful transmission. Fig. 1 shows a typical transmission sequence. By regenerative theorems,<sup>4</sup> one can compute the throughput efficiency of this scheme by finding the ratio of the expected number of successfully delivered data blocks to the expected number of transmissions in a cycle. Let L and M denote, respectively, the expected cycle length and the expected number of data blocks successfully delivered in a cycle. Then it can be shown that

$$L = 2(s+1) + \frac{1}{1-q} \left[ (1-p)t_{10}^s + qt_{11}^s \right] + \frac{1}{1-p} \left[ pt_{00}^s + (1-q)t_{01}^s \right]$$
(4)

After some manipulations, we obtain

$$L = 2(s+1) + \frac{(1-p) + (1-q)(p+q-1)^{s+1}}{(1-q)[2-(p+q)]} + \frac{(1-q) + (1-p)(p+q-1)^{s+1}}{(1-p)[2-(p+q)]}$$
(5)

Similarly, it can also be shown that

$$M = 1 + \frac{(1-q) + (1-p)(p+q-1)^{p+1}}{(1-p)[2-(p+q)]}$$
 (6)

Therefore, the throughput efficiency of the Sastry modification is given by

$$\eta_s = \frac{(1-p)(1-q)\{[2-(p+q)] + (p+q-1)^{s+1}\} \\ + (1-q)^2}{2(1-p)(1-q)\{(s+1)[2-(p+q)] + (p+q-1)^{s+1}\} \\ + (1-p)^2 + (1-q)^2}$$

(c) Moeneclaey and Bruneel's scheme: For the Moeneclaey and Bruneel scheme, <sup>3</sup> each data block is transmitted continuously to the receiver until an ACK is received. The throughput efficiency for this scheme can also be computed by using regenerative theorems. In reality, we have

$$\eta_m = \frac{(1-q)[2-(p+q)]}{(s+1)[2-(p+q)](1-q)+(1-p)[1-(p+q-1)^{s+1}]}$$

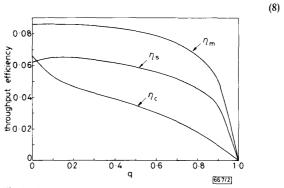


Fig. 2 Throughput efficiency as function of q for p = 0.1 and s = 10

Discussion: When p+q=1, we have  $\eta_c=(1-q)/(sq+1)$ ,  $\eta_s=(1-q)/[2sq(1-q)+1]$ , and  $\eta_m=(1-q)/[s(1-q)+1]$ , in accordance with the results obtained previously under the independent error model. It was shown that Sastry's modification is always outperformed by either the classic go-back-N scheme or the Moeneclaey and Bruneel scheme under the independent error model. The same phenomenon under Markov error models can be observed in Figs. 2-4. Therefore, it seems also unrealistic to use Sastry's modification for Markov channels. Besides, the Moeneclaey and Bruneel

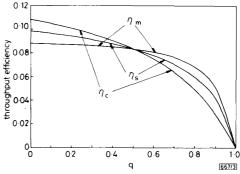


Fig. 3 Throughput efficiency as function of q for p = 0.5 and s = 10

scheme is suitable for noisy channels (i.e. small p and large q), especially when round-trip delay is large.

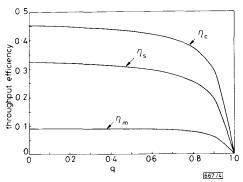


Fig. 4 Throughput efficiency as function of q for p=0.9 and s=10

For a Markov channel with state transition matrix T, the equivalent block error probability is equal to (1-p)/(2-p-q) for the corresponding independent error model. It was proved that the classic go-back-N ARQ strategy is more efficient for a Markov system than the corresponding system

under the independent error model if and only if  $p+q \ge 1$ . One can easily verify that the conclusion is opposite for the Moeneclaey and Bruneel scheme, i.e. the Moeneclaey and Bruneel scheme is more efficient for a Markov system than the corresponding system under the independent error model if and only if  $p+q \le 1$ .

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## OPTIMISED MULTIPLE QUANTUM WELL PHASE MODULATOR

Indexing terms: Optoelectronics, Optical modulation, Phase modulation. Integrated optics

The quantum confined Stark effect (QCSE) in GaAs/AlGaAs multiple quantum well structures has been shown by several authors to cause large changes of refractive index<sup>1,2</sup> and absorption coefficient<sup>3,4</sup> near the band edge which are potentially of use for phase modulation applications. Hence in a phase modulator there is necessarily a trade-off between the magnitude of the phase modulation, highest near the exciton features, and that of the amplitude modulation. In this letter we report a pin phase modulator that has been optimised for maximum phase change relative to the contrast (ratio of transmission in the 'on' state to that in the 'off' state), within certain constraints: the maximum tolerable contrast and the maximum zero bias absorption loss were chosen to be 2 and 50%, respectively. The parameters optimised were the operating wavelength, the number of quantum wells in the active layer and the position of that layer, the operating bias, the device length and the background doping in the intrinsic region. The optimised device was measured to have a phase modulation figure of merit of 27°/V mm at the optimum bias of 6 V, giving a phase change in a 385  $\mu$ m-long device of 63 at this bias. The corresponding contrast was approximately 1.4. Photocurrent measurements show that the zero bias absorption loss is also within the 50% limit.

The optimisation process uses data from a single quantum well, normal incidence photodiode to obtain the electric field dependence of an isolated quantum well so as to be able to model the behaviour of waveguide devices.<sup>5</sup> Photocurrent data from a normal-incidence pin photodiode incorporating a single 84 Å GaAs quantum well in the Al<sub>0.3</sub>Ga<sub>0.7</sub>As intrinsic region was used to obtain the zero bias absorption coefficient, the change in the absorption coefficient and, via the Kramers-Krönig relation, the change in refractive index for various biases over the wavelength range 820-880 nm.6 From the spectra obtained the operating wavelength was optimised for maximum phase change relative to contrast in the absence of waveguiding, bearing in mind that the final device had to exhibit a contrast of less than 2 and a zero bias of less than 50%. The optimum wavelength was found to be 863 nm, some 17 nm below that of the heavy hole excition feature. This set of data was then used to optimise a waveguide device with a 1 μm-thick core of the same refractive index as 84 Å/60 Å GaAs/Al<sub>0.3</sub>Ga<sub>0.7</sub>As multiple quantum well material. The core could be of this material, of the alloy Al<sub>0·12</sub>Ga<sub>0·88</sub>As, which has the same index, or some combination of the two. The cladding index was then chosen so that the device would be

single-mode, and so the cladding material was of Al<sub>0.15</sub>Ga<sub>0.85</sub>As. The modelling was done for TE propagation, since only in this case are the selection rules the same as in the transverse case.<sup>7</sup> The details of the model are described in Reference 6. Various device parameters were then varied to optimise for maximum ratio of phase change to contrast within the defined operating constraints.

The optimised structure is shown in Fig. 1. The position and thickness of the quantum well layer, as well as being chosen to reduce amplitude modulation, are chosen so that the device is tolerant of deviations of the background doping, since although the optimum doping was found to be about  $5 \times 10^{16} \, \mathrm{cm}^{-3}$  it is difficult in practice to control this value sufficiently accurately. The GaAs capping layer is present to facilitate the formation of good ohmic contacts. The structure was grown by atmospheric pressure MOVPE and the contacts were evaporated onto the *p*-type capping layer (Cr/Au) and the *n*-type substrate (Sn/Au). The contacts were alloyed at a temperature of 440°C for 2 min, and then the contacted material was fabricated into slab waveguides approximately 300  $\mu$ m in length.

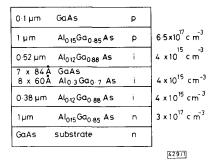


Fig. 1 Structure of optimised phase modulator

Light from a Coherent 599-01 dye laser using Styryl 9M dye, pumped by a Coherent Innova I-90 argon ion laser, was end-fire coupled into the devices, and contrast as a function of DC bias was measured using conventional lock-in techniques. The results for a 385 µm long device are shown in Fig. 2, along with the contrast predicted from the model. The phase change as a function of bias in the same device was measured by placing the sample in one arm of a Mach-Zehnder interferometer and observing the shift of the fringe pattern. The shifted and unshifted fringe patterns, imaged on a TV camera, were stored in a frame store and each was averaged. By taking the complex Fourier transform of the shifted and unshifted patterns one can determine the phase shift from the rotation in the complex plane of the vector corresponding to the spatial