

# 國立交通大學

電子工程學系 電子研究所碩士班

## 碩士論文

應用於混模電路測試之利用混合基底與內插法  
的快速傅立葉轉換



**FFT Based on Mixed-Radix and  
Interpolation for Mixed-Signal Testing**

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中華民國 九十四年 六月

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## 摘要

在混模電路測試上，快速傅立葉轉換被廣泛地用來將時域訊號轉成頻域訊號以取得待測電路的傳輸參數，例如頻率響應、相位響應及訊號雜音比等等。但是，轉換過程需要大量的運算並且受制於取樣點數必須是 2 的冪次方。

在此論文中，吾人提出一個利用混合基底 2/4、3 及 5 的快速傅立葉轉換演算法以增加運算速度並且增加可應用的取樣點數為 2、3 及 5 的冪次組合，此外再透過對取樣數據作內插的方法將取樣點數擴增為任意數。吾人先討論單一基底的快速傅立葉轉換，接著再將演算法擴充到混合基底。一個用來將取樣數據重新排列的演算法被找出來，這使得混合基底的快速傅立葉轉換可行。然後，吾人再將內插法的觀念與混合基底的快速傅立葉轉換結合，使混合基底的快速傅立葉轉換能操作在任意取樣點數。透過模擬，吾人從理論與實驗上對不同演算法作運算速度及取樣點數的比較。最後，吾人提出兩個實驗結果作為範例，一個是對多音訊號取樣，另一個是方波訊號。實驗結果顯示：在內插法造成訊號雜音比下降的可接受範圍內，演算法確實有速度上的優勢。

# FFT Based on Mixed-Radix and Interpolation for Mixed-Signal Testing

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In mixed-signal testing, FFT (fast Fourier transformation) is used widely to transform the time domain signal into the frequency domain signal to obtain the transmission parameters such as frequency response, phase response, signal to noise ratio (S/N), etc, for the circuit under test (CUT). However, the process is involved with large computation efforts and the limitation on the number of sampling points which need to be of the power of 2.

In this thesis, we propose a FFT algorithm based on the mixed radices of 2/4, 3 and 5 to increase the computation speed and to increase the applicable number of sampling points to the power of 2, 3 and 5, and furthermore to any number of sampling points by using an interpolation technique to interpolate the sampled data. We first discuss the single-radix FFT and then expand the algorithm to mixed-radix FFT. A re-ordering algorithm for the sampled data is found which makes the mixed-radix FFT possible. We then bring out the concept of

interpolation to be in conjunction with the mixed-radix FFT to make the mixed-radix FFT algorithm be able to be applied to any number of sampling points. We compare the computation speed and the numbers of sampling points for different algorithms both theoretically and experimentally via simulation. Finally, experimental results on two examples, one is a multi-tone signal and the other is a square wave signal, are presented. They show that the algorithm can really demonstrate the speed advantage with acceptable degradation on S/N introduced by the interpolation process.



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# Chapter 1 Introduction

## 1.1 Motivation

Mixed-signal circuits are widely used in the present electronic system, especially, in SoC. To test mixed-signal circuits [1], Digital Signal Processing approach is usually adopted, which is shown in Figure 1.1. An appropriately selected test patterns are applied the mixed signal circuit under test (CUT) and its response is transformed into the frequency domain by using the discrete Fourier transform (DFT) technique. The transformed spectrum is analyzed to deduce the transmission parameters such as frequency response, S/N, etc, for the CUT. It can then analyze the obtained parameters by comparator to mark the CUT as good or faulty.

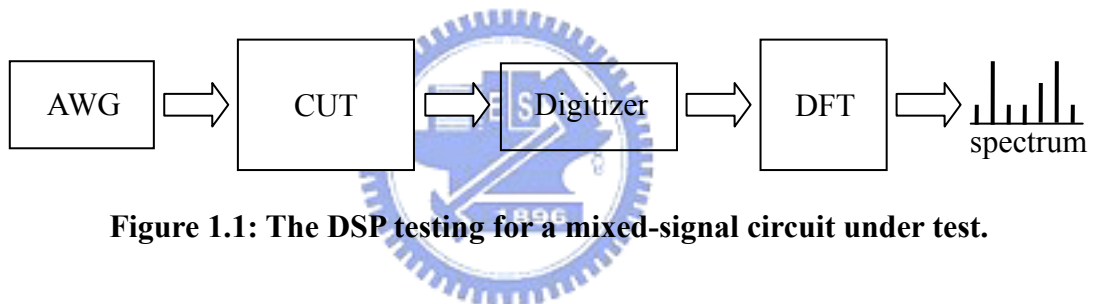


Figure 1.1: The DSP testing for a mixed-signal circuit under test.

In the above, the DFT is a computation-intensive process which is proportional to  $N^2$ , where  $N$  is the number of sample points of the output response of the CUT. In 1965, a much faster algorithm was developed by Cooley and Tukey [2], which is called the fast Fourier transform (FFT) algorithm [3]-[7]. The computation complexity for the algorithm (radix-2 FFT) is proportional to  $M \log_2 N$ .

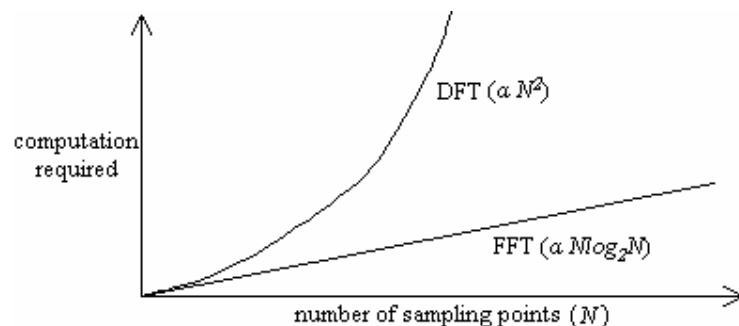


Figure 1.2: A comparison of computation complexity between DFT and FFT.

Figure 1.2 plots the comparison of computation complexity of the DFT and FFT algorithms.

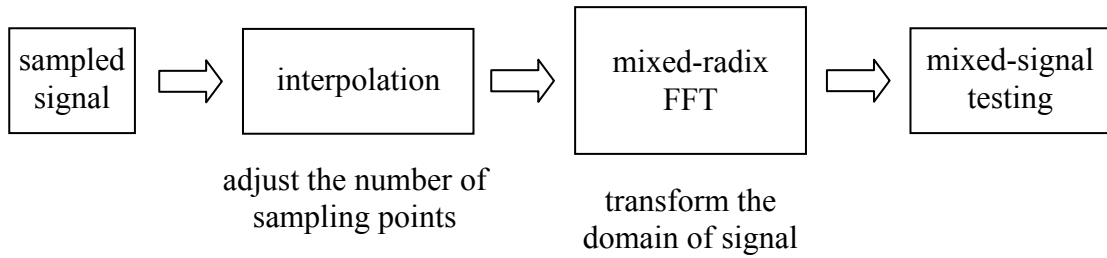
However, radix-2 FFT has an important limitation, that is, the number of sampling points must be of the power of 2, i.e.,  $2^k$ , where  $k$  is a natural number. This means that the numbers of sampling points to be able to be applied by using the FFT algorithm are very limited. For example as shown in Table 1.1, for a sampling system which can hold 10000 data points, there are only 13 numbers of sampling points can be used the FFT.

$k$	1	2	3	4	5	6	7	8	9	10	11	12	13
$2^k$	2	4	8	16	32	64	128	256	512	1024	2048	4096	8196
distance	-	2	4	8	16	32	64	128	256	512	1024	2048	4096

**Table 1.1: The numbers of sampling points which can be applied by the radix-2 FFT under the value 10000.**

In this thesis work, it is to investigate to apply the FFT algorithm to, in addition to the number of sampling points that is of radix 2, other numbers which could be of other radices such as 3, 4, and 5 and a mixed of the above numbers. This will increase the range of application of FFT. In addition, to apply the FFT in a mixed-radix fashion may have a faster computation speed since its number of computation stages is fewer than that of the radix-2 FFT when their numbers of sampling points are similar.

Furthermore, although mixed-radix FFT can increase the numbers of sampling, which can be applied by the FFT algorithm, significantly, the numbers are still too few. Hence we will investigate a technique to use interpolation to adjust the sampling points to be able to be applied the mixed-radix FFT, as shown in Figure 1.3, to further increase the FFT efficiency.



**Figure 1.3: A method that combines interpolation with mixed-radix FFT.**

However, by using the interpolation to the number of sampling points to match up the mixed-radix FFT, we will introduce errors to the final computation results, which are expressed in worse S/N's. When the distribution of the numbers of sampling points is even enough, the difference between the original number and the interpolated number will be small. However, this will usually lead to a more complicated mixed-radix FFT. So we will study the trade-off between the decreased S/N and the increased computation complexity.

## 1.2 Review of Previous Works

To speed up the FFT operation, there were several algorithms proposed. An FFT using radix-3, 6, and 12 algorithms was used in an ordinary complex plane and the numbers of additions and multiplications were shown to have significantly reduced [8]. Radix-2/8 FFT algorithm was also demonstrated to be able to save real multiplications and have much lower arithmetic complexity than the radix-2/4 FFT algorithm [9]. The form of radix- $p/p^2$  algorithm is called a split-radix algorithm, which is better than radix- $p$  algorithm on length- $p^m$  DFT's. It was shown that whenever a radix- $p^2$  outperforms a radix- $p$  algorithm, then a radix- $p/p^2$  algorithm will outperform both of them [10][11]. As to the limitation of numbers of sampling points, there were two algorithms that can do the FFT for any number of sampling points [12][13]. The proposed algorithms can flexibly compute the discrete Fourier transforms of length  $q \times 2^m$  where  $q$  is an odd integer. Comparisons with previously reported algorithms show that substantial savings on arithmetic operations can be made. Furthermore, a wider range of

choices on different sequence lengths is naturally provided.

In this thesis, we propose a mixed radixes FFT based on mixed radixes 2/4,3,5 incorporating with an interpolation technique to speed up the FFT while still maintaining a good S/N.

### **1.3 Outline of This Thesis**

This thesis is organized as follows. Chapter 2 presents the research on single-radix FFT. Chapter 3 presents the research on mixed-radix FFT. Chapter 4 presents the radix-2/4/3/5 FFT and interpolation. In chapter 5, conclusions are given.





## Chapter 2 Single-Radix FFT

Before deriving the mixed-radix FFT algorithm, we first study the single-radix FFT algorithm. We first study the decimation-in-time algorithm for the general case and then apply it to radix-3 FFT, radix-4 FFT and radix-5 FFT. As for the radix-2 FFT [6][7], since it is well known in the general literature, we will not mention it here in detail.

### 2.1 Decimation-In-Time Algorithm

Decimation is the process of breaking down a group of data into several of its constituents. Decimation-in-time involves breaking down a signal which is a set of discrete time data into several smaller set of discrete time data which represent the same signal.

Let us consider  $N (= L^S)$ -point DFT where  $N$  is the number of total sampled data points,  $L$  is a decimating factor and  $S$  is the corresponding exponent. We split the  $N$ -point data sequence  $x(n)$  into  $L$  data sequences,  $f_1(n), f_2(n), \dots, f_L(n)$ , with  $N/L$  points for each sequence. That is,

$$\begin{aligned} f_1(n) &= x(Ln) \\ f_2(n) &= x(Ln+1) \\ &\vdots \\ f_L(n) &= x(Ln+L-1), \quad n = 0, 1, \dots, \frac{N}{L}-1 \end{aligned}$$

$f_1(n), f_2(n), \dots, f_L(n)$  are the decimated sub-sequences of  $x(n)$  by the factor  $L$ .

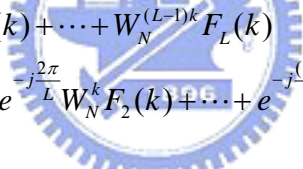
Now the  $N$ -point DFT can be expressed in terms of the DFT's of the decimated sub-sequences as follows:

$$\begin{aligned}
X(k) &= \sum_{n=0}^{N-1} x(n)W_N^{kn}, \quad k = 0, 1, \dots, N-1 \\
&= \sum_{m=0}^{N/L-1} x(Lm)W_N^{Lmk} + \sum_{m=0}^{N/L-1} x(Lm+1)W_N^{(Lm+1)k} + \dots + \sum_{m=0}^{N/L-1} x(Lm+L-1)W_N^{(Lm+L-1)k}
\end{aligned}$$

Where  $W_N^L = W_{N/L}$ . With this substitution, the equation can be expressed as

$$\begin{aligned}
X(k) &= \sum_{m=0}^{N/L-1} f_1(m)W_{N/L}^{km} + W_N^k \sum_{m=0}^{N/L-1} f_2(m)W_{N/L}^{km} + \dots + W_N^{(L-1)k} \sum_{m=0}^{N/L-1} f_L(m)W_{N/L}^{km} \\
&= F_1(k) + W_N^k F_2(k) + \dots + W_N^{(L-1)k} F_L(k), \quad k = 0, 1, \dots, N-1
\end{aligned}$$

Where  $F_1(k), F_2(k), \dots, F_L(k)$  are the  $N/L$ -point DFT's of the sequences  $f_1(m), f_2(m), \dots, f_L(m)$ , respectively. Since  $F_1(k), F_2(k), \dots, F_L(k)$  are periodic with period  $N/L$ , we have  $F_1(k+N/L) = F_1(k), F_2(k+N/L) = F_2(k), \dots, F_L(k+N/L) = F_L(k)$ . Hence the equation may be expressed as

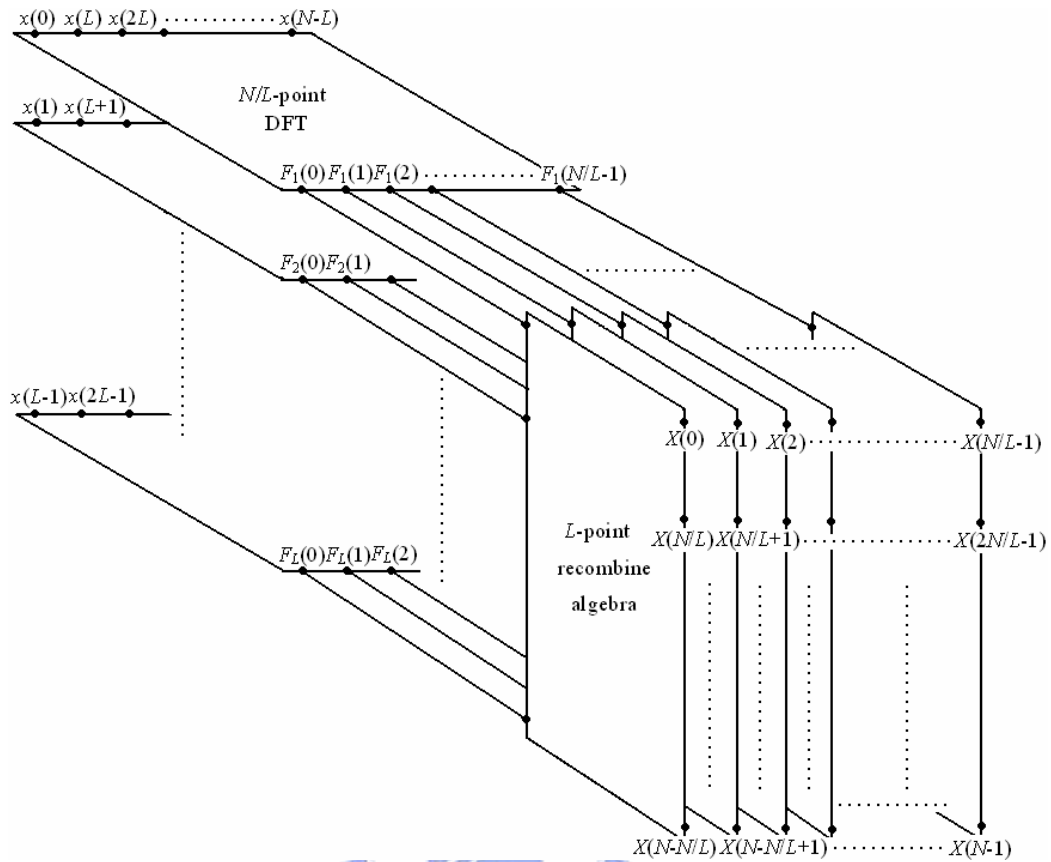


$$\begin{aligned}
X(k) &= F_1(k) + W_N^k F_2(k) + \dots + W_N^{(L-1)k} F_L(k) \\
X(k + N/L) &= F_1(k) + e^{-j\frac{2\pi}{L}} W_N^k F_2(k) + \dots + e^{-j\frac{(L-1)2\pi}{L}} W_N^{(L-1)k} F_L(k) \\
&\vdots \\
X(k + (L-1)N/L) &= F_1(k) + e^{-j\frac{(L-1)2\pi}{L}} W_N^k F_2(k) + \dots + e^{-j\frac{(L-1)^2 \cdot 2\pi}{L}} W_N^{(L-1)k} F_L(k) \\
&k = 0, 1, \dots, N/L - 1
\end{aligned}$$

We can observe that the direct computation of  $F_1(k)$  requires approx  $(N/L)^2$  complex multiplications. The same applies to other  $N/L$ -point DFT's. Furthermore, there are  $(L-1)N$  additional complex multiplications required to compute other parts. Hence the entire computation of  $X(k)$  requires  $L(N/L)^2 + (L-1)N = N^2/L + LN$  complex multiplications. The first step results in a reduction of the number of complex multiplications from approximate  $N^2$  to  $N^2/L + LN$ , which is about a factor of  $L$  for a large  $N$ .

This process of decimating the signal can easily be visualized. The first breakup into

several  $N/L$ -point DFT's can be shown as Figure 2.1.



**Figure 2.1: The first step in the decimation-in-time algorithm.**

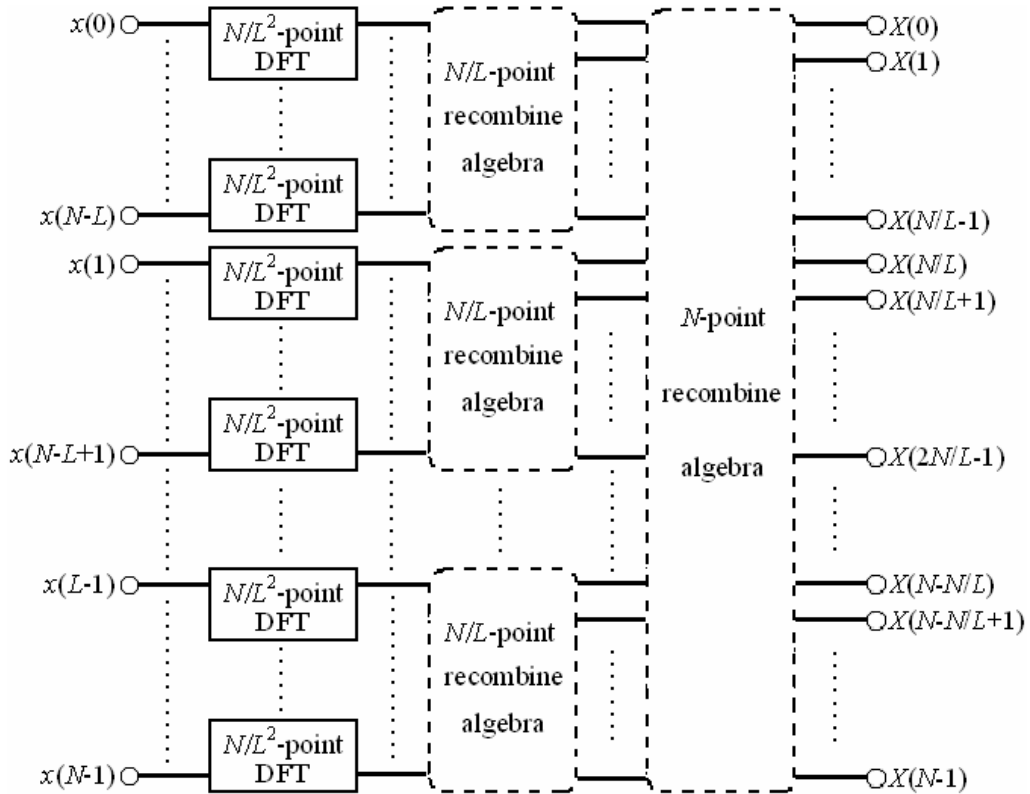
We can continue to split the  $N/L$ -point data sequence  $f_i(n)$  into  $L$  data sequences with  $N/L^2$  points for each sequence  $g_{i1}(n), g_{i2}(n), \dots, g_{iL}(n)$ , that is,

$$\begin{aligned}
 g_{i1}(n) &= f_i(Ln) \\
 g_{i2}(n) &= f_i(Ln+1) \\
 &\vdots \\
 g_{iL}(n) &= f_i(Ln+L-1), \quad i=1,2,\dots,L; \quad n=0,1,\dots,\frac{N}{L^2}-1
 \end{aligned}$$

By computing  $N/L^2$ -point DFT's, we obtain the  $N/L$ -point DFT's  $F_1(k), F_2(k), \dots, F_L(k)$  from the relationship:

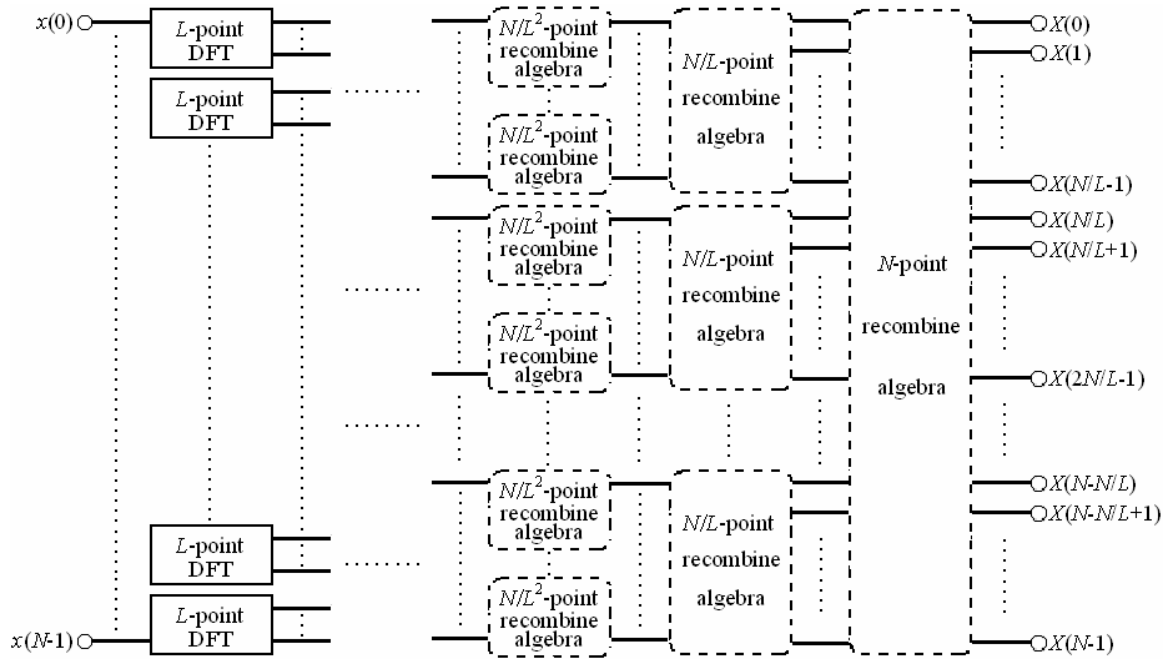
$$\begin{aligned}
F_i(k) &= G_{i1}(k) + W_{N/L}^k G_{i2}(k) + \dots + W_{N/L}^{(L-1)k} G_{iL}(k) \\
F_i(k + N/L^2) &= G_{i1}(k) + e^{-j\frac{2\pi}{L}} W_{N/L}^k G_{i2}(k) + \dots + e^{-j\frac{(L-1) \cdot 2\pi}{L}} W_{N/L}^{(L-1)k} G_{iL}(k) \\
&\vdots \\
F_i(k + (L-1)N/L^2) &= G_{i1}(k) + e^{-j\frac{(L-1) \cdot 2\pi}{L}} W_{N/L}^k G_{i2}(k) + \dots + e^{-j\frac{(L-1)^2 \cdot 2\pi}{L}} W_{N/L}^{(L-1)k} G_{iL}(k) \\
&\quad i = 1, 2, \dots, L; \quad k = 0, 1, \dots, N/L^2 - 1
\end{aligned}$$

The decimation process is taken another stage by breaking down the  $N/L$ -point DFT's into  $N/L^2$ -point DFT's as shown in Figure 2.2.



**Figure 2.2: The second step in the decimation-in-time algorithm.**

The decimation of the data sequence can be repeated again and again until a series of  $L$ -point DFT's as shown in Figure 2.3 is reached.



**Figure 2.3: The last step in the decimation-in-time algorithm.**

Since each stage takes approx  $L^2 \times (N/L) = LN$  complex multiplications and there are  $\log_L N (= S)$  stages, the total number of complex multiplications is  $LM \log_L N$ . This means that this decimation approach has reduced the number of complex multiplications from approx  $N^2$  to  $LM \log_L N$ . In section 2.2 we will derive the  $L$ -point FFT which can reduce a little more complex multiplications and then replace the  $L$ -point DFT with the  $L$ -point FFT.

Another important observation is on the order of the input sequence after it is decimated  $(S-1)$  times. For example, if we consider the case where  $N = 9$  and  $L = 3$ , we know that the decimation yields the sequence in the order of:  $\{x(0), x(3), x(6), x(1), x(4), x(7), x(2), x(5), x(8)\}$ . An algorithm which can re-order the input data sequence will be discussed in section 2.3.

## 2.2 $L$ -point FFT

We have derived the decimation-in-time algorithm which can reduce the amount of computation significantly for any single-radix FFT. But the decimated result is not the

simplest since the equation of the  $L$ -point DFT has conjugate pairs that can be simplified.

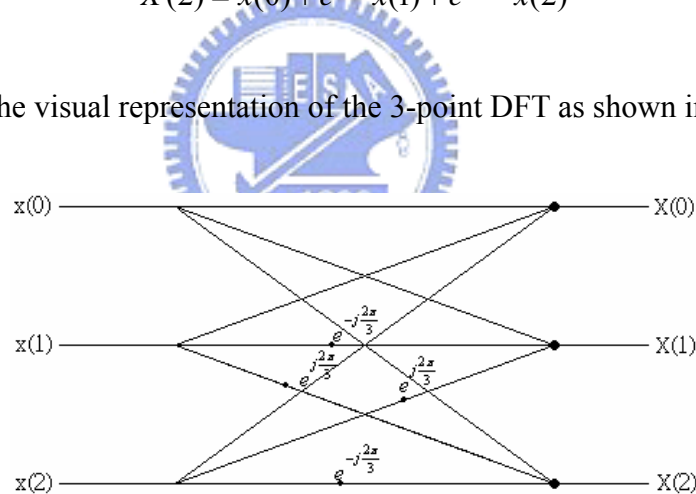
Here we will show the process of deriving the 3-point FFT from the equation of the 3-point DFT as an example and then show the 4-point FFT and 5-point FFT. Finally, we will have a table to compare the amount of computation of these three algorithms with that of the radix-2 FFT.

### 2.2.1 3-point FFT

Let us consider the case of the 3-point DFT. According to equations:

$$\begin{aligned} X(0) &= x(0) + x(1) + x(2) \\ X(1) &= x(0) + e^{-j\frac{2\pi}{3}} x(1) + e^{j\frac{2\pi}{3}} x(2) \\ X(2) &= x(0) + e^{j\frac{2\pi}{3}} x(1) + e^{-j\frac{2\pi}{3}} x(2) \end{aligned}$$

We can obtain the visual representation of the 3-point DFT as shown in Figure 2.4.



**Figure 2.4: The visual representation of the 3-point DFT.**

The 3-point DFT requires 4 complex multiplications and 6 complex additions. Thus if we consider the computation of the  $N_3 (= 3^S)$ -point DFT, the decimation approach will require  $(4N_3/3)\log_3 N_3$  complex multiplications and  $2N_3\log_3 N_3$  complex additions.

Let us pay attention to the equations of 3-point DFT, and then we can find that there are only two multiplications,  $\exp(j2\pi/3)$  and  $\exp(-j2\pi/3)$ . Since these two complex values are conjugate, they have equal real-part and inverse imaginary-part. Hence we can rewrite the

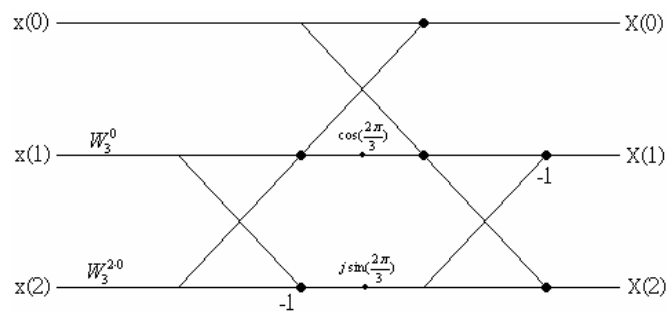
equations as follows:

$$X(0) = X(0) + W_3^0 X(1) + W_3^{2 \cdot 0} X(2)$$

$$X(1) = X(0) + \cos\left(\frac{2\pi}{3}\right) \left[ W_3^0 X(1) + W_3^{2 \cdot 0} X(2) \right] - j \sin\left(\frac{2\pi}{3}\right) \left[ W_3^0 X(1) - W_3^{2 \cdot 0} X(2) \right]$$

$$X(2) = X(0) + \cos\left(\frac{2\pi}{3}\right) \left[ W_3^0 X(1) + W_3^{2 \cdot 0} X(2) \right] + j \sin\left(\frac{2\pi}{3}\right) \left[ W_3^0 X(1) - W_3^{2 \cdot 0} X(2) \right]$$

According to these new equations, we can obtain the visual representation of the 3-point FFT as shown in Figure 2.5. We call this a “Butterfly Structure” of the radix-3 FFT.

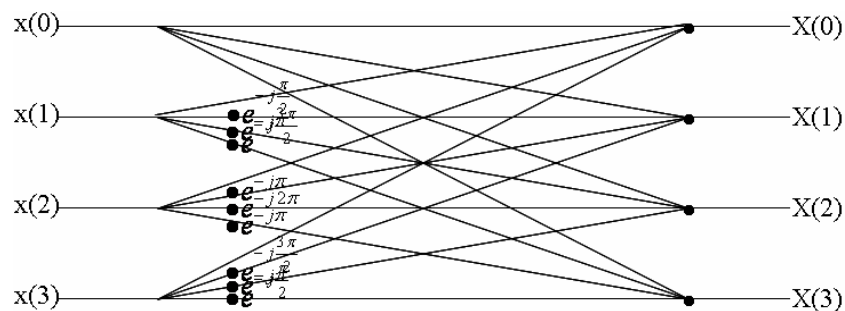


**Figure 2.5: The visual representation of the 3-point FFT (Butterfly structure).**

The 3-point FFT requires 3 complex (2 complex + 2 real) multiplications and 6 complex additions. Thus if we consider the computation of the  $N_3$ -point FFT, the decimation approach will require  $N_3 \log_3 N_3$  complex multiplications and  $2N_3 \log_3 N_3$  complex additions.

## 2.2.2 4-point FFT

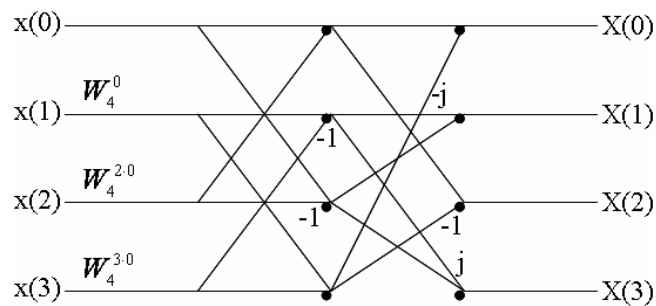
We directly consider the visual representation of the 4-point DFT as shown in Figure 2.6.



**Figure 2.6: The visual representation of the 4-point DFT.**

The 4-point DFT requires 9 complex multiplications and 12 complex additions. Thus if we consider the computation of the  $N_4 (= 4^S)$ -point DFT, the decimation approach will require  $(9N_4/4)\log_4 N_4$  complex multiplications and  $3N_4\log_4 N_4$  complex additions.

And then we consider the visual representation of the 4-point FFT as shown in Figure 2.7. We call this a “Butterfly Structure” of the radix-4 FFT.

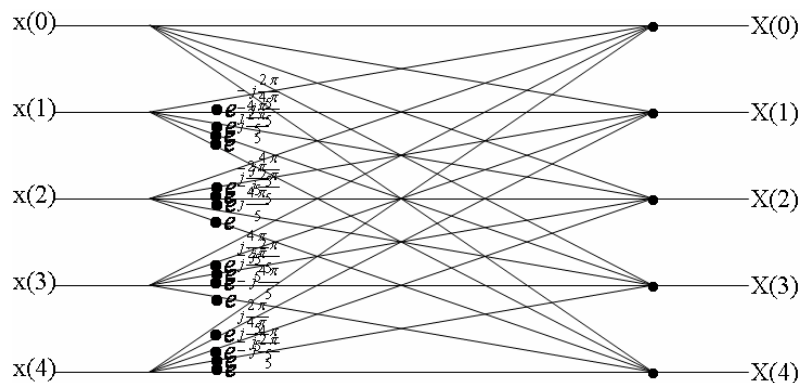


**Figure 2.7: The visual representation of the 4-point FFT (Butterfly Structure).**

The 4-point FFT requires 3 complex multiplications and 8 complex additions, thus if we consider the computation of the  $N_4$ -point FFT, the decimation approach will require  $(3N_4/4)\log_4 N_4$  complex multiplications and  $2N_4\log_4 N_4$  complex additions.

### 2.2.3 5-point FFT

We directly give the visual representation of the 5-point DFT as shown in Figure 2.8.



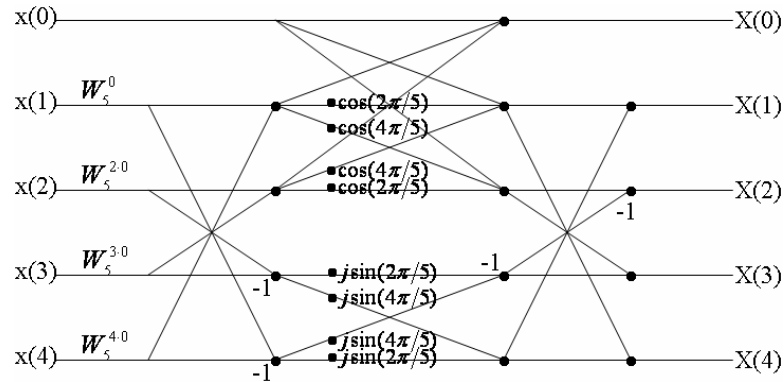
**Figure 2.8: The visual representation of the 5-point DFT.**

The 5-point DFT requires 16 complex multiplications and 20 complex additions, thus if



we consider the computation of the  $N_5 (= 5^S)$ -point DFT, the decimation approach will require  $(16N_5/5)\log_5 N_5$  complex multiplications and  $4N_5\log_5 N_5$  complex additions.

And then we give the “Butterfly Structure” of the 5-point FFT as shown in Figure 2.9.



**Figure 2.9: The visual representation of the 5-point FFT (Butterfly Structure).**

The 5-point FFT requires 8 complex (4 complex + 8 real) and 16 complex additions. Thus if we consider the computation of the  $N_5$ -point FFT, the decimation approach will require  $(8N_5/5)\log_5 N_5$  complex multiplications and  $(16N_5/5)\log_5 N_5$  complex additions.

Obviously, the number of complex multiplications is reduced again by replacing  $L$ -point DFT with  $L$ -point FFT. We have final comparison as shown in Table 2.1.

algorithm	radix-2 FFT	radix-3 FFT	radix-4 FFT	radix-5 FFT
complex multiplication	$(N/2)\log_2 N$	$M\log_3 N$	$(3N/4)\log_4 N$	$(8N/5)\log_5 N$
complex addition	$M\log_2 N$	$2M\log_3 N$	$2M\log_4 N$	$(16N/5)\log_5 N$

**Table 2.1: The comparison of computation complexity of different FFT algorithms.**

Although radix-2 FFT has the minimum number of complex multiplication  $(N/2)$  for each stage, however other single-radix FFT's have fewer number of computation stages. Thus other algorithms could have better performance in speed than radix-2 FFT in certain circumstances. This will be discussed further in section 2.4.

## 2.3 Re-Ordering Algorithm for Input Sequence

The process of decimating the signal in the time domain has made input samples need to be re-ordered. For a 9-point signal, the original order of the samples is 0, 1, 2, 3, 4, 5, 6, 7 and 8. But after decimating by radix-3 FFT the order becomes 0, 3, 6, 1, 4, 7, 2, 5 and 8. The order can be obtained by representing the number in the ternary form as follows in Figure 2.10. In the figure, once the numbers are represented in the ternary form, the digits of the representing ternary bits are reversed. The new numbers represented by the reversed digits are the new sequence which is to be applied to the decimated FFT.

<u>ORIGINAL INPUT</u>			<u>RE-ORDERED INPUT</u>	
<u>Decimal</u>	<u>Ternary</u>		<u>Ternary</u>	<u>Decimal</u>
0	00		00	0
1	01		10	3
2	02		20	6
3	10		01	1
4	11	➔	11	4
5	12	<i>bit-reversal</i>	21	7
6	20		02	2
7	21		12	5
8	22		22	8

Figure 2.10: The process of re-ordering the input samples for radix-3 FFT.

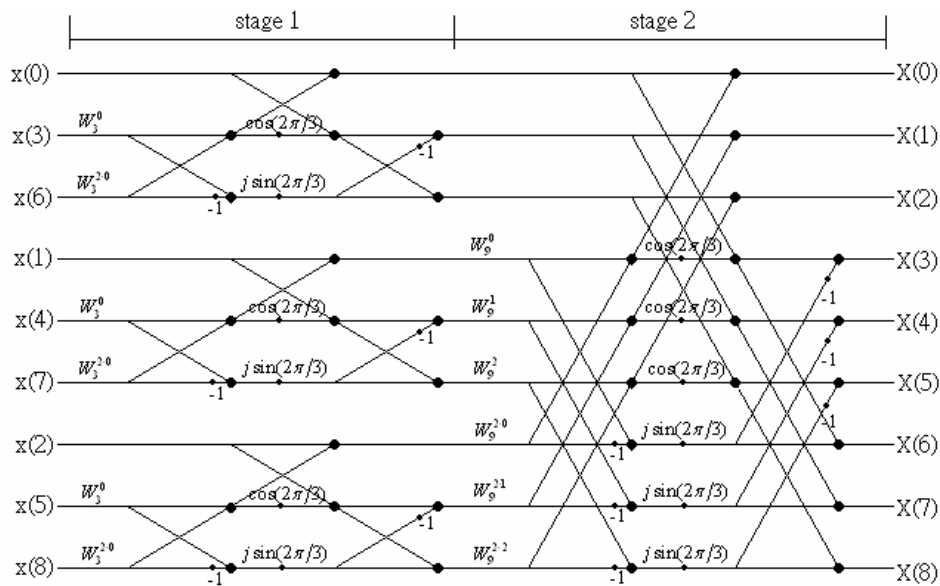


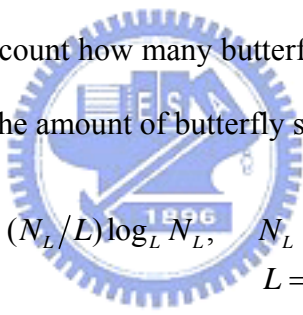
Figure 2.11: 9-point decimation-in-time FFT algorithm.

For the illustrative purpose, we depict the computation of 9-point FFT in Figure 2.11. The computation is performed in two stages, beginning with the computations of three 3-point FFT's, and finally one 9-point recombine algebra.

As for the radix-4 FFT or radix-5 FFT, we can also apply the same bit-reversal procedure to re-order the in put sequence. In the process, we have to remember to first represent the numbers in the quaternary and quinary form respectively.

## 2.4 Speed Ratio of Radix- $L$ FFT to Radix-2 FFT

In section 2.2, we have a table to compare the computation complexity of different single-radix FFT's. In this section, however we will use the idea of butterfly structure to obtain the same result. We will count how many butterfly structures are used for an algorithm. We have an equation to obtain the amount of butterfly structures for each algorithm, that is,



$$\text{radix-}L \text{ FFT: } (N_L/L)\log_L N_L, \quad \begin{array}{l} N_L \text{ is of the power of } L \\ L = 2, 3, 4, 5 \end{array}$$

Where  $N_L/L$  is the number of butterfly structures at each stage and  $\log_L N_L$  is the number of stages for an algorithm. The computation time of a butterfly structure ( $T_L$ ) is:

$$\text{radix-}L \text{ FFT: } T_L = \frac{\text{total run-time}}{\text{amount of butterfly structures}} = \frac{T_{run}}{(N_L/L)\log_L N_L}$$

Because the computation time of the process of re-ordering is very small and can be ignored, so we can directly use the total run-time of the radix- $L$  FFT to calculate  $T_L$ .

We use MATLAB to run the radix- $L$  FFT, record the total run-time and the amount of butterfly structures for each algorithm, and then we can obtain  $T_L$  for each algorithm as shown in Table 2.2.

$T_L$	$T_2$	$T_3$	$T_4$	$T_5$
us	1.76	2.59	3.47	4.69

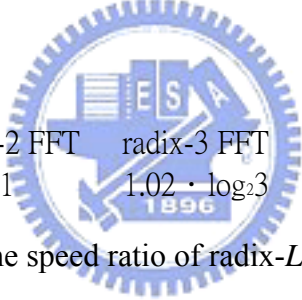
**Table 2.2: The computation time of a butterfly structure for each algorithm.**

After obtaining  $T_L$  for each algorithm, we can start to calculate the speed ratio. Let us assume the numbers of sampling points  $N_2$  is very close to  $N_3$ ,  $N_4$  and  $N_5$  in value, then according to the following equation,

$$\text{speed ratio } (L \text{ to } 2) = \frac{1/T_{run-L}}{1/T_{run-2}} = \frac{1/T_L(N_L/L)\log_L N_L}{1/T_2(N_2/2)\log_2 N_2} \approx \frac{T_2 \times L}{T_L \times 2} \cdot \log_2 L$$

when  $N_2 \approx N_L$

We can easily calculate the speed ratio of radix- $L$  FFT to radix-2 FFT as shown in Table 2.3.



algorithm	radix-2 FFT	radix-3 FFT	radix-4 FFT	radix-5 FFT
speed ratio	1	$1.02 \cdot \log_2 3$	$1.01 \cdot \log_2 4$	$0.94 \cdot \log_2 5$

**Table 2.3: The speed ratio of radix- $L$  FFT to radix-2 FFT.**

Obviously, the computation speed ratio of radix- $L$  FFT to radix-2 FFT is very close to  $\log_2 L$  which is larger than one. So, radix-3 FFT, radix-4 FFT and radix-5 FFT have a better performance than does the radix-2 FFT.

## 2.5 Numbers of Sampling Points

As discussed, radix- $L$  FFT's have a speed performance improvement over the radix-2 FFT, there is another issue which needs to be considered. That is the number of numbers of sampling points which can be applied by FFT's. In the following, we will discuss this issue.

For the radix-2 FFT, the number of sampling points which can be applied by the FFT must

be a number equal to  $2^k$ , where  $k$  is a natural number. Similarly for the radix-3 FFT, the number of sampling points must be one of the series  $3^k$ . Table 2.4 lists the numbers of sampling points that can be applied by the radix-3 FFT under the value 10000. There are only 8 numbers of sampling points. This is fewer than that of the radix-2 FFT, which is 13. The higher radix FFT, the fewer this number. For radix-4 and radix-5 FFT, the number becomes 6 and 5 respectively. This is a drawback for using a higher radix FFT.

$k$	1	2	3	4	5	6	7	8
$3^k$	3	9	27	81	243	729	2187	6561
distance	-	6	18	54	162	486	1458	4374

**Table 2.4: The numbers of sampling points which can be applied by the radix-3 FFT under the value 10000.**



## Chapter 3 Mixed-Radix FFT

As we have discussed in the previous chapter, the FFT's other than radix-2 have the advantages of speed improvement but have the drawback that the numbers of sampling points become less. A good solution for bypassing the drawback is that we can apply FFT's for a set of sampling points by dividing the computation stages into several groups and apply the FFT of different radix for each group of computation stages. We call this the "mixed radix-FFT" algorithm.

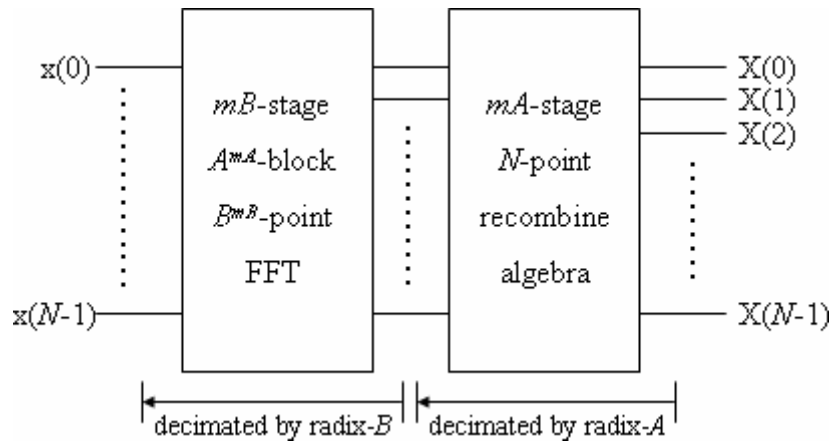
In this chapter, we will derive the mixed-radix FFT based on the results obtained in the previous chapter. First we will discuss the case of two radices, i.e., radix- $A/B$  FFT, and then the case of three radices, i.e., radix- $A/B/C$  FFT.

### 3.1 Radix- $A/B$ FFT

Radix- $A/B$  FFT has two factors  $A$  and  $B$ . According to the order of permutation of these two factors, there will be several forms of algorithm that can be used to decimate the data sequence. For example, if the data sequence has  $N (= A^2 \times B)$  points, there will be three forms of permutation of  $A$  and  $B$ , i.e.,  $AAB$ ,  $ABA$  and  $BAA$ . We only consider the permutation  $AA \dots ABB \dots B$  since all other permutations can be applied the same analysis.

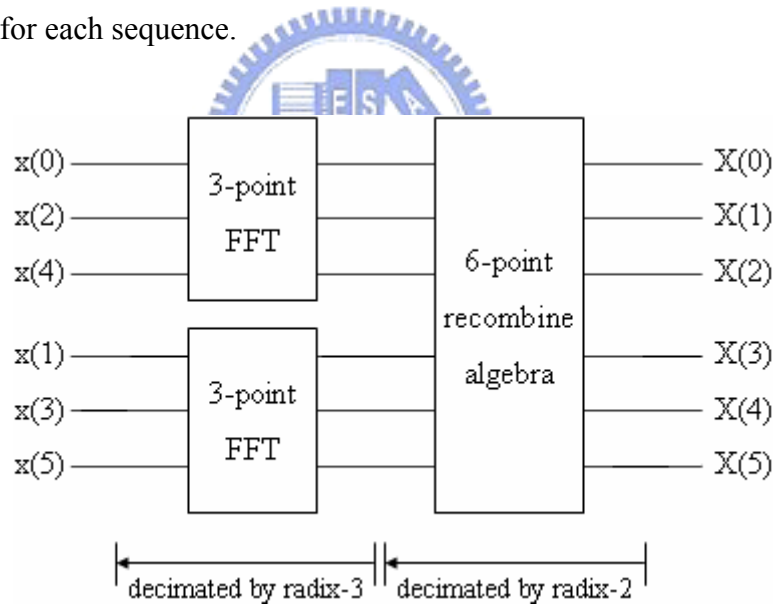
#### 3.1.1 Decimation-In-Time Algorithm

Let us consider the data sequence  $x(n)$  with  $N (= A^{m_A} \times B^{m_B})$  sampling points. First we decimate the data sequence by a factor of  $A$  repeatedly until the data sequence is split into  $A^{m_A}$  data sequences of which each has  $B^{m_B}$  sampling points. And then we decimate each of these data sequences further by a factor of  $B$  repeatedly until these data sequences each is a series of  $B$ -point FFT's. This process can be shown as in Figure 3.1.



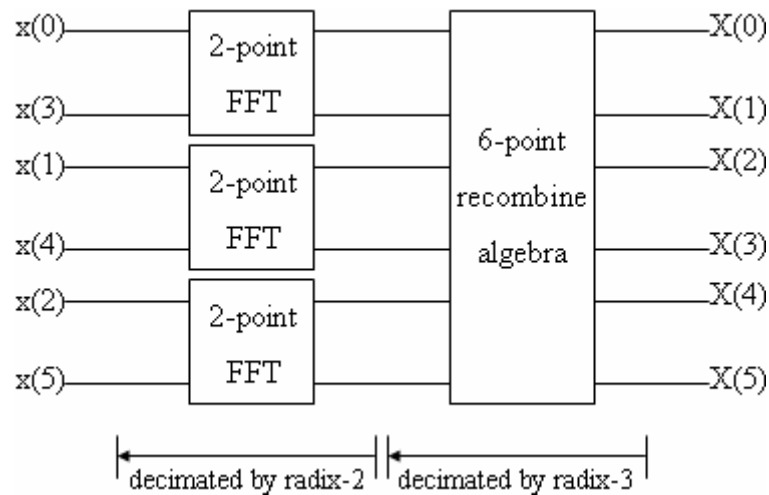
**Figure 3.1: The radix-A/B FFT decimation.**

In the figure, ‘ $mA$ -stage’ means the data sequence is decimated by the radix-A algorithm ( $mA-1$ ) times, ‘ $mB$ -stage’ means the data sequences are decimated by the radix-B algorithm ( $mB-1$ ) times, and ‘ $A^{mA}$ -block and  $B^{mB}$ -point’ means there are  $A^{mA}$  data sequences with  $B^{mB}$  sampling points for each sequence.



**Figure 3.2: The 6-point FFT computed by the radix-2/3 algorithm.**

For an example, Figure 3.2 depicts the 6-point FFT computed by the radix-2/3 algorithm which is stated above. We observe that the data sequence is finally decimated into two 3-point FFT’s. Besides using the radix-2/3 algorithm, we can also use the radix-3/2 algorithm to do the same 6-point FFT. For this case, the data sequence is finally decimated into three 2-point FFT’s as shown in Figure 3.3.



**Figure 3.3: The 6-point FFT computed by the radix-3/2 algorithm.**

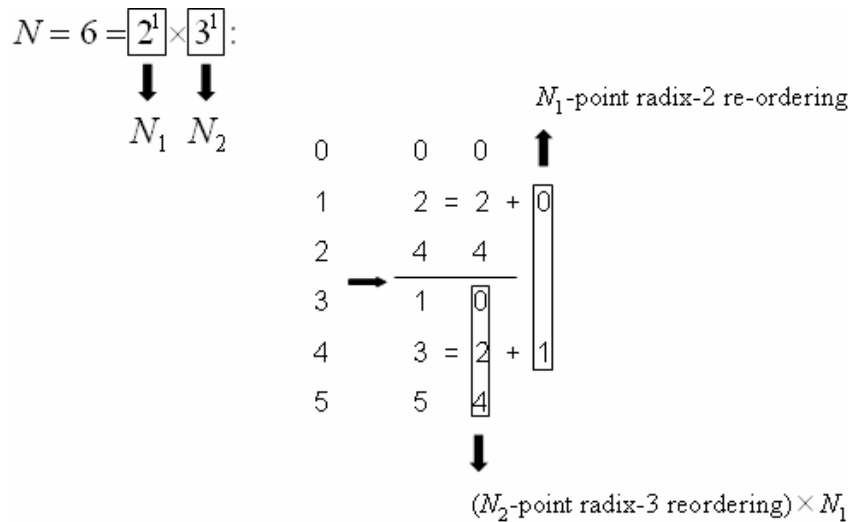
### 3.1.2 Re-Ordering Algorithm for Input Sequence

The re-ordering algorithm used for the radix- $A/B$  FFT is very different from that used for the single-radix FFT. Since there are two factors in the algorithm, the previous bit-reversal procedure can not be used directly.

Here we consider the process of the decimation-in-time algorithm. We first use the radix- $A$  to decimate the data sequence to reach several blocks. We observe the order of the data sequence in every block and try to find a mathematical relation to describe such result. And then we use the radix- $B$  to decimate each block which has  $B^{mB}$  sampling points. Before decimating, we re-assign numbers beginning from 0 to the data sequences in each block and record the true order and the new number of each data. After this, we use the bit-reversal (radix- $B$  re-ordering) procedure to re-order the data in each block. Finally, we can obtain the input order according to the previous record.

However the above method is a little bit complicated. We can directly find a mathematical relation to describe the input order. The following Figure 3.4 depicts the 6-point radix-2/3 re-ordering algorithm. This algorithm can be extended to radix- $A/B$ .





**Figure 3.4: The 6-point radix-2/3 re-ordering algorithm.**

In the figure, ‘0, 1, 2, 3, 4, 5’ means the original order, ‘0, 2, 4, 1, 3, 5’ means the decimated order and the part written on the right of the equal sign is the re-ordering algorithm. ‘ $N_1$ -point radix-2 re-ordering’ means using bit-reversal (radix-2 re-ordering) to re-order the sequence of numbers  $0 \dots N_1 - 1$  where  $N_1$  is of the power of 2. ‘ $(N_2$ -point radix-3 re-ordering)  $\times N_1$ ’ means using bit-reversal (radix-3 re-ordering) to re-order the sequence of numbers  $0 \dots N_2 - 1$  where  $N_2$  is of the power of 3, and then multiplied by the sequence  $N_1$ .

One thing to be noted is that: For applying the radix- $A/B$  FFT or the radix- $B/A$  FFT, the order of  $A$  and  $B$  will give different run-time. Table 3.1 compares the computation speed for the radix- $A/B$  FFT and the radix- $B/A$  FFT under the different conditions.

	$A^{m^A} > B^{m^B}$	$B^{m^B} > A^{m^A}$
radix- $A/B$ FFT	slow	fast
radix- $B/A$ FFT	fast	slow

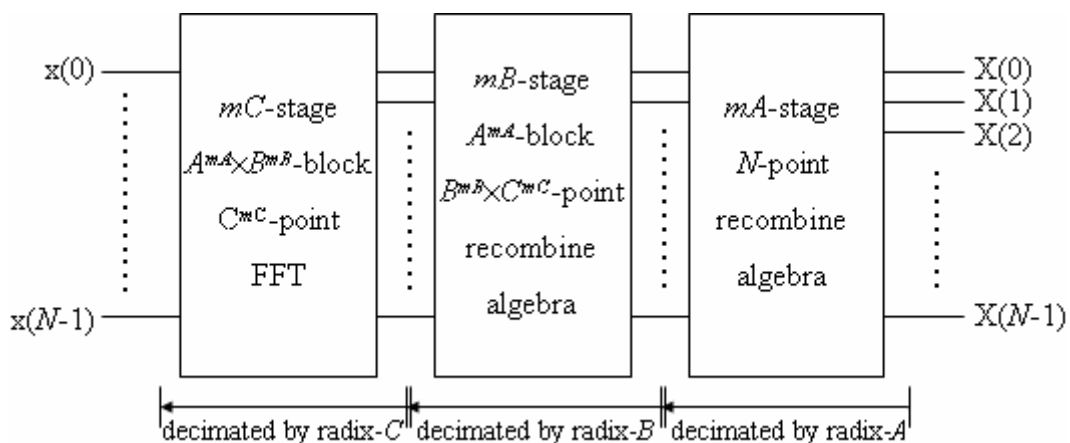
**Table 3.1: A comparison of the radix- $A/B$  FFT and the radix- $B/A$  FFT.**

### 3.2 Radix- $A/B/C$ FFT

For the three radices  $A$ ,  $B$  and  $C$ , there are permutations of radix- $A/B/C$  FFT, radix- $A/C/B$  FFT, radix- $B/A/C$  FFT, radix- $B/C/A$  FFT, radix- $C/A/B$  FFT and radix- $C/B/A$  FFT. Here we only discuss for the case of the radix- $A/B/C$  FFT, and all other permutations of radix order are the same.

### 3.2.1 Decimation-In-Time Algorithm

Let us consider the data sequence  $x(n)$  with  $N (= A^{m_A} \times B^{m_B} \times C^{m_C})$  sampling points. First the data sequence is decimated by a factor of  $A$  repeatedly until the data sequence is split into  $A^{m_A}$  data sequences for which each sequence has  $(B^{m_B} \times C^{m_C})$  sampling points. And then each of these data sequences is decimated by the factor of  $B$  repeatedly until each sequence is split into  $B^{m_B}$  data sequences for which each sequence has  $C^{m_C}$  sampling points. Hence, now we have altogether  $(A^{m_A} \times B^{m_B})$  data sequences for which each sequence has  $C^{m_C}$  sampling points. Finally, each of these sequences is decimated by the factor of  $C$  repeatedly until it reaches a series of  $C$ -point FFT's. In other words, the radix- $A/B/C$  FFT uses the radix- $A$  algorithm, the radix- $B$  algorithm and the radix- $C$  algorithm in turn to decimate the DFT computation as shown in Figure 3.5.

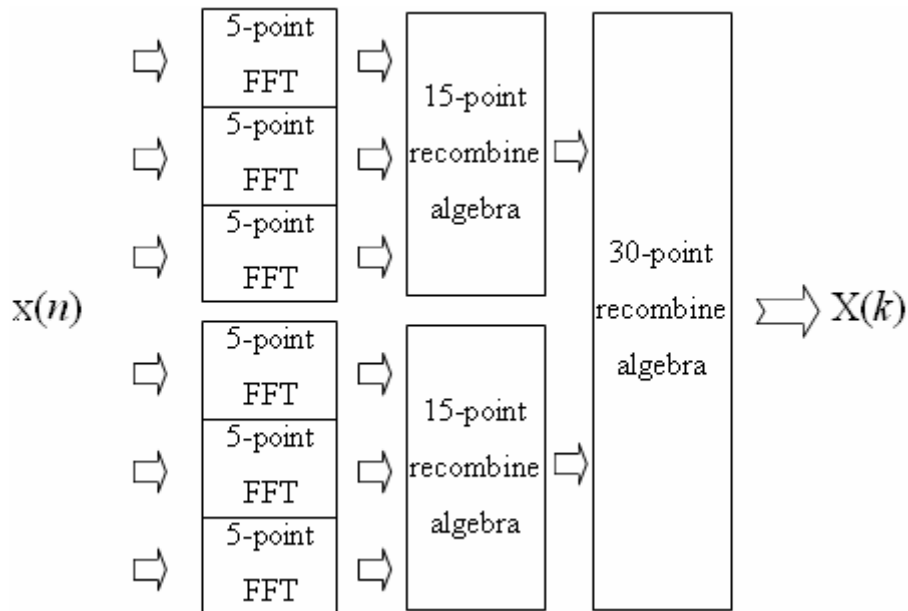


**Figure 3.5: The Radix- $A/B/C$  FFT decimation.**

In the figure, ' $m_A$ -stage' means the data sequence is decimated by the radix- $A$  algorithm

( $mA-1$ ) times, ' $mB$ -stage' means the data sequences are decimated by the radix- $B$  algorithm ( $mB-1$ ) times, ' $mC$ -stage' means the data sequences are decimated by the radix- $C$  algorithm ( $mC-1$ ) times, ' $A^{mA} \times B^{mB}$ -block and  $C^{mC}$ -point' means there are  $(A^{mA} \times B^{mB})$  data sequences with  $C^{mC}$  sampling points for each sequence and ' $A^{mA}$ -block and  $(B^{mB} \times C^{mC})$ -point' means there are  $A^{mA}$  data sequences with  $(B^{mB} \times C^{mC})$  data points for each sequence.

For example, Figure 3.6 depicts the 30-point FFT computed by the radix-2/3/5 algorithm. We observe that the data sequence is decimated into two 15-point FFT's first and then these two 15-point FFT's are decimated into three 5-point FFT's, respectively.

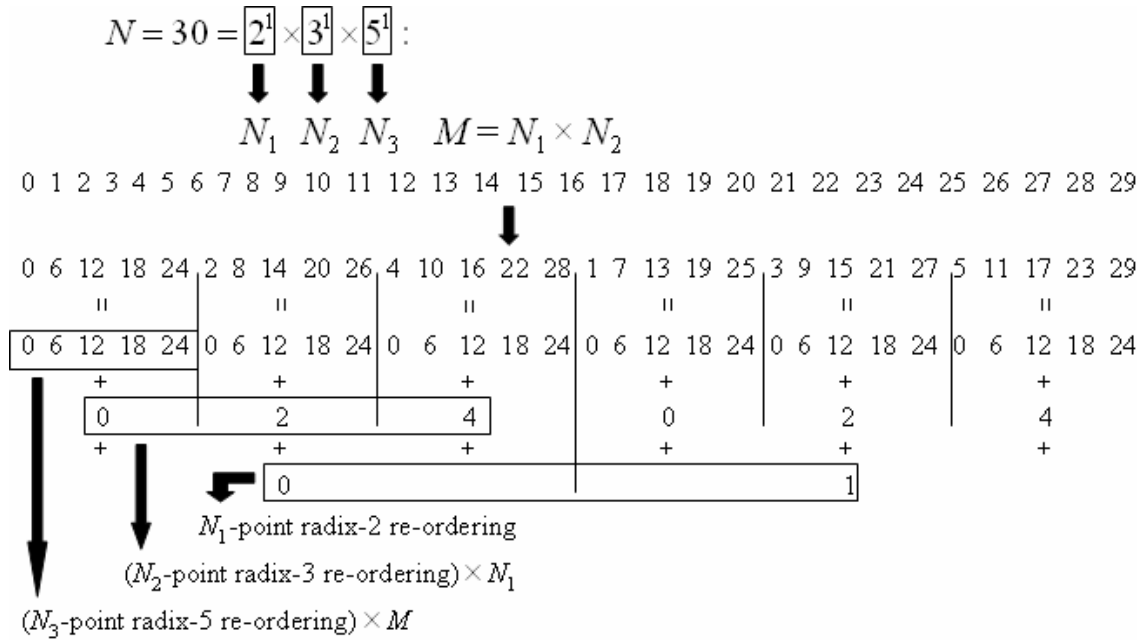


**Figure 3.6: The 30-point FFT computed by the radix-2/3/5 algorithm.**

This 30-point FFT can also be done by using other five algorithms of different order of 2, 3 and 5. Although the structures of computation and the order of input sequences are different, the obtained results are still the same.

### 3.2.2 Re-Ordering Algorithm for Input Sequence

Similar to section 3.1.2, here again we will find a mathematical relationship to describe the input order directly. We use radix-2/3/5 as an example to re-order a 30-point sequence in Figure 3.7. The procedure can be extended to the general case of the radix- $A/B/C$ .



**Figure 3.7: The illustration of a 30-point radix-2/3/5 re-ordering algorithm.**

In the figure, the first sequence is in the original order, the second sequence is in the decimated order and the part written under the equal sign is the re-ordering algorithm. ‘ $N_1$ -point radix-2 re-ordering’ means using bit-reversal (radix-2 re-ordering) to re-order the sequence of numbers  $0 \dots N_1 - 1$  where  $N_1$  is of the power of 2, and ‘ $(N_2$ -point radix-3 re-ordering)  $\times N_1$ ’ and ‘ $(N_3$ -point radix-5 re-ordering)  $\times M$ ’ are similar to those discussed in section 3.1.2.

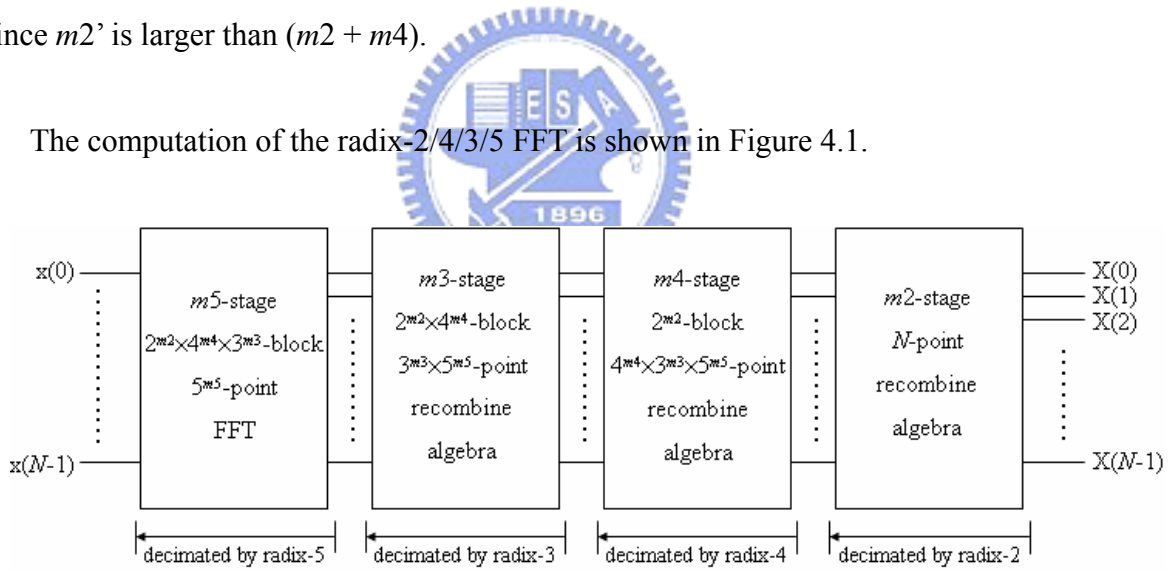
## Chapter 4 Radix-2/4/3/5 FFT and Interpolation

Now we can use the radix-2/3/5 FFT to replace the radix-2 FFT, but we still want to speed up the radix-2/3/5 FFT and extend the numbers of sampling points to more numbers. To do this, we will replace the radix-2 algorithm in the radix-2/3/5 FFT with the radix-2/4 algorithm, thus we will obtain the radix-2/4/3/5 FFT. Furthermore, we will add the concept of interpolation to extend the applicable numbers of sampling points to any number.

### 4.1 Radix-2/4/3/5 FFT Algorithm

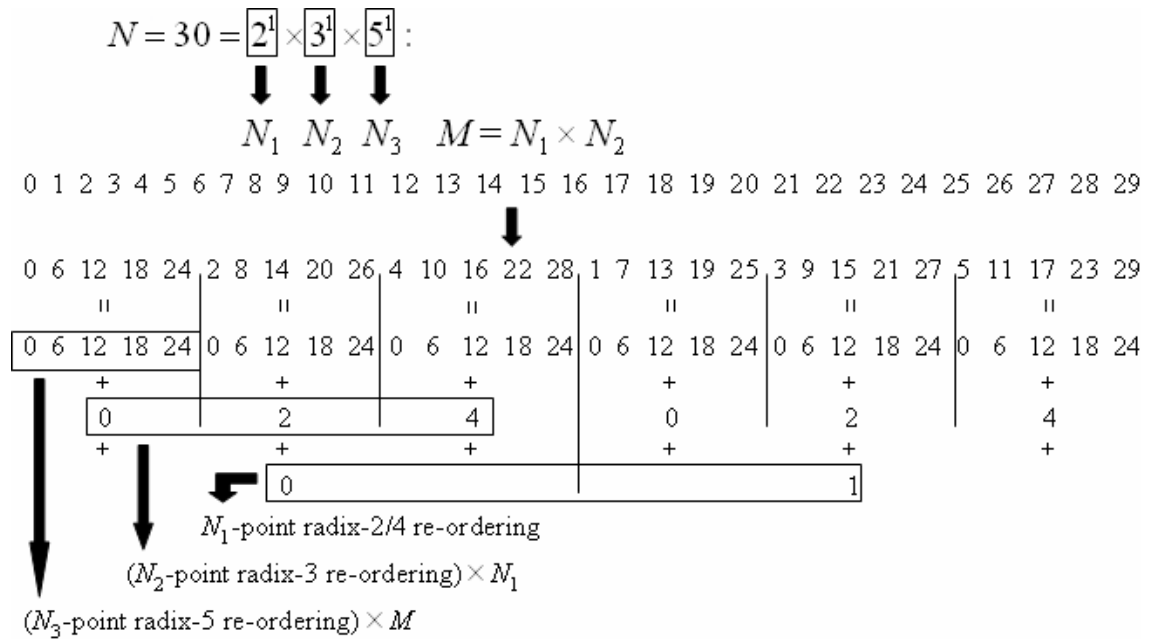
For a data sequence  $x(n)$  with  $N (= 2^{m_2'} \times 3^{m_3} \times 5^{m_5})$  sampling points, we replace the radix-2 algorithm with the radix-2/4 algorithm, in other words, we consider the data sequence  $x(n)$  to be  $N (= 2^{m_2} \times 4^{m_4} \times 3^{m_3} \times 5^{m_5})$  sampling points. This reduces the number of computation stages since  $m_2'$  is larger than  $(m_2 + m_4)$ .

The computation of the radix-2/4/3/5 FFT is shown in Figure 4.1.



**Figure 4.1: The radix-2/4/3/5 FFT decimation.**

As for the re-ordering algorithm, it is the same as that of Figure 3.8, except that the  $N_1$ -point radix-2 re-ordering must be replaced with the  $N_1$ -point radix-2/4 re-ordering as shown in Figure 4.2.



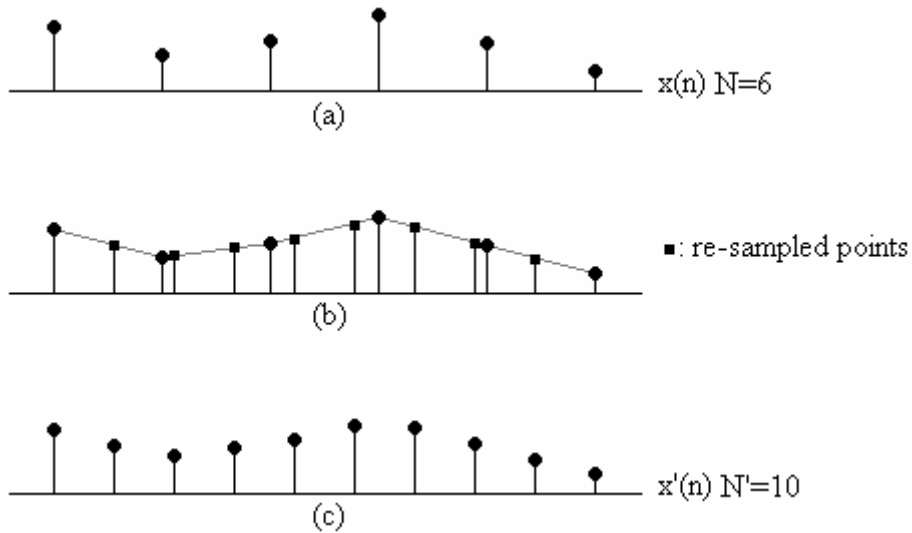
**Figure 4.2: The 30-point radix-2/4/3/5 re-ordering algorithm.**

In the figure, the ‘radix-2/4 re-ordering’ is the radix- $A/B$  re-ordering algorithm discussed in section 3.1.2.



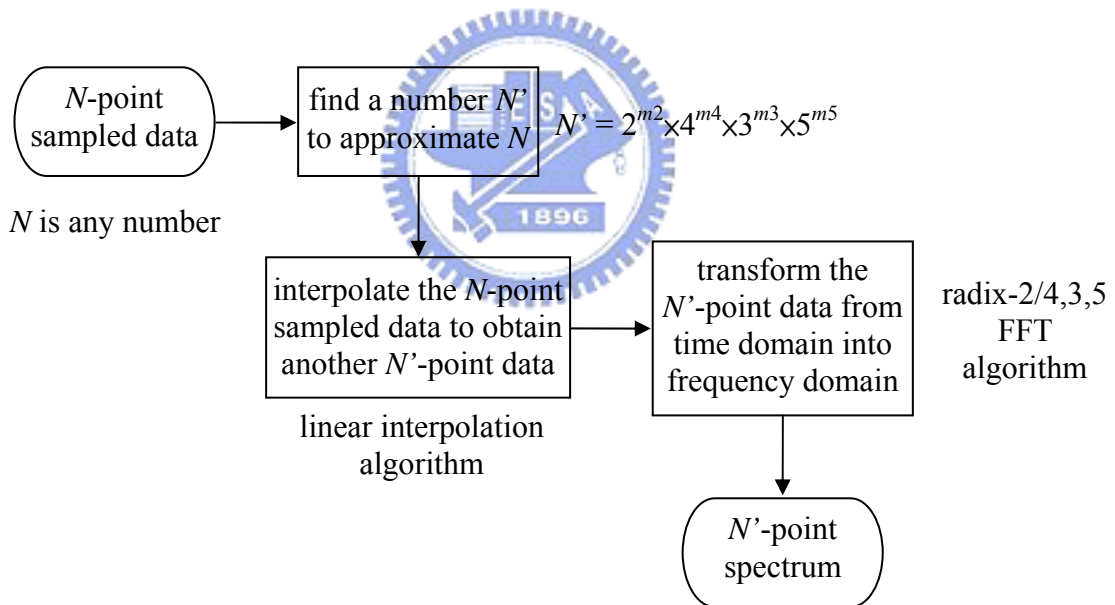
## 4.2 Interpolation Algorithm

The purpose of this algorithm is that it is to re-sample the sampled data so that the number of sampling points can be applied by the radix-2/4/3/5 FFT. For example, if there is a data sequence  $x(n)$  for a signal with  $N$  ( $=$  any number) sampling points, interpolation can re-sample  $x(n)$  and then we can obtain another similar sequence  $x'(n)$ , which expresses the same signal, with  $N'$  ( $= 2^{m_2} \times 4^{m_4} \times 3^{m_3} \times 5^{m_5}$ ) sampling points. In order to shorten the run-time of interpolation and reduce the difference between  $x(n)$  and  $x'(n)$ , we use the linear interpolation which is shown in Figure 4.3. In the figure, there is a data sequence  $x(n)$  with 6 sampling points in (a). We use a straight line passing through two points of  $x(n)$  to approximate the actual curve and then represent the interpolated signal with a new set of sequence  $x'(n)$  with 10 sampling points placed in equal time interval as shown in (c).



**Figure 4.3: The illustration of linear interpolation.**

The following Figure 4.4 depicts the flow chart of an  $N$ -point FFT operated by the new algorithm.



**Figure 4.4: The flow chart of  $N$ -point FFT operated by the new algorithm.**

### 4.3 Simulation Results

To illustrate the advantage of the radix-2/4/3/5 FFT on the applicable number of sampling points, we do the following experiment: Within the range from 1 to 10000, we choose 50 groups of  $N_s$  evenly. The numbers of sampling points ( $N_s$ ) and the data sequences are

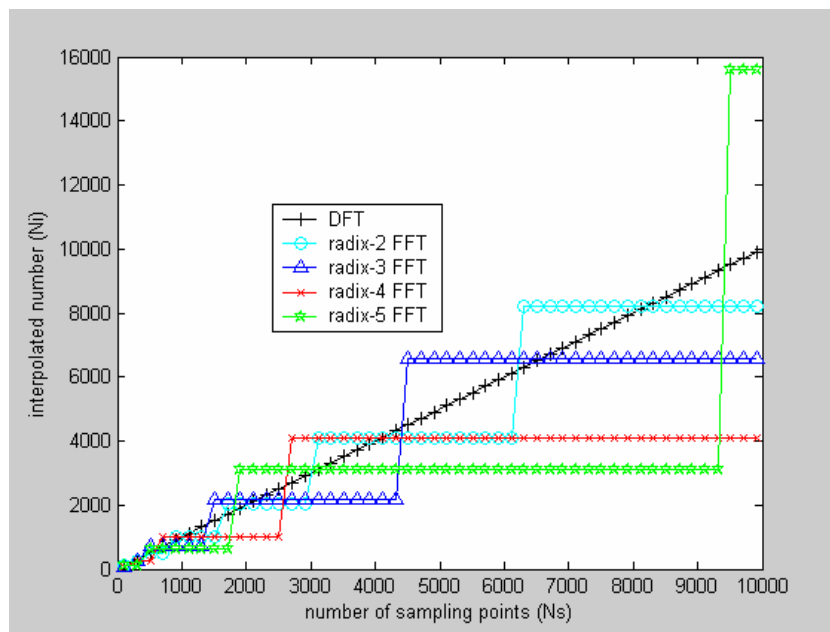
obtained from the procedure as follows:

```

for (i=1:50)
    Ns(i)=111+(i-1)*200;
    for (j=1:Ns(i))
        x(i,j)=j;
    end
end
end

```

We plot the required numbers of sampling points for which the FFT algorithms can be applied in terms of the number of sampling points for the single-radix FFT's in Figure 4.5 and for the mixed-radix FFT's in Figure 4.6.

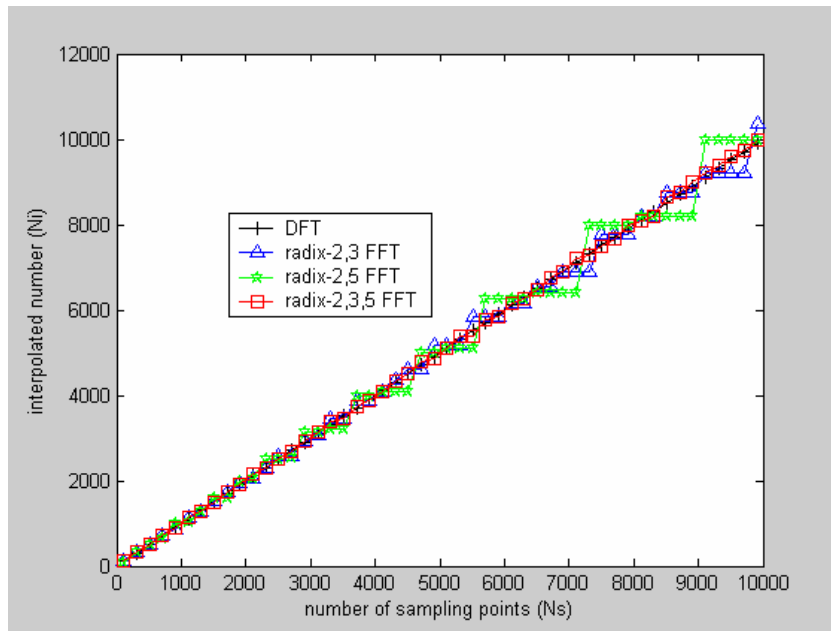


**Figure 4.5: The comparison of applicable number ( $N_i$ ) of sampling points which can be applied by the single-radix FFT's.**

In the figure, 'DFT' means using discrete Fourier transform which can be applied with any number of sampling points. Thus 'DFT' is a straight line. The radix-2 FFT has the most numbers of sampling points than any other single-radix FFT as discussed in section 2.5 since it has the minimum step-level in the figure. The larger the decimating factor is, the larger the

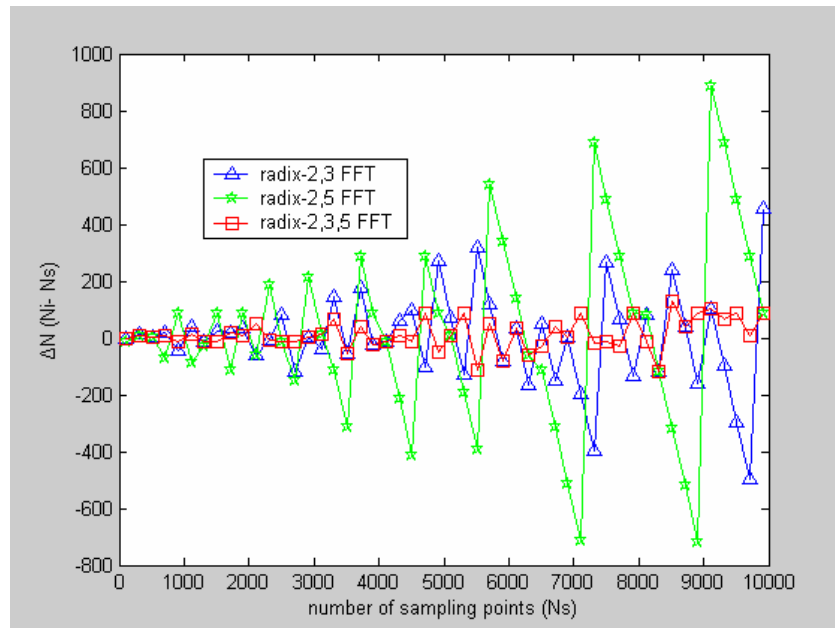


step-level is. Hence, it is not a good idea to use the single-radix FFT's.



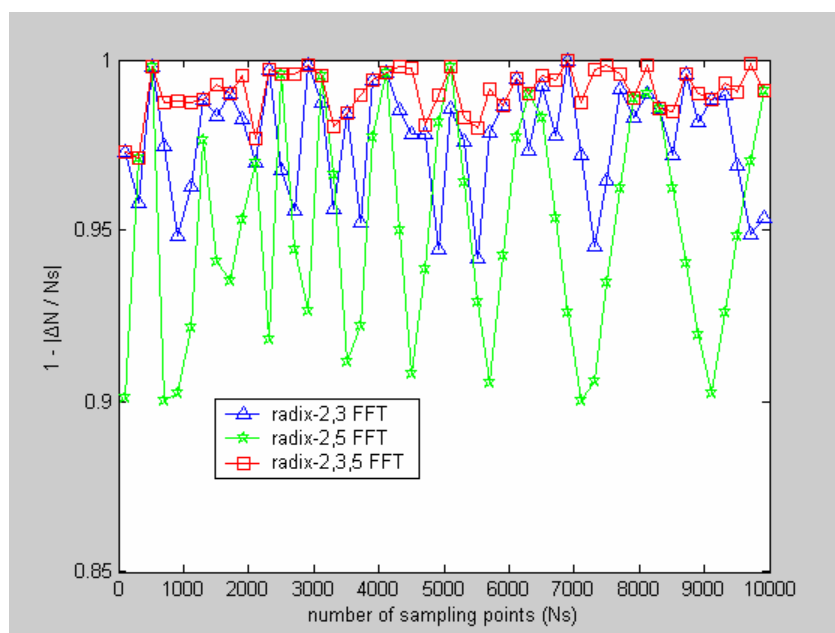
**Figure 4.6: The comparison of applicable number ( $N_i$ ) of sampling points which can be applied by the mixed-radix FFT's.**

Figure 4.6 shows the similar plots for the mixed-radix FFT's. In the figure, 'radix-2,3 FFT' includes two algorithms, i.e., radix-2/3 FFT and radix-3/2 FFT, since they have the same number of re-sampled points. And so does for the 'radix-2,5 FFT' and 'radix-2,3,5 FFT'. From the figure, we can find that 'radix-2,3 FFT', 'radix-2,5 FFT' and 'radix-2,3,5 FFT' all follow the straight line of 'DFT', and as  $N_s$  is larger, the step-levels of 'radix-2,3 FFT' and 'radix-2,5 FFT' are larger too. As for the 'radix-2,3,5 FFT', it has the smallest step-level and approaches that of the 'DFT'.



**Figure 4.7: The comparison of  $\Delta N (N_i - N_s)$  for different mixed-radix FFT's.**

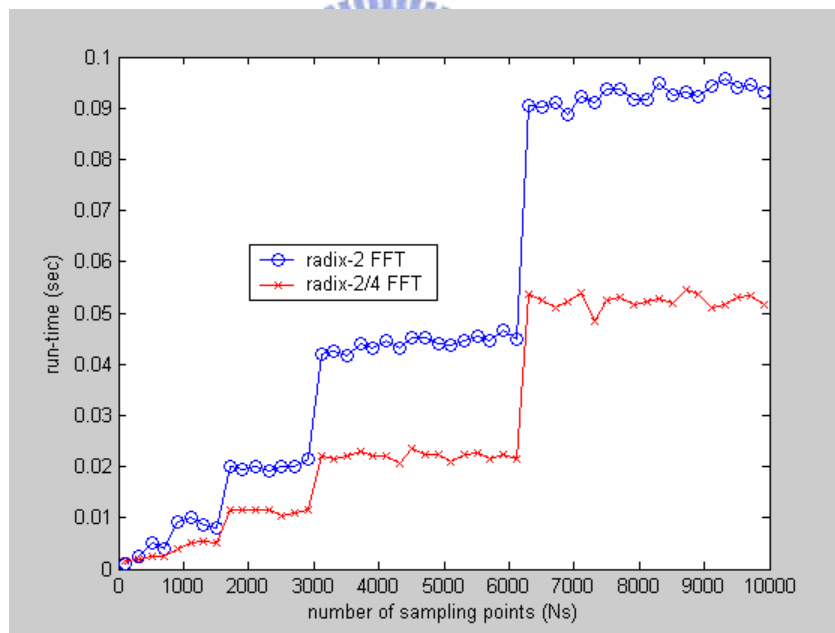
Figure 4.7 shows the  $\Delta N (N_i - N_s)$  plots which are derived from data from Figure 4.6 for each mixed-radix FFT. Obviously, 'radix-2,5 FFT' has the maximum amplitude and 'radix-2,3,5 FFT' has the minimum amplitude. To further demonstrate the above property is to plot the data in the form of  $1 - |\Delta N / N_s|$  versus the number of sampling points as shown in Figure 4.8.



**Figure 4.8: The comparison of  $1 - |\Delta N / N_s|$  for different mixed-radix FFT's.**

Since interpolation modifies the original signal and introduces effectively noise to the signal, it is interesting to investigate how the S/N of the signal will be affected by the interpolation process. The  $1 - |\Delta N/N_s|$  figure reflects indirectly the S/N of the processed signal. In the figure, it can be seen that the radix-2,3,5 FFT has the best S/N than other two algorithms. However, as for other two algorithms, their S/N's are not too bad since their values of  $1 - |\Delta N/N_s|$  are all larger than 0.9.

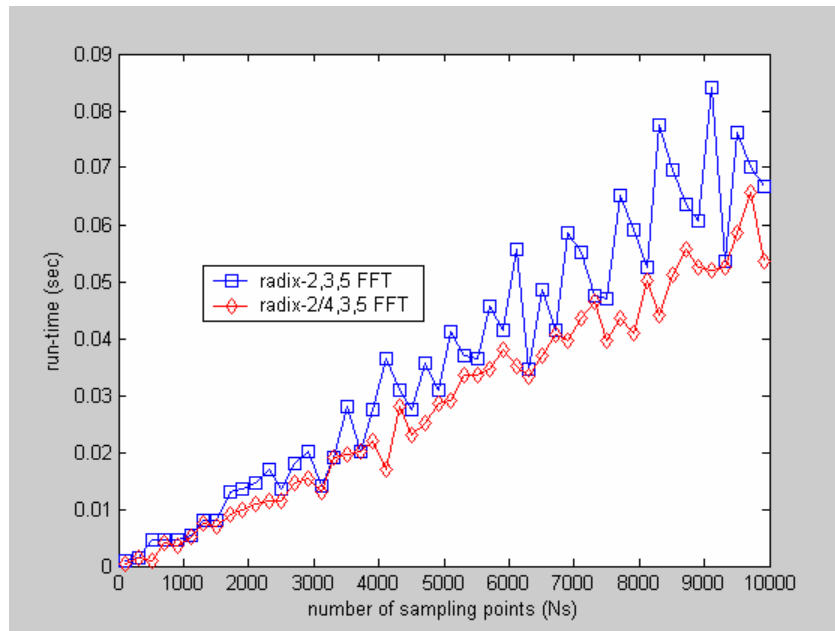
Next, we will compare the run-time for different FFT algorithms. First we compare the run-times of the radix-2 FFT and the radix-2/4 FFT as shown in Figure 4.9 and then compare the run-times of the radix-2,3,5 FFT and the radix-2/4,3,5 FFT as shown in Figure 4.10 since they have the same numbers of re-sampled points, respectively.



**Figure 4.9: The comparison of run-times of radix-2 FFT and radix-2/4 FFT.**

Obviously, the radix-2/4 FFT is faster than radix-2 FFT and the run-time of radix-2/4 FFT is almost half of that of radix-2 FFT. This result is consistent with that of Table 2.3. In this experiment, for the radix-2/4 FFT, only the first stage is decimated by the factor 2, and all other stages are decimated by the factor 4. So, the radix-2/4 FFT is almost the same as the

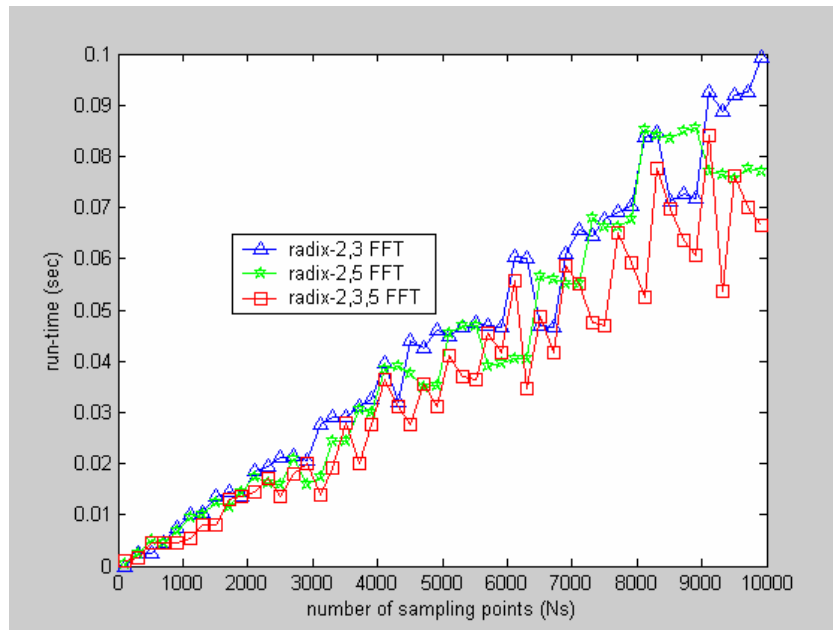
radix-4 FFT, except that the number of sampling points that can be applied to radix-2/4 FFT is of the power of 2, not of the power of 4.



**Figure 4.10: The comparison of run-times of radix-2,3,5 FFT and radix-2/4,3,5 FFT.**

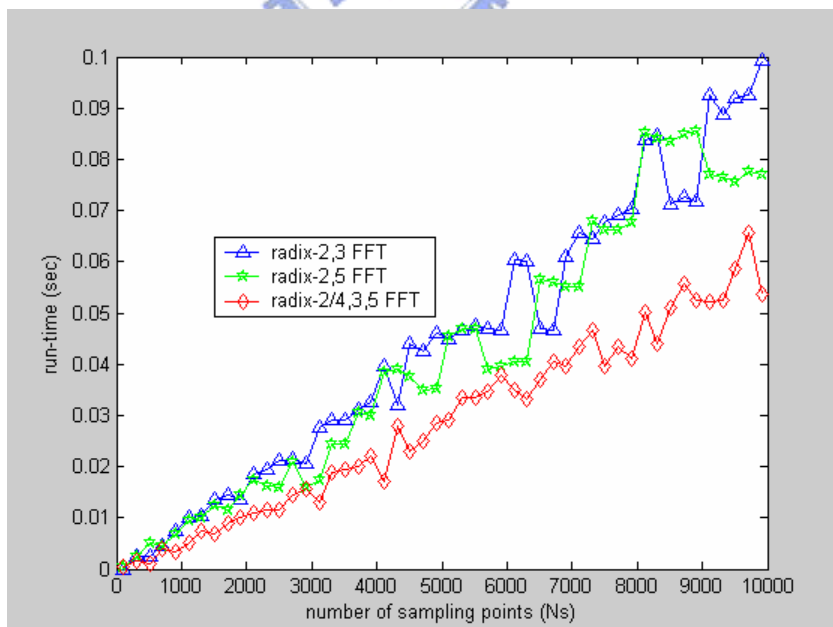
In this experiment, the radix-2/4,3,5 FFT was obtained by replacing the radix-2 algorithm of the radix-2,3,5 FFT with the radix-2/4 algorithm. From this figure, we can see that the radix-2/4,3,5 FFT is faster than the radix-2,3,5 FFT, but there are some positions that these two algorithms overlap. This is because the numbers of re-sampled points are  $(2 \times 3^{m^3} \times 5^{m^5})$  where the radix-2/4,3,5 FFT is the same as the radix-2,3,5 FFT.

The comparison of run-times of radix-2,3 FFT, radix-2,5 FFT and radix-2,3,5 FFT that have similar numbers of re-sampled points is shown in Figure 4.11 . Obviously, radix-2,3 FFT and radix-2,5 FFT have similar performance of computation speed but the radix-2,3,5 FFT has a better computation speed performance than the above two algorithms generally.



**Figure 4.11: The comparison of run-times of radix-2,3 FFT, radix-2,5 FFT and radix-2,3,5 FFT.**

And the comparison of run-times of radix-2,3 FFT, radix-2,5 FFT and radix-2/4,3,5 FFT is shown in Figure 4.12. Obviously, radix-2/4,3,5 FFT is much better than other two algorithms.



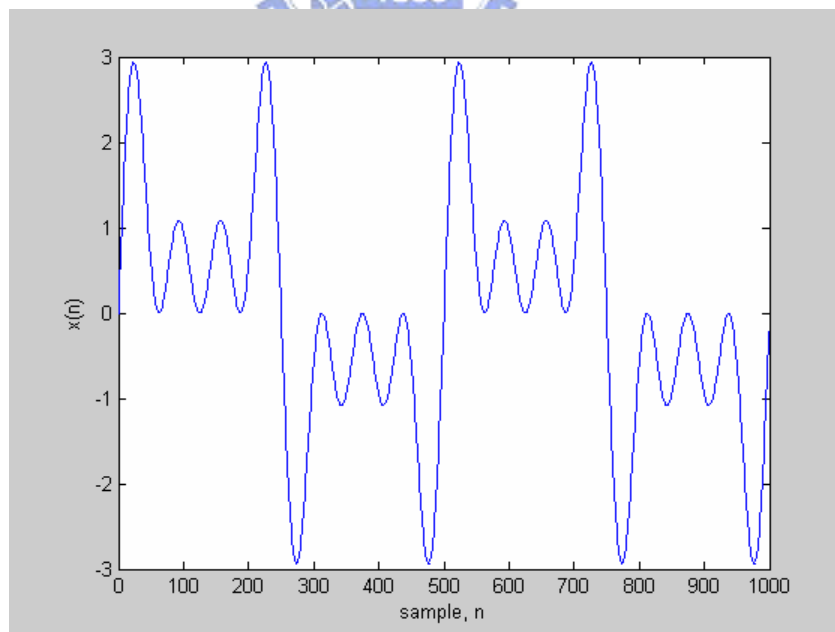
**Figure 4.12: The comparison of run-times of radix-2,3 FFT, radix-2,5 FFT and radix-2/4,3,5 FFT.**

As a conclusion, based on results of Figures 4.6, 4.7, 4.8 and 4.12, we can conclude that the radix-2/4,3,5 FFT cooperating with interpolation is the best algorithm in doing FFT which satisfy both the requirements of the performance of computation speed and the applicable numbers of sampling points.

## 4.4 Examples

In this section we further use two examples applied with the above developed FFT algorithms to show their effectiveness. One example is to transform a multi-tone signal which is often used as a test input for a mixed signal testing, and the other is a square wave signal applied with the algorithm. The results are also shown in terms of the applicable number of sampling points, the run-time, and the introduced noise due to the interpolating the data points on the signal.

### 4.4.1 Multi-tone Signal



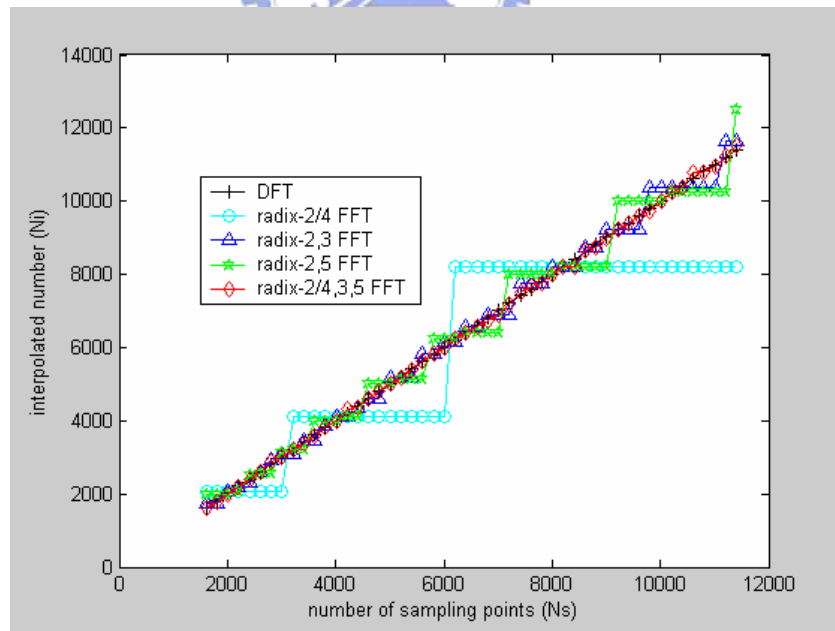
**Figure 4.13: The waveform of the multi-tone signal.**

This signal contains 4 tones, the frequency and magnitude of each tone and the sampling

frequency are shown as follows and the signal is plotted in Figure 4.13. The frequencies are described in the form of bin number.

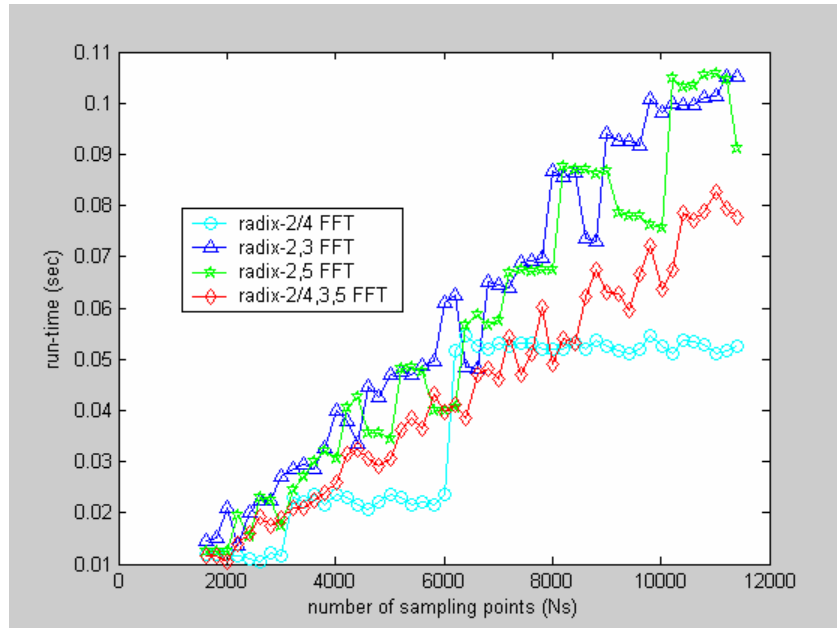
```
f=[1 3 5 7]*115;
m=[1 1 1 1];
nyquist=max(f)*2;
for (i=1:50)
    fs(i)=nyquist+1+(i-1)*200;
end
```

We plot the required numbers of sampling points for which the FFT algorithms can be applied in terms of the number of sampling points for the FFT's in Figure 4.14. Still, we can see that the radix-2/4,3,5 algorithm is the best algorithm which can take the most of number of applicable data points while other algorithms all have limited number of applicable data points.



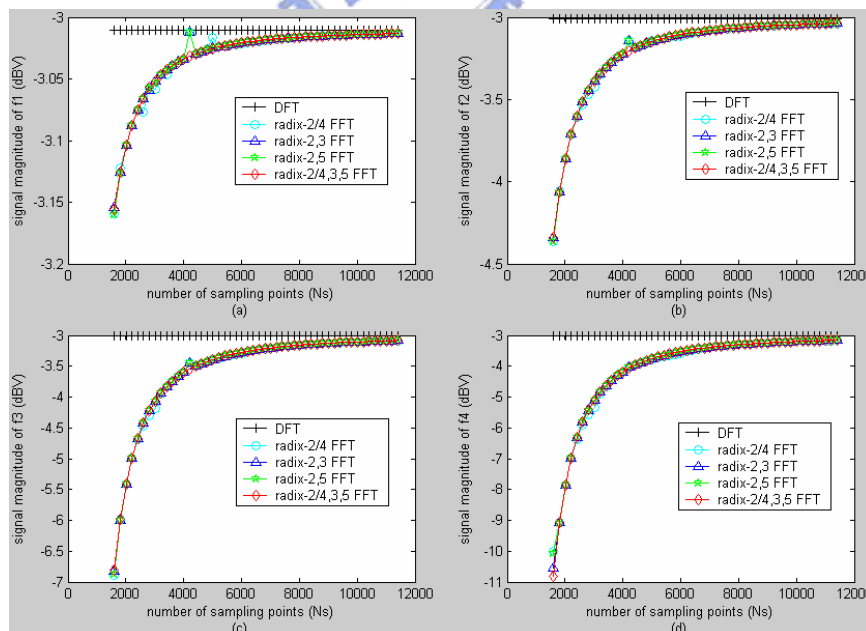
**Figure 4.14: The comparison of applicable number ( $N_i$ ) of sampling points which can be applied by the FFT's.**

Next, we will compare the run-time for different FFT algorithms as shown in Figure 4.15.



**Figure 4.15: The comparison of run-times of different FFT algorithms.**

In the figure, we can see that the radix-2/4,3,5 algorithm has the best run-time among all the algorithms except the radix-2/4 algorithm. The radix-2/4 algorithm has the least run-time because the number of data points are less than that of the radix-2/4,3,5 algorithm.

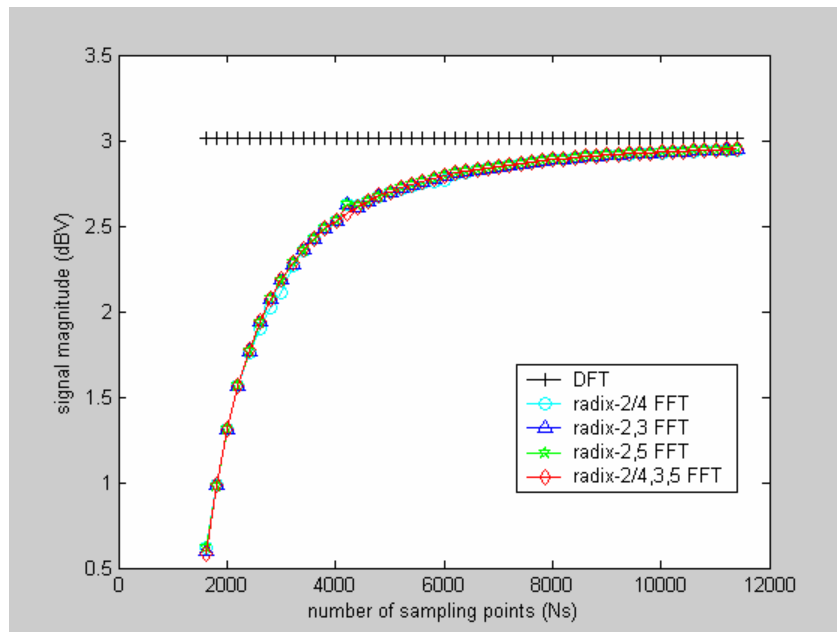


**Figure 4.16: The comparisons of magnitudes of each tone of different FFT algorithms**

**(a) 1<sup>st</sup> tone, (b) 2<sup>nd</sup> tone, (c) 3<sup>rd</sup> tone, and (d) 4<sup>th</sup> tone.**



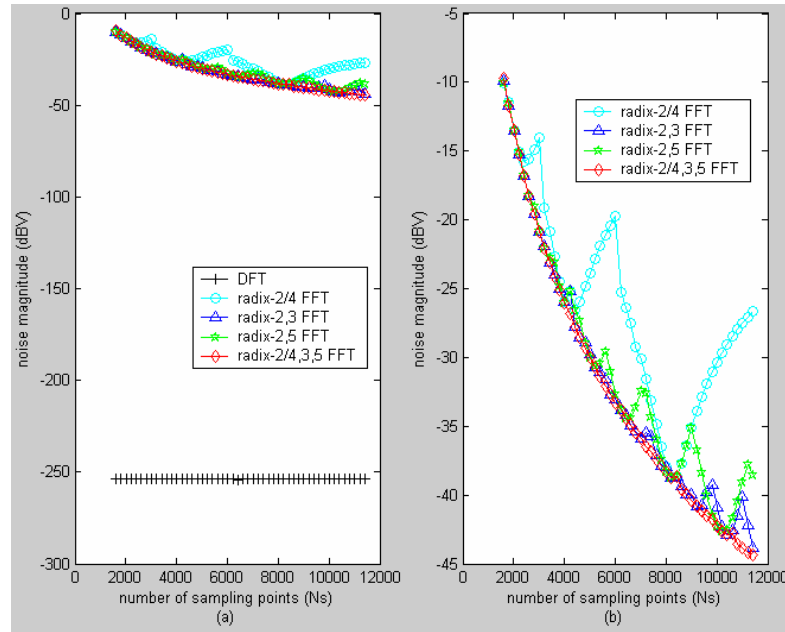
Figure 4.16 plots the obtained magnitudes of each tone in terms of the number of data points for each algorithm. For reference, the magnitudes obtained by using the DFT, which can be considered to be the ideal case transformation, are also plotted. It is seen that the magnitudes of the tones obtained through all FFT algorithms increase with the number of data points since more data points more accurate the obtained values. Figure 4.17 shows the combined magnitude of the four tones where we see the same trend.



**Figure 4.17: The comparison of magnitudes of signal of different FFT algorithms.**

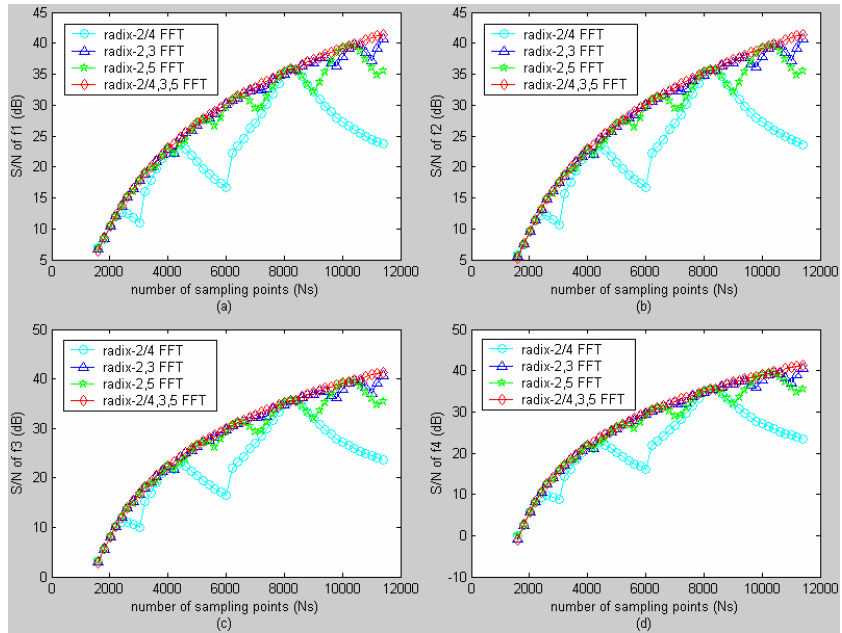
Figure 4.18 plots the introduced noises in terms of the number of data points due to interpolating process for each algorithm. Figure 4.18 (a) also includes the noise obtained for the original four-tone signal if the DFT algorithm is used as the transforming algorithm. It is seen that for all FFT algorithms, much noises are introduced on the transformed spectrums due to the interpolating process which makes the interpolated signal different from the original signal. Figure 4.18 (b) shows the expanded plot of the Figure 4.18 (a)'s plot to show the noise introduced for each FFT algorithm. It is seen that for all algorithms, except the radix-2/4,3,5 one, the noises oscillate with the number of data points since the number of

interpolating data points increase periodically with respect to the number of interpolated data points. The introduced noise for the radix-2/4,3,5 algorithm decreases monotonically with respect to the number of interpolated data points.

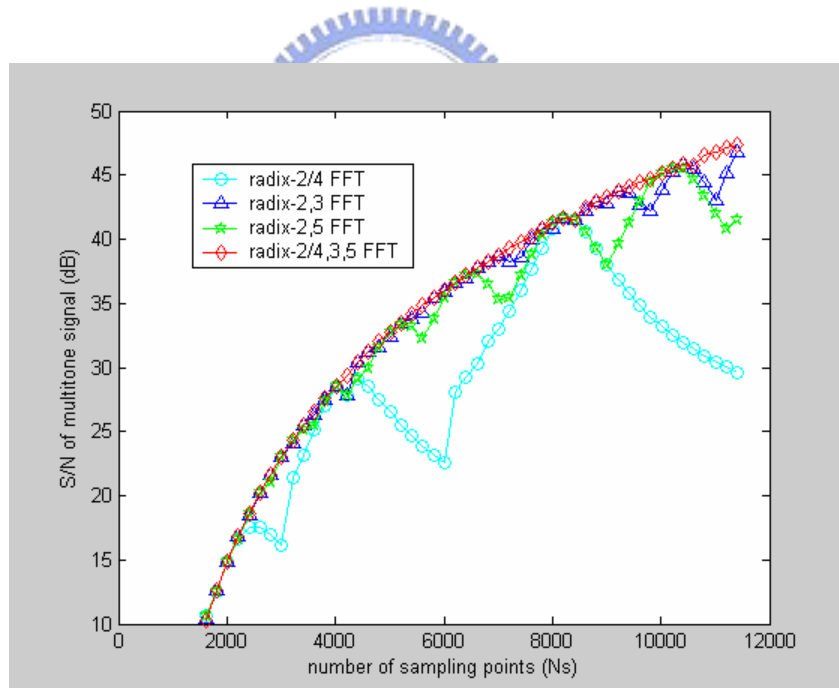


**Figure 4.18: The comparisons of magnitudes of noise of different FFT algorithms (a) with DFT value as a reference, and (b) without DFT value.**

If we express the above in terms of S/N of the signal with respect to the noise induced by the interpolation process, the difference between different algorithms can be more clearly seen. Figure 4.19 plots the S/N of each tone and Figure 4.20 plots the S/N for the combined waveform of tones. It is seen that the radix 2/4,3,5 algorithm has the largest S/N.



**Figure 4.19: The comparisons of S/N's of each tone of different FFT algorithms  
(a) 1<sup>st</sup> tone, (b) 2<sup>nd</sup> tone, (c) 3<sup>rd</sup> tone, and (d) 4<sup>th</sup> tone.**

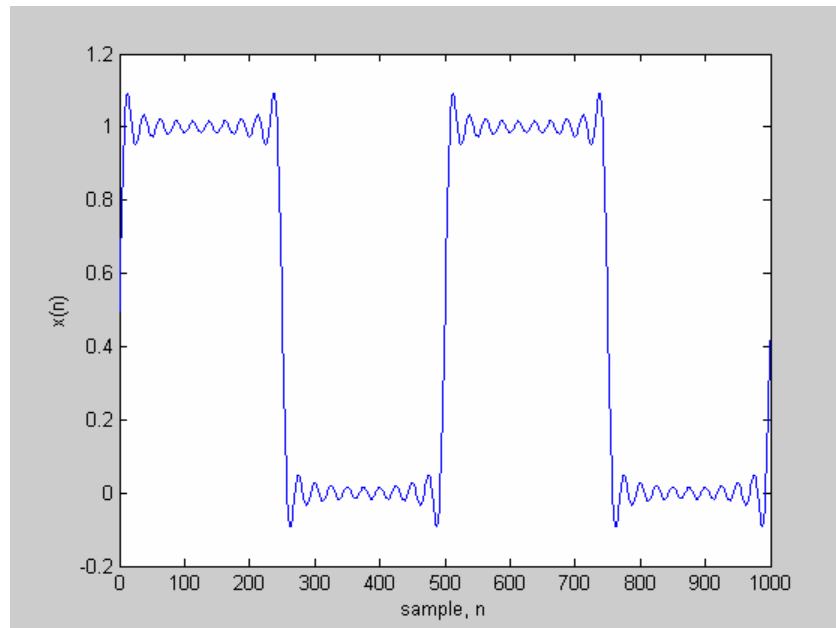


**Figure 4.20: The comparison of S/N's of signal of different FFT algorithms.**

#### 4.4.2 Square Wave

A square wave signal is a signal composed of infinite number of frequencies and is a good signal for making the experiment of applying the developed algorithms. Here we use ten tones

to approximate the square wave. Figure 4.21 shows the magnitude of each tone, the DC value, and the sampling frequency for the signal analyzed.

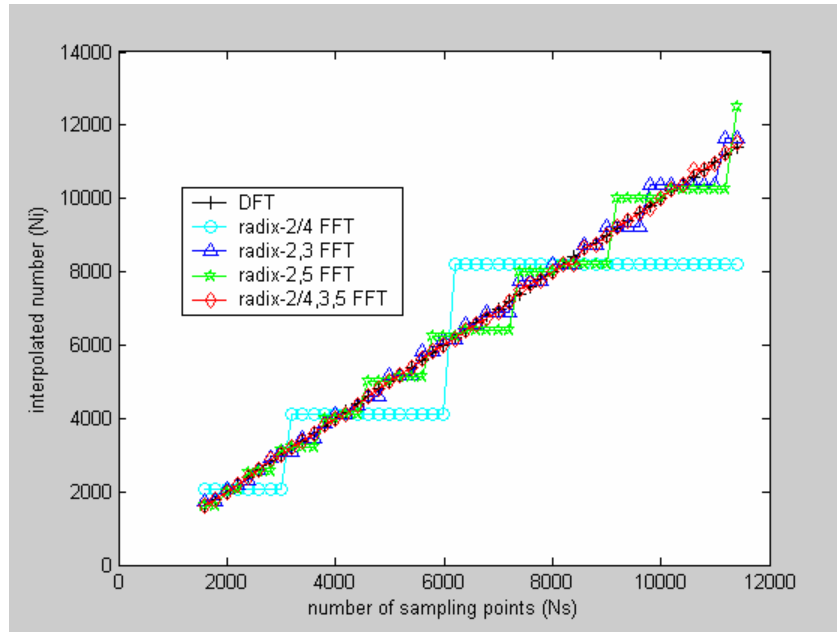


**Figure 4.21: The waveform of the square wave.**

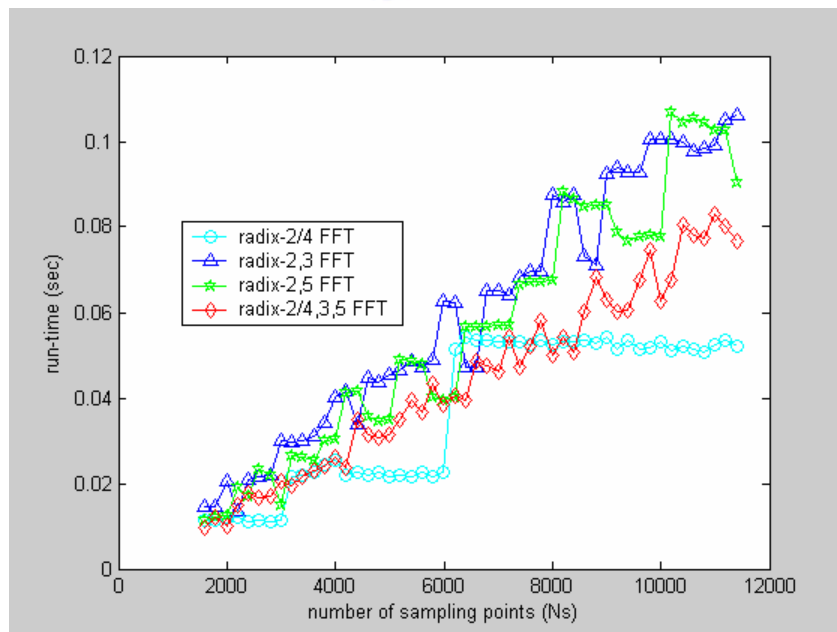
```
f=[1 3 5 7 9 11 13 15 17 19]*42;
m=[0.637 0.212 0.127 0.091 0.071 0.058 0.049 0.043 0.038 0.034];
DC=0.5;
nyquist=max(f)*2;
for (i=1:50)
    fs(i)=nyquist+1+(i-1)*200;
end
```

Figures 4.22, 23, 24, 25, and 26 are the corresponding respective figures of this example as the Figures 4.14, 15, 17, 18, and 20 of the previous example. We can see the similar phenomena as those observed in the Figures 4.14, 15, 17, 18, and 20. The only difference is that since for the square wave, its waveform is mainly determined by the first several fundamental and harmonics frequency components, the sampling frequency selected for this wave is the Nyquist frequency of the last tone of this square wave. This makes the sampling

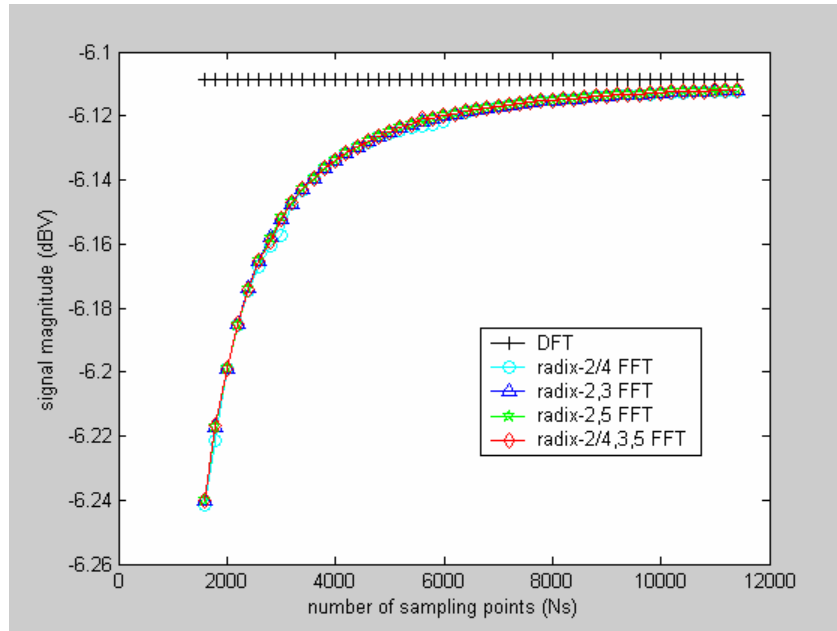
frequency far above than the fundamental tone of this wave and the sampling points are much more than the case of the previous example. This consequently makes the noise introduced due to interpolation less.



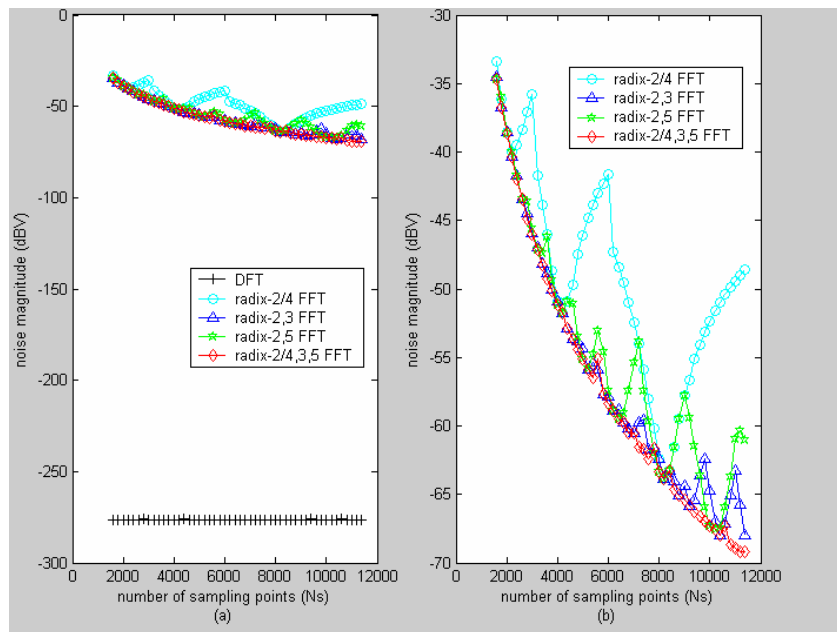
**Figure 4.22: The comparison of applicable number ( $N_i$ ) of sampling points which can be applied by the FFT's.**



**Figure 4.23: The comparison of run-times of different FFT algorithms.**

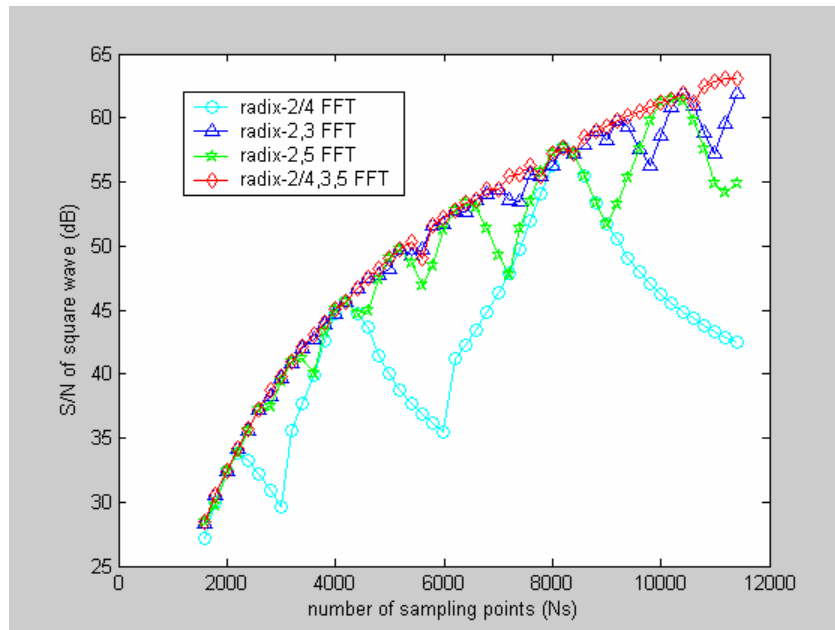


**Figure 4.24: The comparison of magnitudes of signal of different FFT algorithms.**



**Figure 4.25: The comparisons of magnitudes of noise of different FFT algorithms**

**(a) with DFT value as a reference, and (b) without DFT value.**



**Figure 4.26: The comparison of S/N's of signal of different FFT algorithms.**

In summary, from the experiment results on these two examples, the radix-2/4,3,5 FFT exhibits the best performance on the run-time and the S/N over the radix-2,3 FFT and the radix-2,5 FFT algorithms. For the radix-2/4 algorithm, it has a good performance on the run time due to its less number of sampling points, however, it produces too much noise when interpolation is used.

## Chapter 5 Conclusion

In this thesis, we have proposed and studied the mixed-radix FFT algorithm to speed up the conventional radix-2 FFT algorithm. In the work, we first studied the single-radix FFT and derived the decimation-in-time algorithm, the process of re-ordering the data sequence, and discussed the applicable numbers of sampling points which can be applied by the FFT algorithm. Then based on this, we further studied the mixed-radix FFT algorithm. We first derived the algorithm with two radices, i.e., radix- $A/B$  FFT, by directly applying the decimation-in-time algorithm and deriving a mathematical relationship for the decimated order. And then, we derived the radix- $A/B/C$  FFT which is much faster than the conventional radix-2 FFT algorithm. An additional advantage of the algorithm is that the numbers of sampling points which can be applied by the FFT algorithm have been increased a lot.

To further extend the use of the mixed-radix FFT, we adopt the interpolation technique to re-sample the original signal to make the number of sampling points to be applicable by the mixed-radix algorithm. We studied interpolation for the radix-2,3 FFT, the radix-2/4 FFT, the radix-2,5 FFT and the radix-2,3,5 FFT and found that the radix-2,3,5 FFT cooperates with interpolation has the best performance. Furthermore, we extended it to the radix-2/4,3,5 FFT to further improve the computation speed of the algorithm.



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民國八十八年九月至民國九十二年一月

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論文題目： 應用於混模電路測試之利用混合基底與內插法的快速傅立葉轉換

FFT Based on Mixed-Radix and Interpolation for Mixed-Signal Testing