# 國立交通大學

電控工程研究所

#### 碩士論文

# 應用於車用電子之電力線通訊收發器設計 Transceiver Design of Power Line Communication for Automotive Electronics

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中華民國九十九年十一月

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#### 摘 要

近年來車用電子(automotive electronics)越來越受到重視,許多研究或產品多著重在 多媒體以及娛樂等的發展。因此,在汽車自動化的程度越來越高以及車內電子控制單元 也越來越多的情況下,在接線上也就越來越複雜,所以我們將焦點轉移到車內元件彼此 之間的資料傳輸。我們將電力線通信(power line communication)與車用電子結合在一 起,每個裝置都會需要一條線連到車上的電源,這個是沒辦法省略的,所以要是我們可 以同時在電源線上面供給電源又同時傳輸資料的話,便可使接線簡單化,既可降低複雜 度又能省成本。我們利用電源線當作傳輸訊號的通道,為了達到抗拒外界雜訊的干擾, 我們利用了直接序列展頻技術(direct sequence spread spectrum)來對資料做編碼,使得訊 噪比可以提升。我們也設計了一個應用於此技術的收發器,在資料與時脈恢復器上,將 相位選擇與應用在 DSSS 的傳統架構結合在一起。我們也設計了一個具有封包的 32 位 元的串列解串器,使得我們只需要一條電源線就可達到資料的傳輸。整體系統包含了傳 送端、接收端、耦合電路、電源線以及後端的控制電路,控制電路負責的是接收到資料 後產生訊號來控制 LED、馬達以及降壓電路的亮度、閃爍頻率、閃爍的亮度、轉速以及 輸出電壓。我們利用 FPGA 來實現 TX、RX 以及後端的控制電路,然後將 PowerMOS、 LED、馬達、降壓電路、開關等利用印刷電路板來整合,展示了資料傳輸及控制元件的 能力,證明了此電路設計是能夠被實現出來且非常實際的,以及將電力線通信用在車用 電子上的想法是可行的。

關鍵字: 電力線通信(power line communication)、車用電子(automotive electronics)、直接序列展頻技術(direct sequence spread spectrum)

# Transceiver Design of Power Line Communication for Automotive Electronics

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#### **Abstract**

Willey,

In recent years, the importance and diversity of the automotive electronics has attracted more and more attention. Most of the applications are centered on the multimedia or entertainment. Here we focus on the data communication between the electronic devices on the car. We apply the power line communication to the automotive electronics. DC 12V power line of the car is used as the transmission media. We use direct sequence spread spectrum (DSSS) modulation to the data signal in order to resist the noise interference. We design a clock and data recovery (CDR) circuit which combines the phase picking for clock recovery and byte synchronization for data recovery. We realize a transceiver with DSSS and a 32-bit SERDES. The whole system includes a transmitter, a receiver, a coupling circuit and a control circuit. A delta-sigma digital pulse width modulator ( $\Delta\Sigma$ -DPWM) driver is implemented for the power driving. The control circuit adjusts PWM for on-time duty and frequency control of the LEDs, controls the rotation speed of the motors and the output voltage of DC-DC converter. We use FPGA to realize the transmitter and receiver which includes the control circuit. We put PowerMOS, LED, motor, buck converter, switches and coupling circuits on the printed circuit boards. Finally, we demonstrate the data transmission and controlling of the devices. We prove the idea of power line communication for automotive electronics is feasible.

Keyword: power line communication, automotive electronics, direct sequence spread spectrum (DSSS)

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首先要感謝的是我的指導教授蘇朝琴老師,總是告訴我們要不斷的去思考,當問題發生時不要去逃避,表面上看到的東西不一定就是真相,要仔細地去思考背後的一切。做研究最重要的就是要培養我們的思考能力以及解決問題的能力,模擬只是輔助,最重要的是自己的想法。不管是在學習上或是生活中,在老師的教導下,都使我獲益良多。

此外感謝實驗室的所有學長、同學、學弟妹的幫忙與支持。謝謝丸子學長在我一 開始很灰心無助的時候開導了我,適時的點醒了我,告訴了我做研究最重要的事情,讓 我能重新振作起來。謝謝盈杰學長的鼓勵以及協助。多虧了有你的帶領,實驗室的任何 大小事都能完善的解決。謝謝仁乾學長對於 PLL 方面給了我許多的指導以及建議,也 很謝謝你能一直關心我們這些學弟。謝謝庭佑學長對於研究上面的協助,以及很辛苦的 維護實驗室的工作站及硬體設備、使我們能擁有穩定的模擬環境。謝謝洲銘、家齊學長 陪伴了我們快兩年的時間,在研究上一起討論,在生活上一起分享。很懷念晚上一起在 實驗室的日子。謝謝于昇學長在學校陪伴了我們兩年多的日子,但我們還是室友所以生 活上還是不會分開的。在未來還是能一起討論學術,一起去看表演,一起練吉他,一起 看曼聯輸給切爾西吧!謝謝土豆的一起努力,想當初也是因為我人太好才會想跟你一 組,沒想到在課業上也都跟你同一組。想到那時候我們辛苦的做研究以及量測,深怕會 做不出來的時候,到現在能夠跟你一起口試一起畢業,真的是太感動了,以後也要一起 去打球吃飯跟大潤發。也很謝謝海豹、修銘、群育、博祥,在這兩年多以來大家一起討 論功課一起做研究,雖然彼此之間做的東西是不一樣的,但都能互相給意見,在生活上 也增添了許多快樂。謝謝澤勝、璟伊、弘宇、順煜、昶志、馬克、阿 MON、小紅豆、 嘉哲學弟妹們的陪伴,在枯燥的研究生涯裡增添了許多的歡樂色彩。希望大家都能夠研 究順利以及找到自己想做的事。

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# Chapter 1

#### Introduction



#### 1.1 Motivation

As the car technology advanced rapidly, more and more electronic devices are applied on the car, such as LED lighting, electric sunroof, electric windows, etc. are simple ones. Multimedia center, GPS navigator, ABS system, air conditioning thermostat, etc. are complex ones. The cost of electronics in a regular car is about 15% to 25%. For a luxury one or RV, it can be over 50%, therefore, the automotive electronics is becoming increasingly popular and important.

Observing from the automotive electronics these days, the devices on the cars have independent communication agreement, such as CAN (Control Area Network), LIN (Local Interconnect Network), or FlexRay. It means that each device needs a dedicated

communication bus between control centers. Different agreements are used for different local control or different cars. As the device increases, the number of wires increases. The complexity of wiring and routing becomes more and more difficult. In order to achieve the control of devices and decrease the complexity of wiring and routing, we take advantage of the existing power line on the car to achieve the data communication.

Power line communication becomes attractive for transmission of power and data with no extra wiring [2,3,4]. The research in power line communication has been a long time. At the beginning, they are used for measurement of loading and automation on the high voltage. Now the communication network on the medium voltage and low voltage is demanding and popular.

But the difficulty of power line communication is that the signal is interfered by various noises such as impulse noise or continuous wave noise which is mixed of wideband and narrowband. The amplitude sometimes is as large as signal and the channel characteristic is time-variant so it is difficult to modeling.

#### 1.2 Features

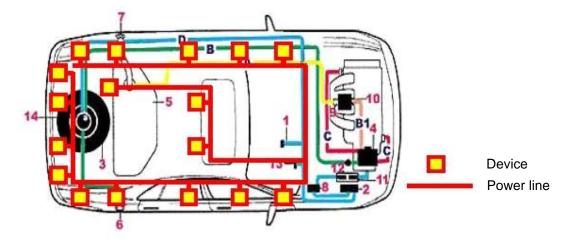


Fig. 1-1: A generalized model of a power line communication for automotive

In this thesis, we propose a method of power line communication for automobile combining a programmable power driving circuit. We use the DC 12V power line existing on

the car to supply power and transmit control data to the electric loads such as LED, motor, DC-DC converter. We use capacitive coupling to transmit and receive data and use inductive coupling to supply power. For the noise interference, we use direct sequence spread spectrum technique to modulate the control data, and we design a transceiver to realize this function.

#### 1.3 Organization

This thesis comprises six chapters. The motivation and features are described in this chapter. Chapter 2 describes the background studies. It starts with an overview of the basic serial link and other basic knowledge. Then, we introduce the architecture of the clock and data recovery (CDR). Finally, we introduce the direct sequence spread spectrum and its method for data recovery.

Chapter 3 addresses the idea of our proposed CDR and analyzes the system-level behaviors of the CDR. We built the transmitter and the receiver for the system behavioral simulation and discuss the effect of jitter and frequency error. Chapter 4 depicts the implementation of sub-circuits of the transmitter and the receiver. We use SPICE to simulate and verify that it can function well.

Chapter 5 shows the realization on the FPGA and PCB. We use FPGA to realize the transmitter and receiver and backend control circuit. The LEDs, motors, buck converters and other front-end discrete components are realized on the PCB. We connect the system to the power line and demonstrate the controlling of the LEDs, motors and buck converters. Finally, chapter 6 concludes this thesis.

## Chapter 2

# **Background Studies**



#### 2.1 Basic Serial Link [5, 6]

The concept of serial communication is the process of sending data one bit at a time, sequentially, over a communication channel or computer bus. This is in contrast to parallel communication which several bits are sent at a time over several parallel channels or buses. A generalized model of serial communication is illustrated in Fig. 2-1. It consists of a transmitter, a communication channel, and a receiver. The transmitter end contains a serializer, a phase-locked loop (PLL), and an output driver. The PLL supplies the multi-phase clock to the serializer which serializes the parallel data bits into one serial sequence. The output driver will change the shape of data according to the channel and output the data. The receiver end contains an input amplifier and a clock recovery circuit and a deserializer. The input amplifier receives the data and restores it. The clock recovery circuit extracts the clock from the input

data and retimes the input data. The data is descrialized into parallel form by a descrializer.

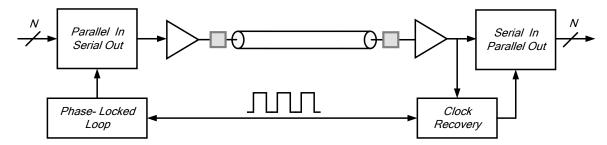


Fig. 2-1: A generalized model of a serial link

#### 2.1.1. Jitter Analysis

Jitter is the time domain behavior of phase noise. It represents the deviation of the zero crossing from their ideal position. In Fig. 2-2,  $x_1(t)$  represents the ideal signal and  $x_2(t)$  represents the actual signal. To quantify jitter we can measure the deviation of each positive or negative transition point of  $x_2(t)$  from its corresponding point in  $x_1(t)$ , i.e.,  $\Delta T_1$ ,  $\Delta T_2$ , etc.

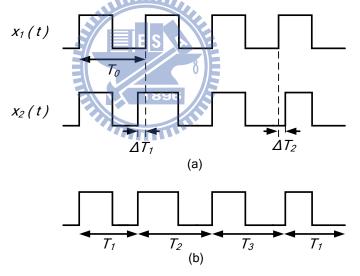


Fig. 2-2: (a) Cycle-to-cycle jitter, (b) variable cycles.

This type of jitter is called absolute jitter. Since the deviations are random, we measure a very large number of  $\Delta T$ 's and determine the root mean square (RMS) value of absolute jitter as:

$$\Delta T_{abs,ms} = \lim_{N \to \infty} \frac{1}{N} \sqrt{\Delta T_1^2 + \Delta T_2^2 + \dots + \Delta T_N^2}$$
 (2.1)

Another type of jitter is called cycle-to-cycle jitter. It is obtained by measuring the difference between every two consecutive cycles of the waveform and taking the RMS value:

$$\Delta T_{cc,rms} \approx \lim_{N \to \infty} \frac{1}{N} \sqrt{T_2 - T_1^2 + T_3 - T_2^2 + \dots + T_N - T_{N-1}^2}$$
 (2.2)

Absolute and cycle-to-cycle jitters are commonly used to characterize the quality of the signals in the time domain. A third type of jitter is called period jitter. It is defined as the deviation of each cycle from the average period of the waveform,  $\overline{T}$ :

$$T_{p,ms} = \lim_{N \to \infty} \frac{1}{N} \sqrt{\overline{T} - T_1^2 + \overline{T} - T_2^2 + \dots + \overline{T} - T_N^2}$$
 (2.3)

Generally speaking, the total jitter has two components: deterministic jitter (DJ) and random jitter (RJ), as shown in Fig. 2-3.

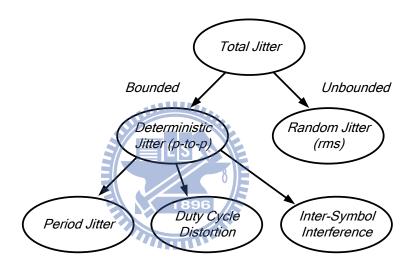


Fig. 2-3: Components of total jitter

The probability density function (PDF) of deterministic jitter is non-Gaussian. So it has bounded peak-to-peak value. There are three types of DJ: periodic jitter, inter symbol interference and duty cycle distortion. The PDF of random jitter is Gaussian. So it is unbounded. It is often generated by different noise sources such as thermal noise, power supply noise, substrate noise, etc. Random jitter is characterized by RMS value because it takes a long time to obtain the peak-to-peak value to achieve statistical significance.

Multiple random jitter sources can be added in an RMS fashion. But it needs a peak-to-peak value to be added to the deterministic jitter to obtain total jitter. Although the Gaussian statistics imply infinite peak-to-peak value, we can use the probability of data error

to transform RJ to DJ. For example, if the specification of error is 10<sup>-12</sup>, then the RMS value of RJ multiplies 14.1 is the peak-to-peak value of RJ. The conversion between RMS and peak-to-peak value of RJ are illustrated in table 2.1

Table 2.1

Probability of Data Error	N= Peak-to-Peak /RMS
10 <sup>-12</sup>	14.1
10 <sup>-13</sup>	14.7
10 <sup>-14</sup>	15.3
10 <sup>-15</sup>	15.9

#### 2.1.2. Jitter Transfer

Jitter transfer function of a device represents the output jitter as the input jitter is varied at different rates. It is defined as the ratio of jitter on the output to the jitter applied on the input of the device under different frequency. It can quantify the jitter accumulation performance of data retiming devices. Fig. 2-4 shows the jitter transfer mask for OC-192.

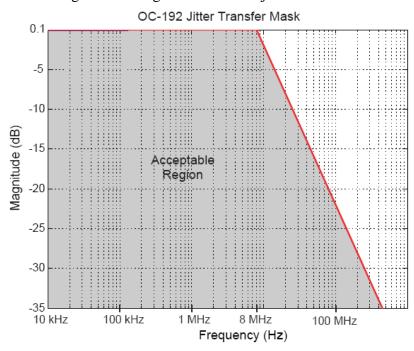


Fig. 2-4: Jitter transfer mask[1]

#### 2.1.3. Jitter Tolerance

Jitter tolerance specifies how much input jitter a device must tolerate without increasing the bit error rate. For example, the Fig. 2-5 shows the mask of jitter tolerance for OC-192. It indicates that the system must retime the input data correctly if the input data jitter is over 15UI when the jitter frequency is less than 2.4 KHz.

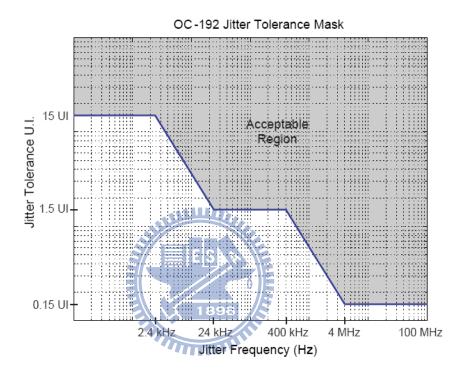


Fig. 2-5: Jitter tolerance mask[1]

#### 2.1.4. Bit Error Rate (BER)

Bit error rate is an index for the system performance. It represents the reliability of the link and it defines the data rate of transportation. If the BER is rising above a specified level, it means that the system could not operate properly at the rate. It is calculated as follow:

$$BER = \frac{\text{number of erroneous bits}}{\text{number of transmitted bits}}$$
 (2.4)

#### 2.2 Architecture of CDR [7]

Fig. 2-6 shows the concept of a clock and data recovery (CDR) circuit. At the receiver side the received data is asynchronous and noisy, so a CDR circuit is used to extract the information of clock from the input data and synchronize the input data.

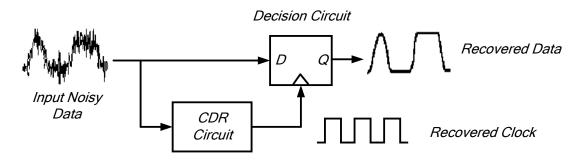


Fig. 2-6: Concept of a CDR circuit

#### 2.2.1. PLL-Based CDR

PLL-based CDR is a closed loop system. In general, it refers to and Nth-order system, where  $N \ge 2$ . A general function block of PLL-based type clock and data recovery circuit is shown in Fig. 2-7[8]. It consists of phase and/or frequency detectors, a charge pump, a loop filter,

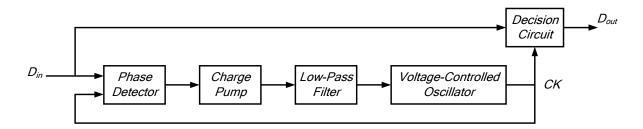


Fig. 2-7: A general function block of PLL-based type CDR [8]

and a voltage-controlled oscillator (VCO). The frequency detectors (FD) are used for the pull-in process and the phase detectors (PD) are used for the lock-in process. According to the PD structure used in a PLL-based CDR, we can categorize it into linear and binary types [9]. The linear PD proposed by Hogge outputs an Up pulse when the rising edge of the clock leads the transition edge of the data. When the rising edge of the clock lags the transition edge of

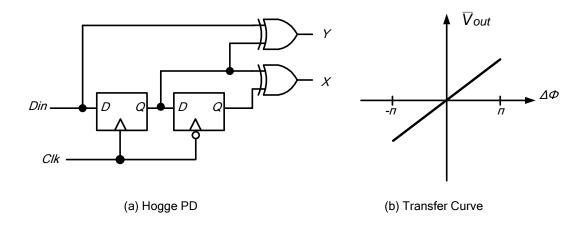


Fig. 2-8: Linear PD proposed by Hogge

the data it outputs a Down pulse. The width of the pulse is the phase difference between the data and the clock. The phase difference is converted into control voltage of the VCO. The high frequency noise in this control voltage is filtered out by the loop filter. The loop adjusts the frequency of the VCO by the control voltage until the phase is exactly the same as the input data. The VCO will provide the sampling phase to the decision circuit to retime the input data.

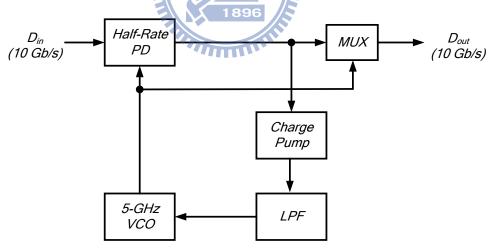


Fig. 2-9: A PLL-based CDR with linear PD [8]

Fig. 2-8 shows an example of a CDR with linear PD proposed by Jafar Savoj. It incorporates a half-rate PD to retime and demultiplex the input data. The half-rate PD produces an error signal and a reference signal. The error signal represents the phase difference. But it is effective only with the reference signal. Because the random nature of the

data and the periodicity of the clock make the average of the error signal which depends on the pattern of the input data. So a reference signal must also be generated to reduce the dependence by averaging.

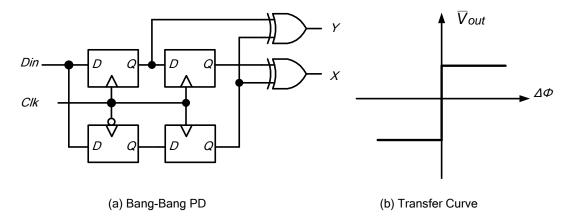


Fig. 2-10: Bang-bang PD proposed by Alexander

The binary PD proposed by Alexander is also called a bang-bang PD. As we can see from the transfer curve, its output pulse only tells that the clock is lead or lag no matter what the phase difference amount is. The bang-bang PD is usually used in the all digital structure with digital loop filter or confidence counter and a digitally controlled oscillator. Fig. 2-11 shows an example of a CDR with binary PD proposed by Jafar Savoj.

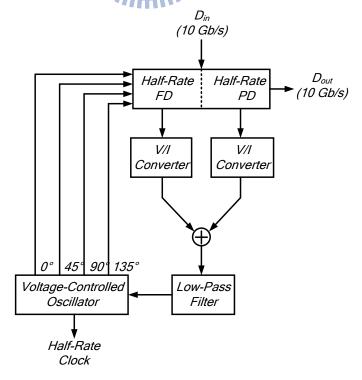


Fig. 2-11: A PLL-based CDR with binary PD [10]

The system utilizes a half-rate binary PD and a half-rate binary FD to indicate both the phase and frequency difference. Based on the half-rate topology, it reliefs the requirement on the circuit and retimes the input data inherently. The FD generates an error signal and controls the VCO frequency toward half of the input data. Then the PD takes over. It locks the VCO phase to the input data and produces a retimed output.

The stability of the feedback system is another issue for the PLL-based CDR. The tracking bandwidth is limited by the stability of the system and we also have to meet the jitter specification such as jitter tolerance, jitter transfer, etc. In general the loop bandwidth is usually less than one-tenth of the internal clock frequency.

#### 2.2.2. Oversampling-Based CDR

A blind oversampling CDR is a feed-forward system. Basically, it uses multi-phase of the clocks generated from PLL to sample the input data. Usually a single data bit will be sampled three times. It is called 3x-oversampling. We can increase the samples to four or higher to eliminate the transition ambiguity. After oversampling, the transition information of

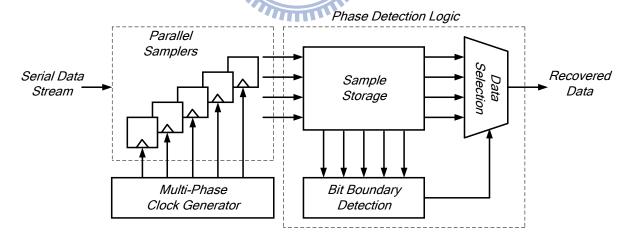


Fig. 2-12: A block diagram of a blind oversampling CDR [11]

the input data must be detected. We can employ the center-picking or majority-voting method to extract the transition edge of the input data and we can pick the optimum sampling phase of the clock. So it is also called the phase picking.

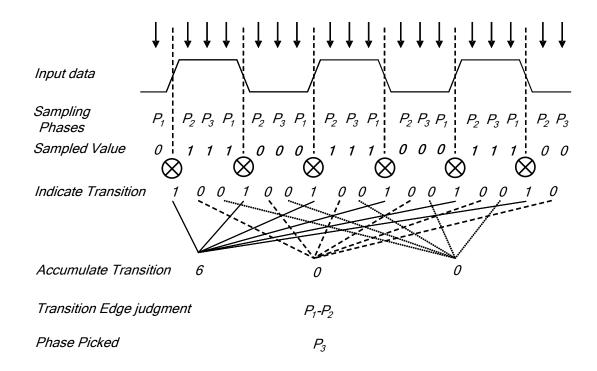


Fig. 2-13: An example of phase decision

For example, the data sampled by 3 phases per bit time is shown in Fig 2-13. It decides the boundary with a portion of a sampled stream. First the neighboring data are XORed. Then the transitions are accumulated. The transition that has the highest accumulated count is the data transition edge. Then we can choice P3 as sampling phase. We can increase the length of the sampled stream to obtain more accurate results but the storage hardware overhead will also be increased.

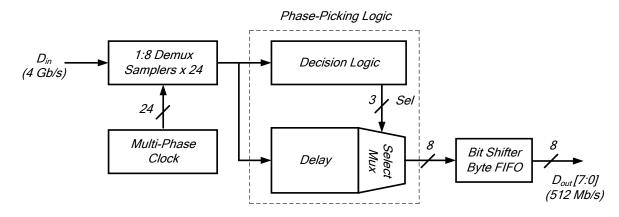


Fig. 2-14: An example of oversampling CDR [12]

The phase picking scheme accompanies static phase offset error on each sampling because neither the data nor the clock phases are adjusted. The maximum phase offset is (0.5UI/OSR), where OSR denotes the oversampling ratio. Although we can increase the OSR to decrease the phase offset but in practical a higher OSR implies high accuracy phase resolution for each sampling, which is always a challenge. A phase-picking CDR tracks the high-frequency jitter of the input data well while it needs extra circuit such as FIFO buffer to tracks the frequency error between the input data and the clock. This architecture eliminates the need on the acquisition time but requires more hardware for executing algorithm and introduces processing latency to the data recovery.

#### 2.2.3. DLL-Based CDR

DLL-based CDR is just like PLL-based CDR but it adjusts the phase of the input data or the clock. It is a closed loop first-order system. So it is a simplified version of PLL-based architecture. According to the subject of the delay adjustment, it can be divided into two types: clock interpolation and data-deskew. The clock interpolation architecture can track the input data continuously by the phase rotation scheme. But it needs additional hardware such as FIFO buffer to overcome the data overflow/underflow problem. For the data-deskew architecture, it adjust the phase of the input data and synchronize it with the phases of the clock. It is limited by the tuning range of the circuit. So it is only suitable for burst-mode application.

#### ■ Clock Interpolation

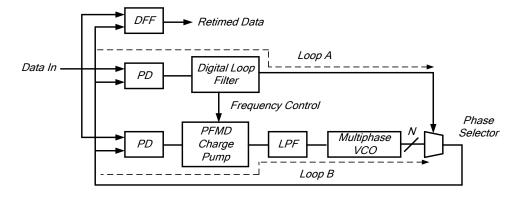


Fig. 2-15: A clock interpolation CDR proposed by Ruiyuan Zhang [13]

Fig. 2-15 shows an example of the clock interpolation CDR proposed by Ruiyuan Zhang. It combines the fast acquisition of a phase selection DLL with the low jitter of a PLL. Loop A is the phase selection loop. It provides fast acquisition with the input data and frequency direction and magnitude estimation for the PLL loop. Loop B, a PLL loop with narrow loop bandwidth to achieve low jitter. The phase switching is disabled after the PLL acquires the phase lock. The phase selector examines the relationship between the data transitions and the clock phases from the VCO, and then the circuit selects the sampling phase which is the farthest from the data transitions.

#### ■ Data-Deskew

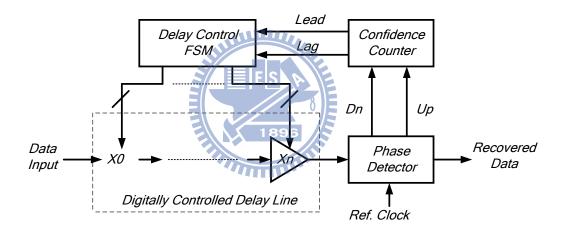


Fig. 2-16: A data-deskew CDR proposed by Hung-Wen Lu [14]

Fig. 2-16 shows an example of the data-deskew CDR proposed by Hung-Wen Lu, the bang-bang PD outputs up and down pulses to give the information of leading and lagging between the clock and the input data. The confidence counter serves as a digital loop filter accumulates the pulses and filter out the high frequency noise. Then it controls the delay through a finite-state machine (FSM). Finally the input data passes through the delay line and synchronizes to the clock. This paper proposes an all digital architecture, so it can provide high flexibility and can be synthesizable.

## 2.3 Direct Sequence Spread Spectrum

#### 2.3.1 Definition of Spread Spectrum

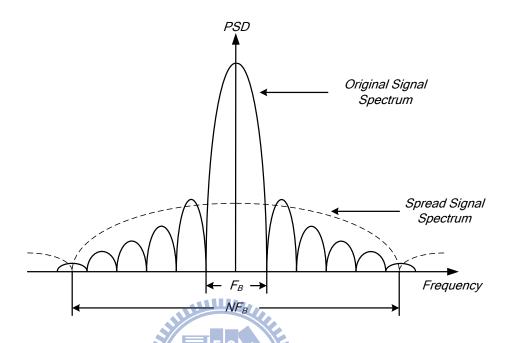


Fig. 2-17: Power spectral density of spread spectrum signal [15]

Spread spectrum is a technique in which a signal is transmitted on a bandwidth larger than the frequency content of the original information by modulating a signal. It generally makes use of the noise-like signal to spread the normally narrowband information signal over a relatively wideband. The receiver correlates the received signals to retrieve the original information signal. The spread spectrum has a lot of benefits, such as: anti-jamming, anti-interference, low probability of intercept, multiple user random access communication, etc. There are two main techniques of the spread spectrum communication: direct sequence and frequency hopping. Here we use the direct sequence spread spectrum. We will illustrate it in the following paragraph.

#### 2.3.2 Direct Sequence Spread Spectrum

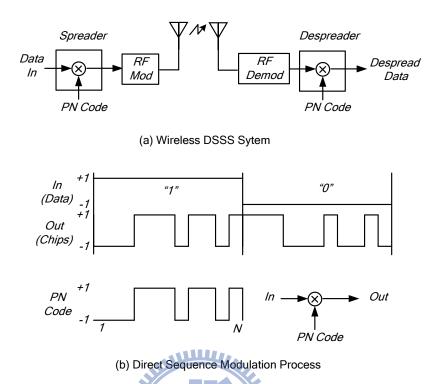


Fig. 2-18: DSSS system and modulation process [16]

Direct sequence spread spectrum (DSSS) transmission multiply the data being transmitted by a noise signal [17]. This noise signal is a pseudorandom sequence of 1 and -1 values at a frequency much higher than that of the original signal. Thereby, it spreads the energy of the original signal into a wider band. The pseudorandom sequence or pseudo noise (PN) code symbols are called chips. This PN code is generated in a deterministic way but act like a random signal. Usually it is generated by a linear feedback shift register (LFSR). The transmitted signal resembles white noise. The noise-like signal can be used to reconstruct the original data at the receiver end. The despread process multiplies it by the same PN code. It calculates the correlation between the transmitted sequence and the local PN code. In a binary direct-sequence system, a chip is a pulse of a DSSS code. Each chip is typically a rectangular pulse of 1 or -1. It is called chips to avoid confusing them with information bits. The chip rate is larger than the data rate. That is, one date bit is represents by multiple chips. The ratio is known as spreading factor or processing gain:

$$PG = \frac{\text{chip rate}}{\text{data rate}}$$

$$= \frac{\text{bandwidth of spread spectrum signal}}{\text{bandwidth of the information signal}}$$
(2.5)

The processing gain also results in enhancing the signal to noise ratio (SNR) on the channel. When the spreaded signal is transmitted in the channel, it is interfered by narrow band noise. At the receiver end, we use the same PN code to despread the data. At the same time, the noise will be spreaded and its power is being reduced. So we can attain a better signal to noise ratio.

#### 2.3.3 Code Synchronization

For despread to work correctly, the transmitted and received sequence must be synchronized. It is called code synchronization. It is a process of aligning the PN code generated from the receiver with the transmitted sequence.

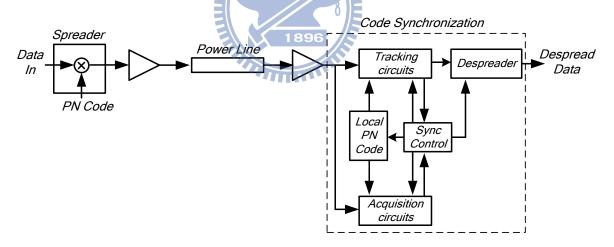


Fig. 2-19: Code Synchronization [18]

Code synchronization is composed of two parts: code acquisition for the coarse synchronization and code tracking for the fine synchronization [21]. The acquisition process involves a search through the region of time-frequency uncertainty and determinate that the locally generated PN code and the received sequence are sufficiently aligned within 0.5 chip or less. Code tracking is the process of maintaining the alignment of the two signals. This is

done by using a feedback loop which constantly serves to reduce the phase error and retain the alignment. An essential function for timing acquisition, tracking and data recovery is to calculate the correlation between the local PN code and received sequence. Usually we use matched filter or correlator to calculate the correlation function of the received signal and local PN code.

#### ■ Code Acquisition

Typically, the methods for PN acquisition can be classified into serial search and parallel search.

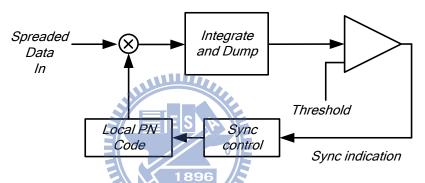


Fig. 2-20: Serial search with correlator [19]

As shown in the Fig. 2-20, this is the serial search with correlator. First, the timing of the local PN code has been set and it is correlated with the transmitted sequence. The integrate and dump block calculates the correlation value over k chips time and dump out the value. The comparator compares the value with a threshold. If the threshold is not exceeded, the sync control will delay the local PN generator usually by 1/2 chip time. The search process is started again. When the transmitted sequence and local PN code are roughly aligned, the delay mechanism will stop and the tracking will be initiated.

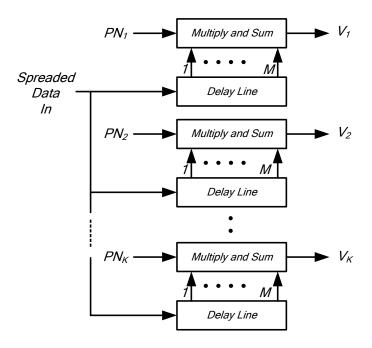


Fig. 2-21: Parallel search with matched filter [19]

Another one is the parallel search with matched filter. It uses a bank of correlators or matched filters. Each one has the same PN code with equally spaced timing. When correlating with the transmitted sequence, one of the correlators will exceed the threshold and indicate the synchronization of both signals. As compared with the serial search method, the parallel search can provide fast acquisition but more hardware is required. Therefore, there is a time-complexity tradeoff.

#### **■** Code Tracking

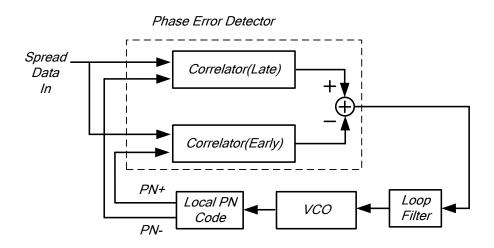


Fig. 2-22: Delay-locked loop [20]

After acquisition or coarse synchronization, code tracking or fine synchronization takes place. There are two well-known methods: delay-locked loop and tau-dither loop. Fig. 2-22 shows the blocks of the delay-locked loop. It uses two correlators and two PN sequences which they are the same sequence with different phase delay. The early phase is typically 1/2 chip time earlier with respect to the transmitted sequence. The other one is the late phase which is 1/2 chip time delay. Using these two sequences to correlate with the transmitted sequence and subtract the results. We can get the information of phase error between the transmitted sequence and local PN code. The loop filter is used to filter out the high frequency noise on the error signal and control the VCO to increase or decrease the frequency smoothly. Therefore, the early and late PN sequence will track the transmitted sequence and let the phase error approach zero.

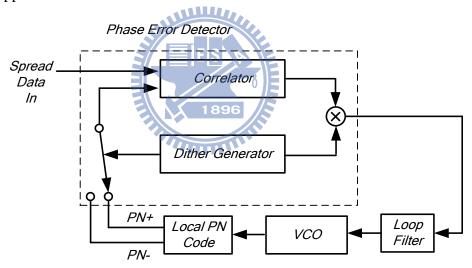


Fig. 2-23: Tau-dither loop [20]

The second one is the tau-dither loop. It is the delay-locked loop with only one correlator. The dither generator controls the early and late sequence and let one of them correlated with the transmitted sequence at each instant of time and then correlates with the result of correlator to get the phase error signal. The next steps are just like the delay locked loops.

## Chapter 3

# **Behavioral Simulation of the Proposed CDR**



# 3.1 The Proposed CDR

According to the application, the CDR is included in a transceiver of power line communication which processes the data at several MHz. So, if we choose a PLL-based CDR, we need a low speed VCO at the same as the reference clock. The resolution issue and lock-in time issue makes the PLL-based CDR not suitable for our system. The data-deskew CDR has the same problems. With oversampling-based CDR, we just need a multi-phase clock generator. The phase decision logic can be implemented as an all-digital structure. It is the most robust and simplest way to design our system. But the oversampling-based CDR needs a large storage component to eliminate the transition ambiguity or noise to recover the data. So as to save the hardware overhead, we use the phase rotation method to realize the clock and data recovery. Basically, we use three phases of clock to sample the data. The

3x-oversampling can get enough information of data transition edge with the least phase resolution. Because the data rate is just several MHz, using more phases of clock with higher phase resolution to sample the data is impractical and uneconomical. From the sampled data, we can get the lead and lag information between clock and data. We rotate the phases of clock backwards or forwards to get the optimum sampling point of the data. Traditionally, we have to collect the information of data transition of several bits to decide whether the phases should be shifted or not to avoid the effect of jitter or other noises. Because we use the DSSS modulation scheme, the data has been split into 11 chips. We have to calculate the correlation value between the chips and the local Barker code. We can get enough information from 11 chips. We can use the correlation value to decide whether we shift the phase of clock. For the code acquisition and code tracking, we use the same phase rotation scheme to achieve that. The chip rate is 11MHz and the data rate is 1MHz for the preliminary design. We generate 3 phases of the clock of 11MHz for oversampling.

#### 3.2 The Block Diagram of Proposed CDR

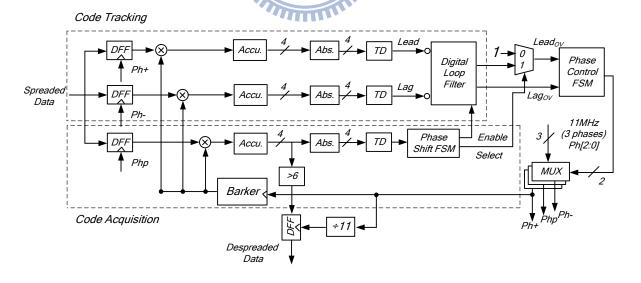


Fig. 3-1: The proposed CDR

Fig. 3-1 shows the block diagram of the proposed CDR. The system can be divided into two parts consists of code acquisition loop and code tracking loop. Firstly, it uses the punctual

phase (Php) of the clock to sample the input data. The punctual phase is one of the three phases of the clock generator. The sampled data are XORed with Barker code and the accumulator sums up the outcome and dumps out every bit time. Because all the calculation is done with unsigned operation, the accumulator just only counts the number of times of one.

```
1: 11 10 9 8 7 6 5 4 3 2 1 0
0: 0 1 2 3 4 5 6 7 8 9 10 11
Transform: 11 9 7 5 3 1 -1 -3 -5 -7 -9 -11

Abs.: 11 9 7 5 3 1 1 3 5 7 9 11
```

Fig. 3-2: Accumulation and absolution results

Fig. 3-2 shows all the case of accumulation results. For the perfect alignment of the sampled data and the local Barker code, the results will be 11 or 0 as for data bit 1 or 0. For one chip misalignment or more, the results will be 5 or 6. So we have to set a threshold value to indicate that whether the sampled data are aligned to the local Barker code. At the same time, we can tolerate that some chips are sampled wrong due to jitter or noise, and still can recover the original data. But at first we must put the threshold value of data bit 1 and 0 at the same benchmark of comparison. As shown in the Fig. 3-2, we transform the 0 to -1, and absolutize the results. So now no matter what the data bit is, we can compare the accumulated results with threshold value fairly. We set the threshold value to 3. The accumulated results must be greater than it to acquire alignment. If the sampled chips do not synchronize to the clock, the select of phase shift FSM will be 0 and let the lead<sub>OV</sub> be 1. It forces phase control FSM to shift the phase of clock in order to proceed next synchronization. In code acquisition, we delay the local Barker code 1/3 chip time when shifting the phase of clock. For the worst case, when the local Barker code is 1 chip time delayed to the sampled chips, it has to shift 33 times of the phase of clock so as to acquire alignment. When the code acquisition is done, the phase shift FSM will enable the code tracking loop and select the output of digital loop filter as the multiplexer (MUX) output. We take the essence of delay locked loops of basic code tracking loop. We employ two phases of clock which are the early (Ph+) and late (Ph-) as compared with the punctual phase (Php). Basically, the blocks in the code tracking are the same in the code acquisition. When there has frequency error between the sampled chips and local clock, the early or late path will fail the threshold test progressively. It implies that the three sampling phases should be shifted to handle the frequency error. In another way, the situation of frequency error is just like the data drifts forward or backward from the local clock since the frequency of both does not match. Therefore, we need to shift the phase of clock in the same way as drifting of the sampled chips. Considering the effect of jitter, we might sample the wrong chips and fail the threshold test and the system would shift the phase of clock. But essentially the system must be capable of dealing with the frequency error and removing the effect of jitter. According to this, a loop filter must be put into the loop in order that the system can smoothly shift the phase of clock due to frequency error and eliminate the interference of jitter. The recovery of data is done by comparing the accumulated result with magnitude threshold value 6 directly. Because we have confirmed that the sampled chips and the local Barker code are synchronized. So now we just need to discriminate the data bit 1 from 0. To sum up, the system performs byte recovery, that is, code acquisition firstly, and then clock recovery, that is, code tracking, and it uses a D flip-flop to retime the despreaded data.

#### 3.2.1. Correlator of the System

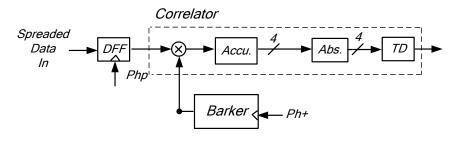


Fig. 3-3: Correlator

The correlator is used to calculate the correlation value between the sampled chips and

the local Barker code. Actually, it has a 4-bit DFF between the accumulator and the absolute value circuit. So the XOR gate, accumulator, and DFF form a digital integrate and dump circuit. It dumps out the accumulated results every bit time, or, every 11 chip time. It goes through the absolute value circuit and the threshold detector. The output of the threshold detector is just 1 or 0 to indicate that whether the accumulation result is greater than the threshold value. If so, the synchronization is completed. Fig. 3-4 shows the behavioral simulation of the correlator.

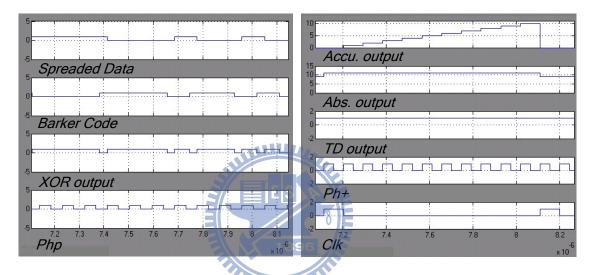


Fig. 3-4: Simulation of correlator

#### 3.2.2. Phase Shift FSM

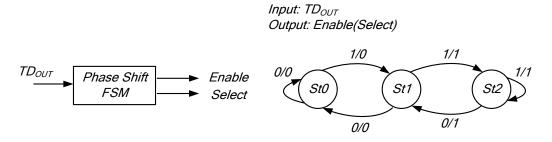


Fig. 3-5: Phase shift FSM

The phase shift FSM in the code acquisition loop controls the phase shifting according to the output of the threshold detector. The state diagram is described that the output of threshold detector has to be 1 two times continuously to set the select and enable signals 1. The local Barker code will synchronize to the sampled chips if it passes the threshold test two times

successively, in case that the noise or jitter causes the system to make wrong decisions. When the burst error happens, it may last more than one bit time. The timing relationship between the samples chips and the clock must be unchanged. So we design a FSM that the misalignment must occurs two times successively in order to restart the synchronization.

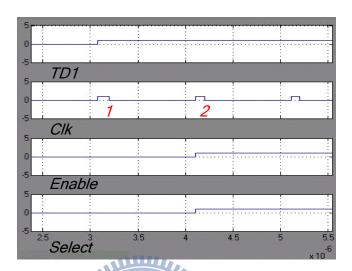
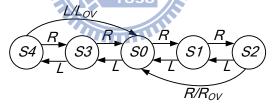
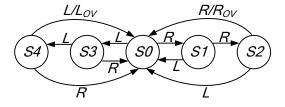


Fig. 3-6: Phase shift FSM

## 3.2.3. Digital Loop Filter



(a) Accumulative type



(b) Continuous type

Fig. 3-7: Types of confidence counter

The position of the data transition will move because of the static phase error (frequency error) or dynamic phase error (jitter). The effect induced by the jitter on the sampled chips will cause the error chips and change the state of the system. We use a confidence counter as

the loop filter. There are two types of confidence counter. One is the accumulative type and another is the continuous type. As shown in Fig. 3-7, the initial state of the confidence counters is set to S0 on both types. R represents the change of the state of the confidence counter to the right; L represents the change of the state of the confidence counter to the left. The R<sub>OV</sub> and L<sub>OV</sub> are the overflow signals indicate that the system should change its state. In the accumulative type, the overflow only happens if three times of R or L are accumulated. In the continuous type, it happens when R or L occurs three times continuously. Considering the situation where the dynamic and static phase error happens to the system. Because we use 3x-oversampling, the phase resolution is very low. The system can tolerate more dynamic phase error. So considering the frequency error only, the phases of clock (Ph+ or Ph-) will move toward the transition edge of data gradually. Therefore, the system shifts the phase of clock if one of the two paths fails the threshold test several times in a continuous way. Consequently, the continuous type confidence counter is more suitable for our system.

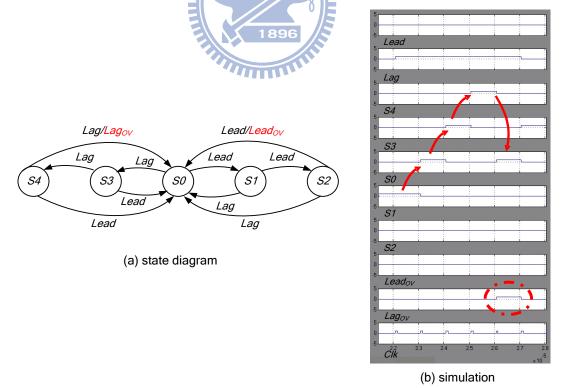


Fig. 3-8: Confidence counter

#### 3.2.4. Phase Control FSM

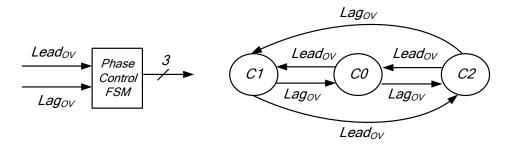


Fig. 3-9: Phase control FSM

The 3-bit output of the phase control FSM control the output of MUX. The states of FSM are encoded as one-hot-state and used as the output signal. It rotates the clock phases counterclockwise or clockwise depending on whether lead<sub>OV</sub> or lag<sub>OV</sub> is 1.

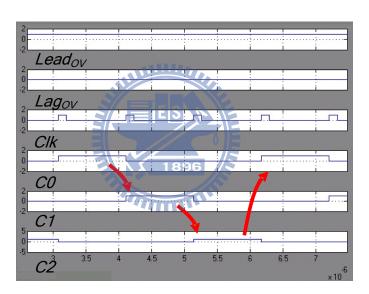


Fig. 3-10: Simulation of phase control FSM

#### 3.2.5. Phase Rotation Scheme

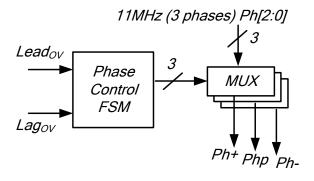


Fig. 3-11: Phase rotation scheme

The phase rotation is done by a FSM and a MUX, as shown in Fig. 3-12. It can be observed that when Ph2 is shifted to Ph1, the waveform of Ph+ is changed and it advances 1/3 chip time as compared with the original phase. Therefore, if Ph2 is shifted to Ph0, and Ph+ is delayed 1/3 chip time. Besides, Ph+ is used for sampling the data. It is also the clock for Barker code. The system clock is divided from it as well. Here we must carefully design the MUX of Ph+ to prevent the race problem.

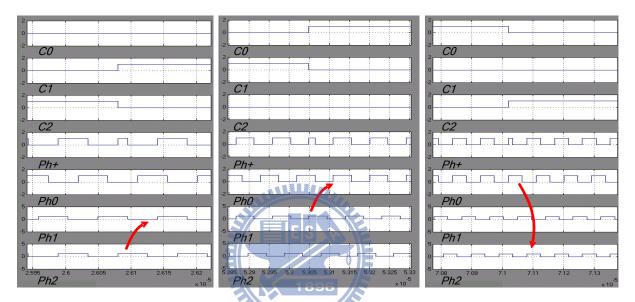


Fig. 3-12: Simulation of phase rotation

#### 3.3 Behavioral Simulation

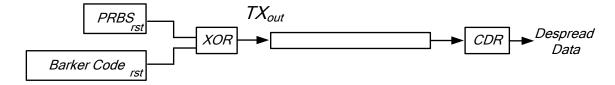


Fig. 3-13: Behavioral simulation of CDR

We use Simulink to execute the behavioral simulation of CDR, the transmitter consists of a Barker code generator, a pseudo random bit sequence generator (PRBS), and a XOR gate.

## 3.3.1. Chips Delay

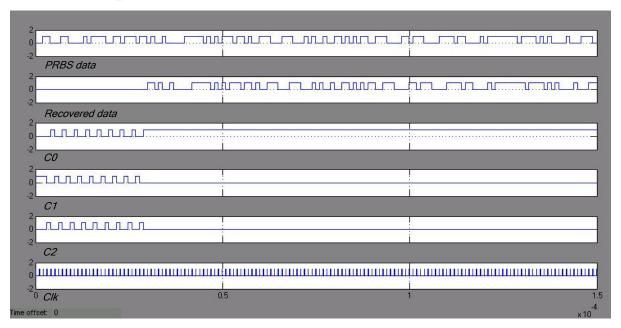


Fig. 3-14: Chips delay test

As Fig. 3-14 shows, we delay the input PRBS data to test whether the code acquisition loop functions well. It shows that the CDR shifts the phases for 26 times, that is, the CDR delays 8 and 2/3 chip time to get synchronization.

#### 3.3.2. Jitter Effect

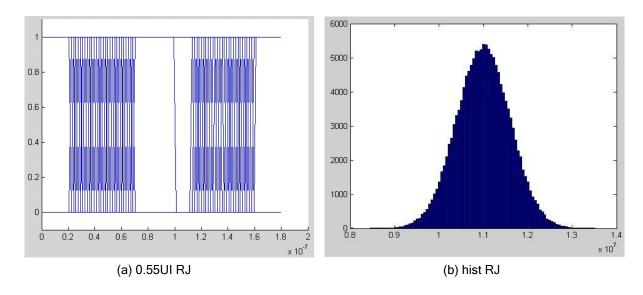


Fig. 3-15: Random jitter

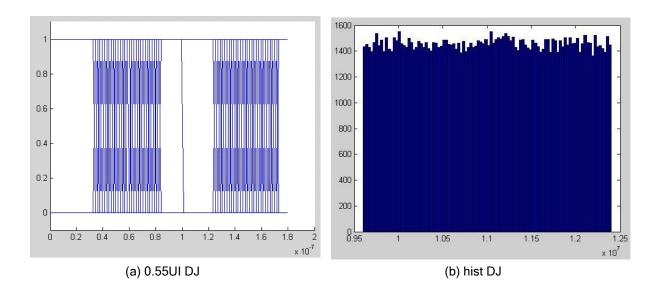


Fig. 3-16: Deterministic jitter

We produce the effect of jitter on the PRBS data at transmitter end to see whether the CDR can recover the data correctly. We put 0.55UI random jitter and 0.55UI deterministic jitter on the data. The RJ is in Gaussian distribution and DJ in is uniform distribution. As Fig. 3-17 shows, the CDR still can recover the data and the phase of clock shifted due to jitter.

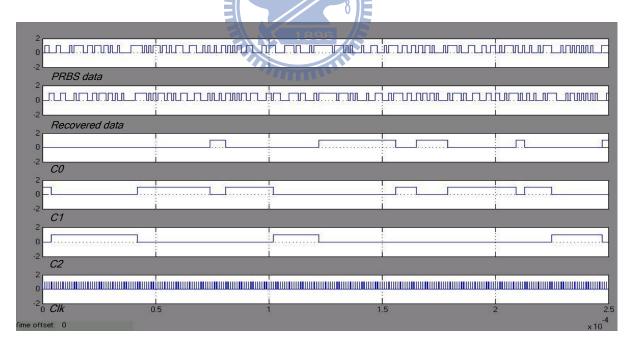


Fig. 3-17: Simulation of RJ and DJ

#### 3.3.3. Frequency Error

The frequency error is a constant frequency difference between the clock of the

transmitter and the clock of the receiver. It causes the data to drift from the sampling clock. So the system must have to shift the phase in the same way in order to fix the error. We set  $\pm 1000$ ppm frequency error at transmitter end. 1000ppm means that there is 1 bit overflow or underflow every 1000 data bits. The run-length in test is 250 data bits, so there will have 1/4 bits overflow or underflow, that is, 2.75 chips overflow or underflow.

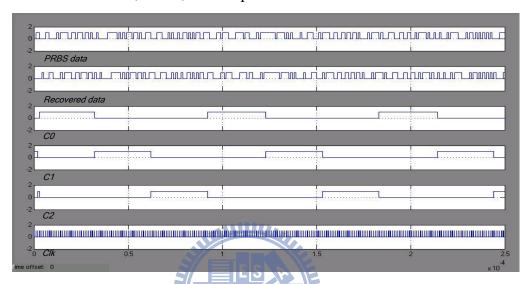


Fig. 3-18: -1000ppm (clock of the data)

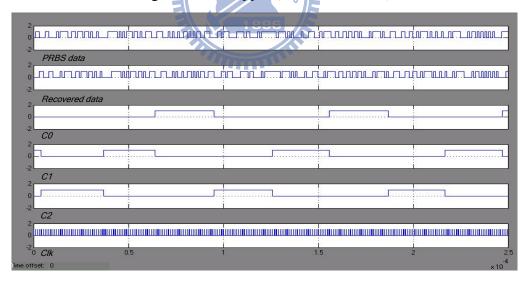


Fig. 3-19: +1000ppm (clock of the data)

As we can see from Fig. 3-18 and 3-19, the phases shift clockwisely or counterclockwisely due to the positive and negative frequency error. When the data is faster than the clock, it will shift counterclockwisely such as C1, to C2 and to C0. We put the jitter and frequency error together into the data and the CDR still can recover the data correctly.

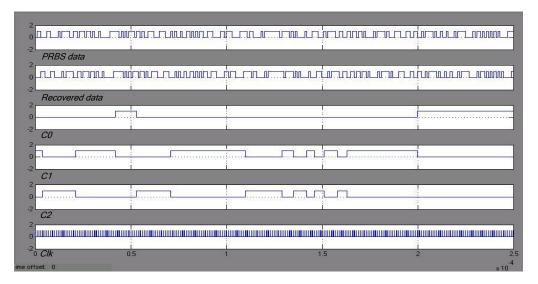


Fig. 3-20: +1000ppm with DJ and RJ

#### 3.4 Transceiver Simulation

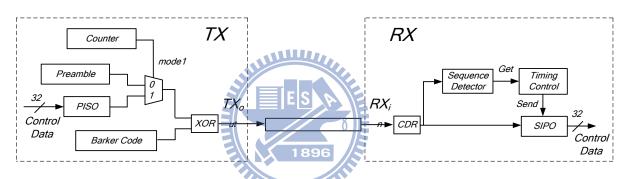


Fig. 3-21: Transceiver system

The goal of our design is to send 32-bit control signal from the transmitter through the power line of the automobile. Then, we despread the chips and send the data to the driver circuit. We know that the data length is 32-bit. So the transmission is burst-mode. Therefore every time a set of control data is sent, the CDR needs to resynchronize each time. Preamble is a sequence of known bits which are sent in each frame. It is used for synchronization. The length of the preamble should be long enough to let the CDR synchronize and recover the data properly. Consequently, at transmitter end, the preamble is sent before the control data and a counter is used for select the output of the MUX. At receiver end, a sequence detector is used to detect the pattern of preamble and control the serial-in parallel-out (SIPO) circuit. A timing control circuit between them is used to control the clock of the SIPO circuit, for it

sends out the control data correctly. Fig. 3-22 shows the transmitter end, and fig. 3-23 shows the receiver end.

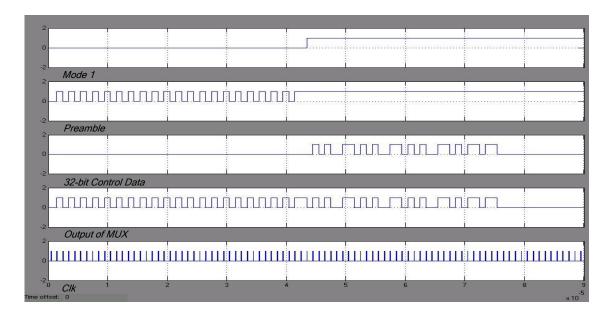


Fig. 3-22: Transmitter



Fig. 3-23: Receiver

As shown in Fig. 3-22, when mode1 becomes high, the data will be serialed out. In Fig. 3-23, we can see that the CDR is in code acquisition at first. To recover the data, the preamble must be detected to set the Get and Send signals and trigger the describing. P0 to P3 are the first four data bits.

## **Chapter 4**

# Circuit Implementation of the Transceiver



### 4.1 Transmitter

#### 4.1.1. Barker Code Generator

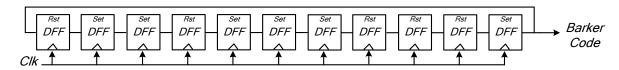


Fig. 4-1: Barker code

As Fig. 4-1 shows, the Barker code generator is a shift register which the output is connected to the input. The sequence of Barker code is {00011101101}, it is implemented by setting and resetting the DFF.

#### 4.1.2. Serializer

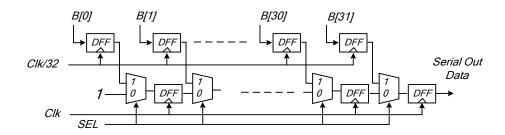


Fig. 4-2: Serializer

The serializer is a conventional type. It consists of two DFFs and a 2-to-1 MUXs. The Clk/32 will load 32-bit control data firstly and SEL will be high for a period of Clk and load the data to the input of DFFs. The DFFs serial out the data. The SEL signal is generated from the circuit shown in Fig. 4-3.

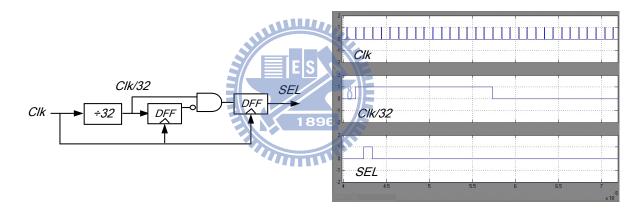


Fig. 4-3: Control signal for serializer

#### 4.1.3. Preamble and Counter

The preamble sequence is {1010......1011}. For the worst case, the synchronization for CDR needs to delay 11 chip times. It is equivalent to 33 times of phase shifting. We overdesign the sequence length to prevent any problem which may influence the synchronization. Therefore, the length is 42-bit. We use a 6-bit synchronous counter to generate the preamble and control the MUX and the serializer of the transmitter. We use the least significant bit (LSB) of the counter as the preamble sequence but the last two 11 pattern is different from the other. When the counter counts to 41, the NOR output will be high and

trigger the DFF to make the OR output high. So, the preamble will be high all the time. The control signals of the MUX and the serializer are produced in the same way. Both signals will be high when the counter counts to 43 and 40.

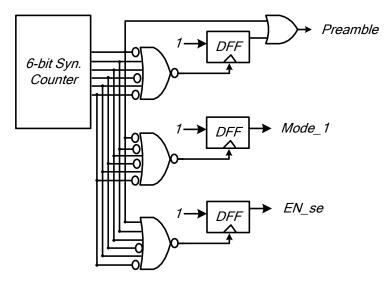


Fig. 4-4: Preamble and counter

#### 4.2 Receiver

#### 4.2.1. Accumulator

The accumulator consists of four half adders and a 4-bit DFFs. The input of accumulator is 1-bit and its output is 4-bit. The carry-out of the half adder connects to the input of the next half adder except the last one. The DFFs store the sums of the half adders for the next time, their output connect to the input of the half adder.

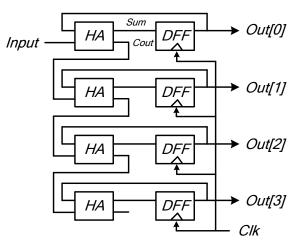


Fig. 4-5: Accumulator

#### 4.2.2. Absolute Value Circuit

According to the Fig. 3-2, we have already known the input and output of the absolute value circuit. It is a combinational circuit. The truth table is as follow:

X <sub>4</sub>	<b>X</b> <sub>3</sub>	$X_2$	X <sub>1</sub>	Y <sub>4</sub>	<b>Y</b> <sub>3</sub>	Y <sub>2</sub>	Y <sub>1</sub>
1	0	1	1	1	0	1	1
1	0	1	0	1	0	0	1
1	0	0	1	0	1	1	1
1	0	0	0	0	1	0	1
0	1	1	1	0	0	1	1
0	1	1	0	0	0	0	1
0	1	0	1	0	0	0	1
0	1	0	0	0	0	1	1
0	0	1	1	0	1	0	1
0	0	1	0	0	1	1	1
0	0	0	1	1	0	0	$ ota{\hspace{-0.03cm}/}{\hspace{-0.03cm}/} $
0	0	0	0	1/	0	ų.	10

$$\begin{split} Y_1 &= 1 \\ Y_2 &= \left( \overline{X_1 \oplus X_4} \right) \bullet \overline{X_2} \bullet \overline{X_3} \\ &+ \left( X_2 \oplus X_3 \right) \bullet \overline{X_1} \bullet \overline{X_4} \\ &+ \left( X_3 \oplus X_4 \right) \bullet X_1 \bullet X_2 \\ Y_3 &= \left( X_2 \oplus X_4 \right) \bullet \overline{X_3} \\ Y_4 &= \left( \overline{X_2} \oplus \overline{X_4} \right) \bullet \overline{X_3} \end{split}$$

Fig. 4-6: Absolute value circuit

#### 4.2.3. Threshold Detector

The threshold detector is a 4-bit magnitude comparator. The output of the comparator will be high when the input is greater than the value we have set. The value is 3 for the code acquisition and code tracking loop and 6 for despread the data.

#### 4.2.4. Confidence Counter

The confidence counter is encoded as an one-hot-state state machine. The input lead and lag signals control the output of the MUX in order to shift 1 left or right. Fig. 4-7 shows the circuit of confidence counter.

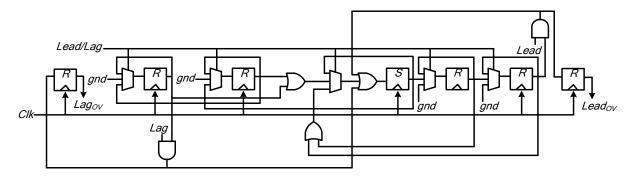


Fig. 4-7: Confidence counter

#### 4.2.5. Phase Control FSM

The circuit of the phase control FSM is similar to the confidence counter, we use the state assignment as its output.

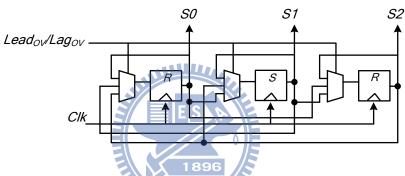


Fig. 4-8: Phase control FSM

#### 4.2.6. Clock Generator

The chip rate is 11MHz and the data rate is 1MHz. We need to generate three phases of 11MHz. The clock generator is a Johnson counter, it uses 66MHz clock to generate 6 phases of 11MHz, and we only take 3 of them as the sampling clocks.

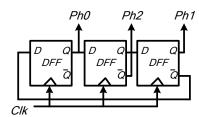


Fig. 4-9: Johnson counter

#### 4.2.7. Delay Adjustment

As we have mentioned before, the phase rotation needs a proper delay to prevent the race

condition. So we design the following circuit to realize the delay signal.

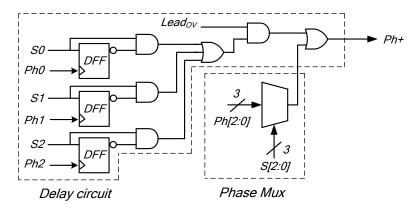


Fig. 4-10: Delay circuit for Ph+ phase

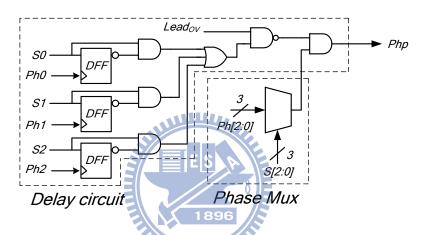


Fig. 4-11: Delay circuit for Php phase

#### 4.2.8. Sequence Detector

The sequence detector is used to detect the sync pattern {1011} in the preamble sequence. Its state diagram is as follows. When it detects the pattern, it will send a Get signal to indicate that the following data are the control signals. The serializer will start to function.

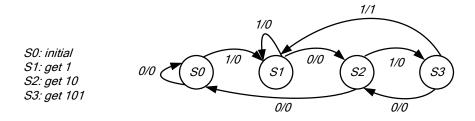


Fig. 4-12: State diagram of sequence detector

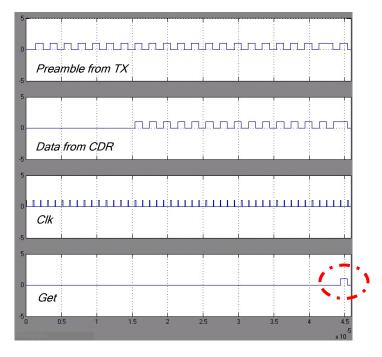


Fig. 4-13: Simulation of the sequence detector

#### 4.2.9. Timing Control

The timing control circuit is used to control the divider and the deserializer. At first, the sequence detector recognizes the preamble and sends the Get signal. Then, the timing control circuit receives the Get signal and enables the divider and the deserializer by the Send signal. The divider divides the retime clock by 32 and generates Clk/32 used in the deserializer. The deserializer starts to receive the 32-bit control data and send to the backend driver circuit in parallel form. After sending out the 32-bit control data, the timing control circuit will disable the divider and deserializer. The deserializer stores the control data until the next new control data is sent.

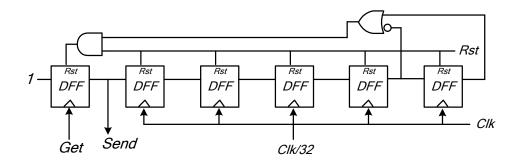


Fig. 4-14: Timing control circuit

#### 4.2.10. Deserializer

The deserializer is a conventional one. The simplest structure is suitable enough for our application.

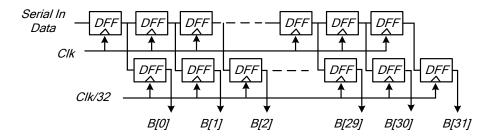


Fig. 4-15: Deserializer

#### 4.3 Circuit Simulation of the Transceiver

We use SPICE to perform the circuit simulation of the entire system. Firstly, the channel is assumed to be ideal. We build up the transmitter and the receiver. Then, we adjust the start time of the transmitter and let the transmitted chips have some delay with respect to the local Barker code at receiver end. We test whether the phase rotation scheme can recover the data correctly or not. As shown in Fig. 4-17, lead<sub>OV</sub> has been high for 14 clock period. It means that the phase rotates 14 times and local barker code is delayed about 4 chip time for the synchronization.

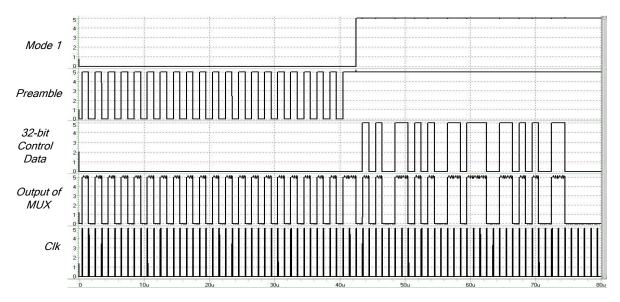


Fig. 4-16: Simulation of the transmitter

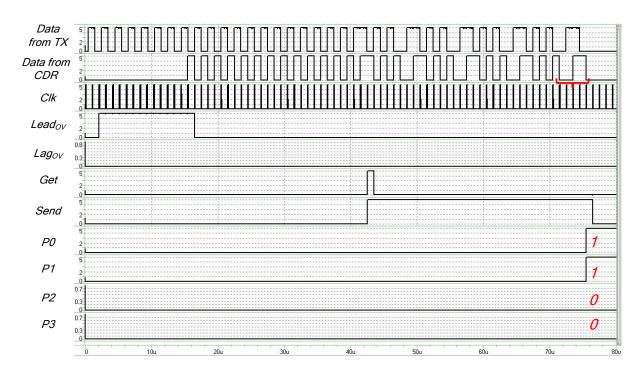


Fig. 4-17: Simulation of the receiver

We can observe that the backend circuit can function well to output the data in parallel.

Next, we change the frequency of clock at the transmitter end to examine if the CDR can shift the phases to correct the frequency error or not.

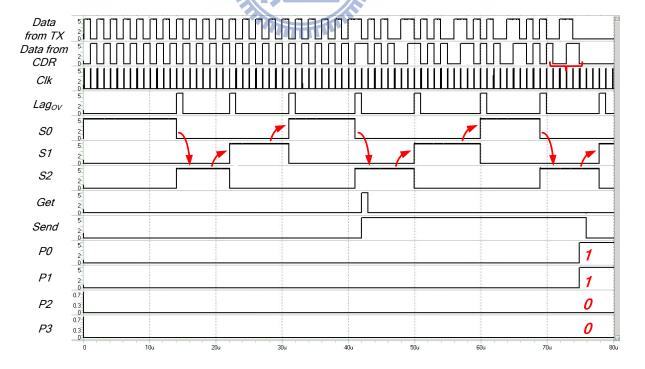


Fig. 4-18: Simulation of the frequency error

We let the clock of the transmitter faster than the clock of receiver. The total run length is

80 data bits. We set the frequency error to 3250ppm. It will underflow 1 chip about every 307.7chips. The total number is about 880 chips. So it will underflow 2.86 chips. As the Fig. 4-18 shows, lag<sub>OV</sub> happens 8 times as the phase rotates 8 times. It means that the system has shifted 2.67 chips. The direction of rotation is from S0, to S2, and to S1 counterclockwisely. All the things are just like we expected.



## Chapter 5

# Realization of the System on FPGA and PCB



#### 5.1 Realization of the Transceiver

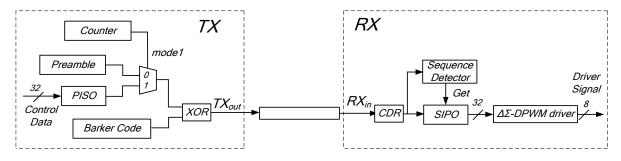
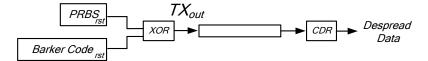


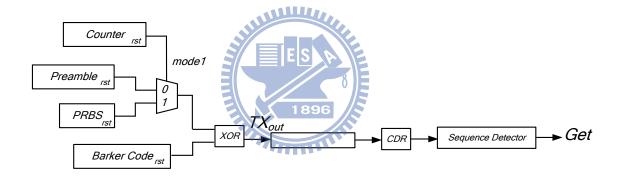
Fig. 5-1: Transceiver with backend control circuit

The whole application is shown in Fig. 5-1. The purpose is to use the delta-sigma digital pulse width modulator ( $\Delta\Sigma$ -DPWM) to generate PWM signal to control the Power MOS switches. The switches are used to control electric motors, LEDs, or buck converters. The electric motors and LEDs are in open loop control. The buck converters are in closed loop

control. These three components have their own operation frequency. Under this frequency, we can change the on-time duty or PWM of motors or LEDs to control the on-off cycling frequency, on-off duty cycle or dimming. The input of backend control circuit is 32-bit. The data transmission method is through the existing power line on the car to deliver the data. We use DSSS modulation to spread data into chips in order to decrease the noise interference. Basically a transmitter and a receiver for DSSS are what we need. Therefore, in order to realize this idea, we build a prototype using FPGA and discrete components on a bread board.



#### (a) Test of data recovery



#### (b) Test of detection of preamble

Fig. 5-2: Test for CDR

As shown in Fig. 5-2, we use the pseudorandom bit sequence (PRBS) encoded by barker code to test whether the CDR can despread the chips correctly. And then, we put the preamble signal before the PRBS data to test the sequence detector. When it receives the recovered data and detects the preamble, it sends out the Get signal to indicate the detection. Besides, we use a single cord wire as the transmission channel.

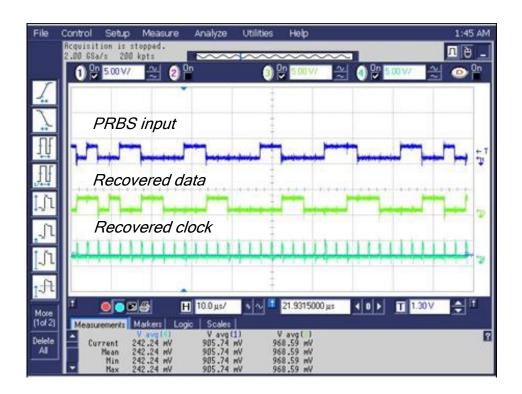


Fig. 5-3: Test of data recovery

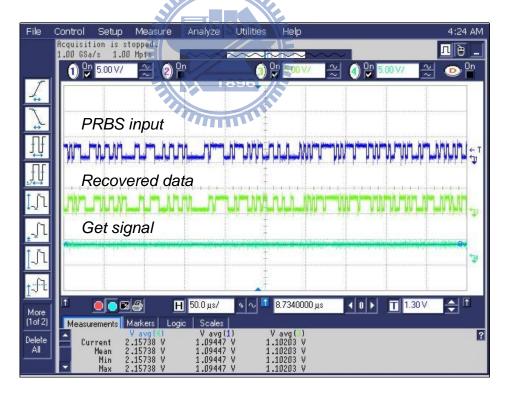


Fig. 5-4: Test of detection of preamble

Fig. 5-3, 5-4 shows that the system can detect the preamble signal and recover the data correctly. Therefore, the transmitter and receiver can function well.

#### 5.2 SERDES and Backend Control Circuit

Now we add a serializer to the transmitter and a deserializer and the  $\Delta\Sigma$ -DPWM control circuit to the receiver, as shown in Fig. 5-5. The simulation procedure is as follows. Set the 32-bit control signal and the transmitter serial out the preamble and the control data. Both of them are encoded by Barker code. Next, the CDR recovers the data and SIPO sends out the 32-bit data in parallel. Finally, the backend control circuit produces the desired PWM control signal. Besides, we use the power line as the transmission channel. Here the DC 12V has not plugged in yet.



Fig. 5-5: Test of SERDES and backend control circuit

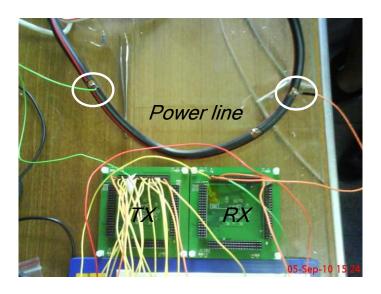


Fig. 5-6: Power line added

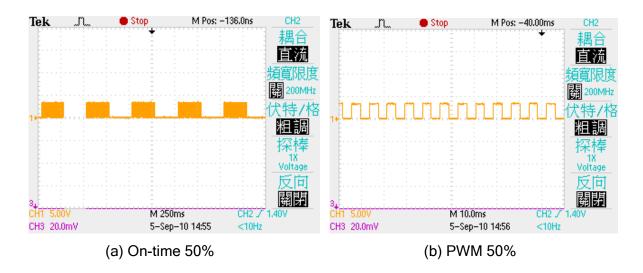


Fig. 5-7: On-time 50% and PWM 50%

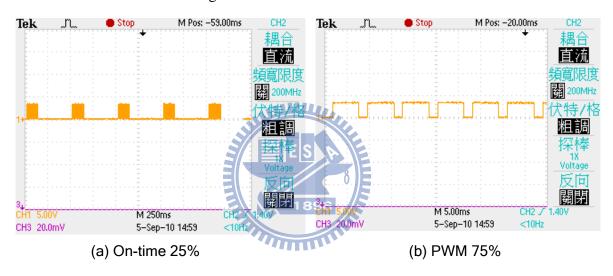


Fig. 5-8: On-time 25% and PWM 75%

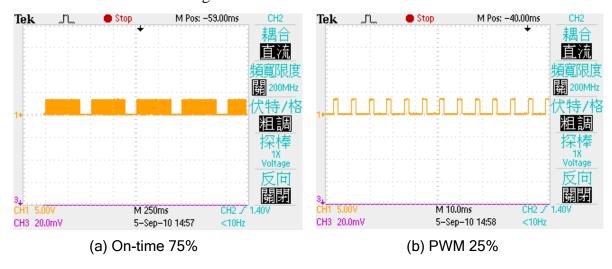


Fig. 5-9: On-time 75% and PWM 25%

Fig. 5-7, 5-8, 5-9 show that the control signals for nMOS switch, we change the on-time

duty and PWM in order to control the LED. The on-time duty allows LEDs to be turn on for 25%, 50% or 75% of a period time. And we can adjust the PWM signal of the on-time, it can change the brightness of LEDs to achieve the function of dimming.

### **5.3** Coupling Circuit with Power Line

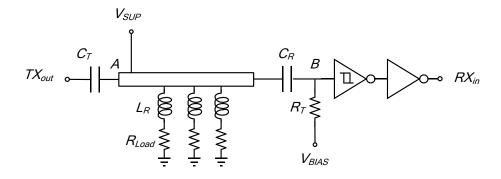


Fig. 5-10: Coupling circuit with power line

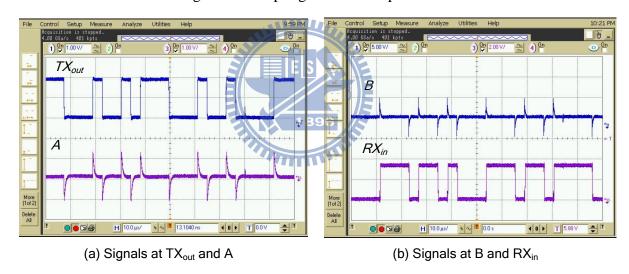


Fig. 5-11: Recovered PRBS signal

Because the power line is used to transmit the DC 12V supply to every electric appliances on the car. In order to transmit the power and data on the same cable, the inductors and capacitors are used here. We use capacitive coupling to transmit signal. It isolates the output of the FPGA from the direct contact to the large voltage on the cable. So, the common mode level can be different on both sides. The inductor is used to filter out the AC signal which is the transmitted signal on the wire. As L and C form a high-pass filter, the signal on the cable becomes pulse shape. Therefore, we use a Schmitt trigger inverter to recover the

pulse signal. We choose  $22\mu$  H of the inductance and  $10\mu$  F of the capacitance. Moreover, we bias the other side of the capacitor to move the pulse signal to the common mode of the inverter. Hence the inverter can recover the signal correctly. In addition, we use cement wire-wound resistors to simulate the load current as the real situation.

As shown in Fig. 5-11, the transmitted signal is recovered to the original signal. It means that the front-end circuit can work accurately. Then we transmit the control signal through the channel and the front-end circuit. Fig. 5-12 shows control signal for the nMOS switch.

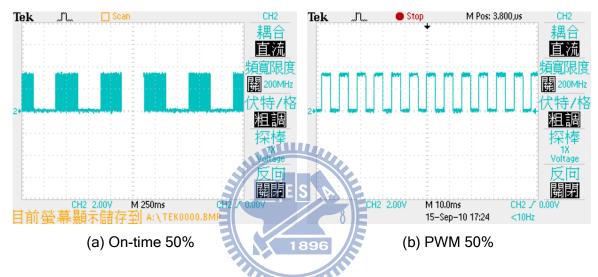


Fig. 5-12: On-time 50% and PWM 50%

## 5.4 Integration of the System

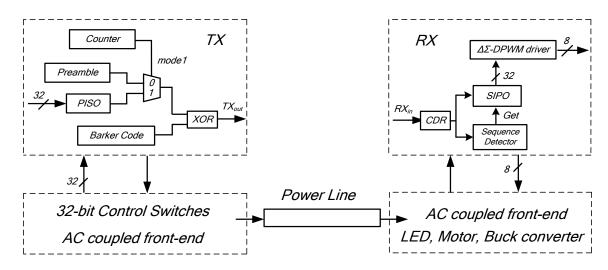


Fig. 5-13: System of the application

We build the entire system including the LED, motor, buck converter and control

switches. All of these circuits are realized on the printed circuit boards except for the front-end circuit. For initial test, the power of the LED, motor and buck converter comes from the DC supply separated from the cable. The cable is still connected with the cement wire wound resistors and low DC voltage.

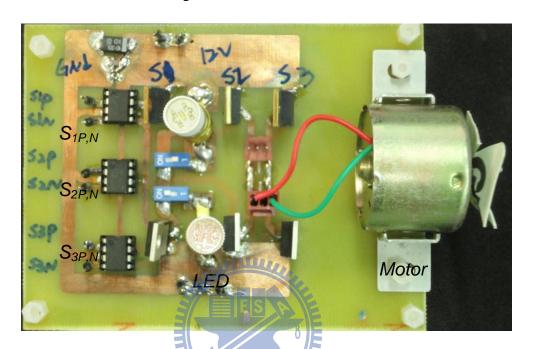


Fig. 5-14: LED, motor, and three pairs of switches

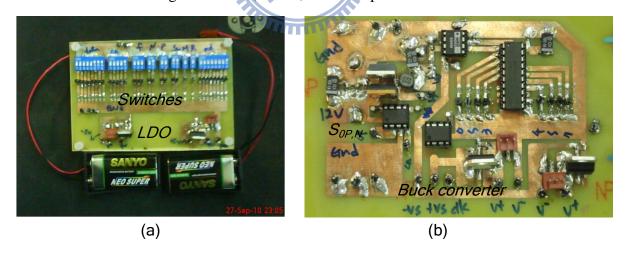


Fig. 5-15: Buck converter, LDO, switches for transmitter and buck converter

The low-dropout regulators are used for switches and FPGA board. Because we need 3.3V for both devices, so we can use regulators to generate 3.3V from any voltage and supply current to the devices. After setting up for the test, we adjust the switches to send out the desire control data. Then we test the brightness, on-off duty cycle and on-off frequency of the

LED, and the rotation speed of the motor, and the target voltage of the buck converter.



Fig. 5-16: Results from buck converter

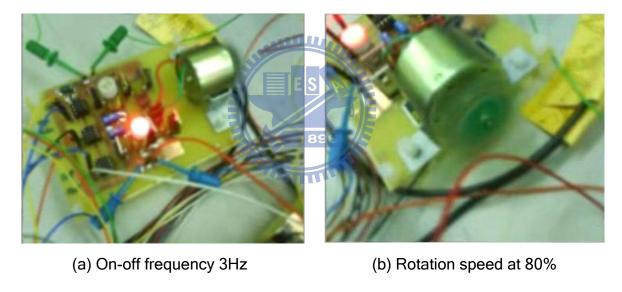


Fig. 5-17: Results from LED and motor

We have tape the process of test. These are the pictures derived from video. Fig. 5-16 shows the output voltage of the buck converter, 1.5V and 9V respectively. Fig. 5-17 shows the LED is flashing at 3Hz and motor is rotating at 80% of the full speed.

#### **■** Final Test

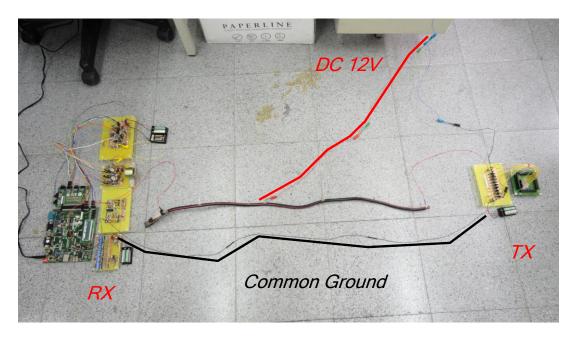


Fig. 5-18: Environment of the final test

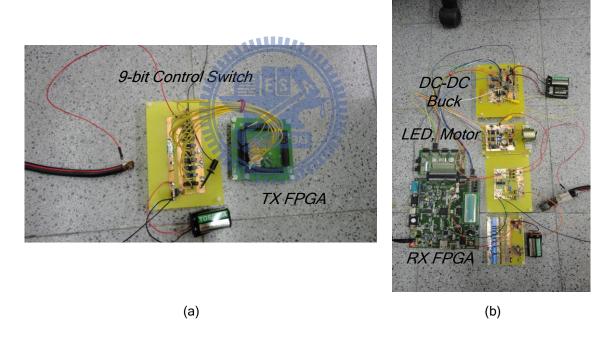


Fig. 5-19: Sub-blocks of the final test

Fig. 5-18 shows the environment set up for the final test. We connect the DC 12V supply to the power line as the practical situation. The contact of the receiver end of the cable is used for data receiving and supplies the power for the LEDs, motors and buck converters at the same time. At the transmitter end, we reduce the 32-bit control signal to only 5-bit control code. We list all the conditions for test, and build a look-up table for that. Consequently, we trim the complexity of circuits and wiring. We have demonstrated and taped the data



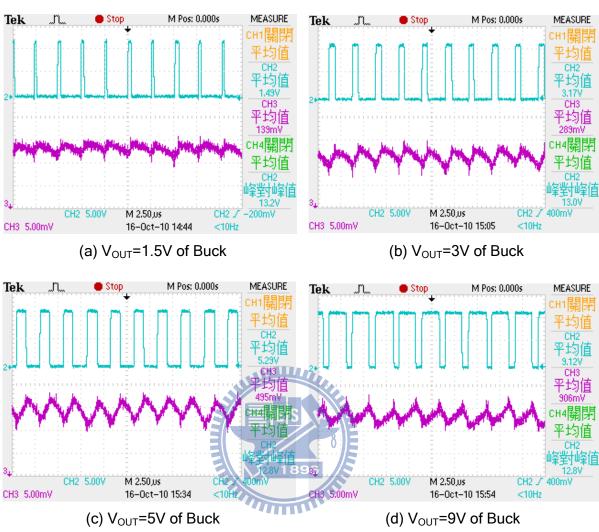


Fig. 5-20 Output voltage of the buck converter

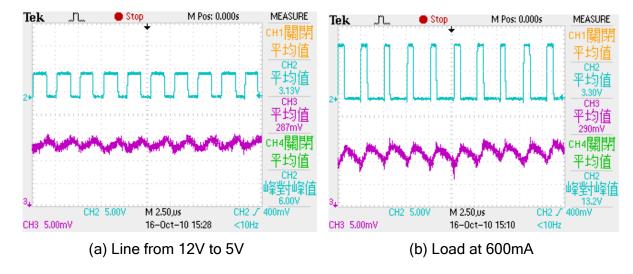


Fig. 5-21: Load and line regulation at output voltage 3V

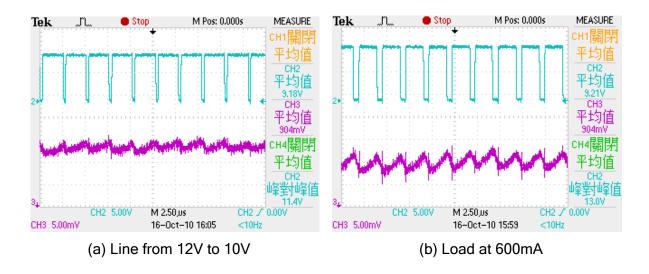


Fig. 5-22: Load and line regulation at output voltage 9V

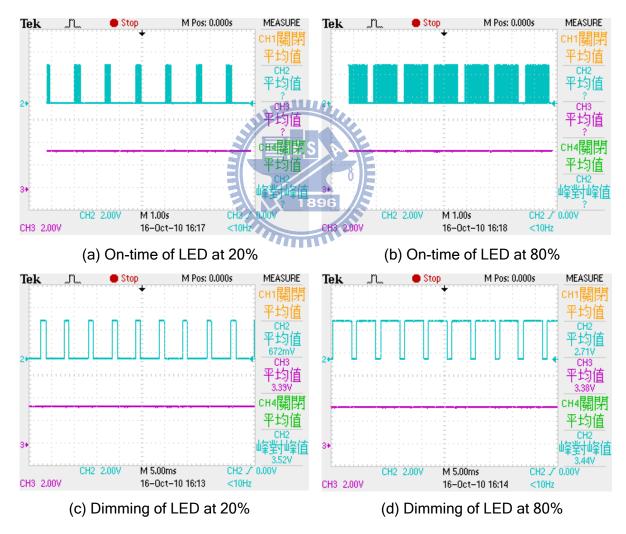


Fig. 5-23: On-time and dimming control of LED

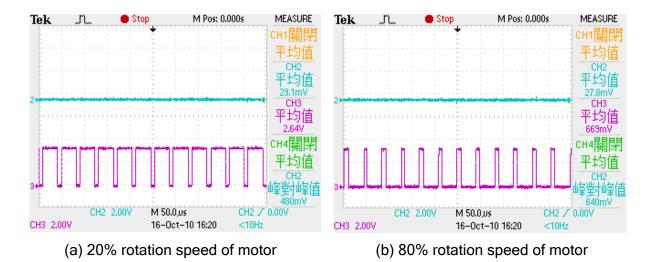


Fig. 5-24: Rotation speed of motor

Fig. 5-20 shows the output voltage of the buck converter from 1.5V to 9V. We test the line and load regulation under these voltages. Fig. 5-21 shows the line regulation from 12V to 5V at output voltage 3V. We can see the duty is changed. Fig. 5-22 shows the line regulation from 12V to 10V at output voltage 9V. The duty cycle is changed as well, and load regulation is the same as the one without the load. Fig. 5-23 shows the control of the LED. Fig. 5-24 shows the control of the motor. Table 2.2 is the look-up table of the control functions.

Table 2.2

Con[4:0]	function	1896	
00000_(0)	1.5V	10000_(16)	LED (NMOS) dimming 80%
00001_(1)	3V	10001_(17)	LED (NMOS) 6Hz
00010_(2)	5V	10010_(18)	LED (NMOS) 3Hz
00011_(3)	9V	10011_(19)	LED (NMOS) 1.5Hz
00100_(4)	LED (PMOS) dimming 20%	10100_(20)	LED (NMOS) 0.75Hz
00101_(5)	LED (PMOS) dimming 50%	10101_(21)	LED (NMOS) on-time 20%
00110_(6)	LED (PMOS) dimming 80%	10110_(22)	LED (NMOS) on-time 50%
00111_(7)	LED (PMOS) 6Hz	10111_(23)	LED (NMOS) on-time 80%
01000_(8)	LED (PMOS) 3HZ	11000_(24)	Motor (s2) (PMOS= on) (NMOS= off)
01001_(9)	LED (PMOS) 1.5Hz	11001_(25)	Motor (s3) (PMOS= off) (NMOS= 20%)
01010_(10)	LED (PMOS) 0.75Hz	11010_(26)	Motor (s3) (PMOS= off) (NMOS= 50%)
01011_(11)	LED (PMOS) on-time 20%	11011_(27)	Motor (s3) (PMOS= off) (NMOS= 80%)
01100_(12)	LED (PMOS) on-time 50%	11100_(28)	Motor (s3) (PMOS= on) (NMOS= off)
01101_(13)	LED (PMOS) on-time 80%	11101_(29)	Motor (s2) (PMOS= off) (NMOS= 20%)
01110_(14)	LED (NMOS) dimming 20%	11110_(30)	Motor (s2) (PMOS= off) (NMOS= 50%)
01111_(15)	LED (NMOS) dimming 50%	11111_(31)	Motor (s2) (PMOS= off) (NMOS= 80%)

## Chapter 6

### **Conclusion**



In this thesis we have proposed a data communication method with power line. We use DC 12V power line of the automobile as the transmission media. We apply the capacitive coupling to transmit and receive the data signal. We can expect that the signal will be corrupted by various noises existing in the channel. So we use direct sequence spread spectrum (DSSS) modulation technique to encode the control data. We have designed a DSSS transceiver to perform the function of spreading and despreading. We choose 11-bit Barker code which is also used in 802.11 wireless communications as the spreading sequence. We have designed a clock and data recovery circuit which combines the phase picking for clock recovery and byte synchronization for data recovery. We utilize the shifting of three phases of the clock to achieve 1/3 chip time delay which is enough for code acquisition. We use a pair of early and late phases to track the frequency error between transmitter (TX) and receiver

(RX), and a digital loop filter is used to filter out the high frequency noise or jitter. The 32-bit SERDES and backend control circuit are realized to perform the serial transmission and PWM control of LED, motor and buck converter. Finally, we use FPGA and PCB to realize the whole system with power line. We have demonstrated and taped the data transmission and controlling of the devices. We have proved the idea of power line communication for automotive electronics is feasible.

The following is the specification table for the DSSS transceiver. The three oversampling phases are generated from the global clock of the FPGA. For consideration of measurement, we lower the chip rate and the data rate to ease requirement of the front-end circuit.

Table 2.3 Specification for the DSSS transceiver

Burst-data length	75 bit		
FPGA global clock	50MHz		
PN code	11-chip Barker code		
Chip rate	260KHz		
Data rate	23.67KHz		
Counter size N	3 (continuous type)		
Inductor (front-end)	22µ H		
Capacitor (front-end)	10µ F		

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