1 Introduction

Let A and B be two n-by-n complex matrices which are unitarily equivalent, that is, $B = U^*AU$ for some unitary matrix U. It is easily seen that matrices A^*A and B^*B have equal traces. Trace of the product of a matrix and its own conjugate is thus an example of a unitary invariant. More generally, consider the multiplicative semigroup W generated by the noncommuting variables x and y. We call an element of W a word and denote it by w(x, y). If there is no confusion over the underlying variables one will also simply denote it by w. For example, $w(x, y) = y^2 x^3 y$ is such a word. In general a word w can always be written as $w = y^{i_1} x^{j_1} \dots y^{i_n} x^{j_n}$ where $i_1, j_1, \dots, i_n, j_n$ are positive integers except that i_1 or j_n may be zero. We say that $i_1, j_1, \dots, i_n, j_n$ are the exponents of a word w and w is of length n if i_1 and j_n are both different from zero. Otherwise w is said to be of length n - 1. If n = 1 and $i_1 = j_1 = 0$, we call w an empty word.

Suppose now we substitute x and y by A and A^* and regarding $w(A, A^*)$ as an n-by-n matrix. Same thing for $w(B, B^*)$. (If the word is empty we assign $w(A, A^*) = w(B, B^*) = I_n$.) One immediately sees that $tr(w(A, A^*)) = tr(w(B, B^*))$ for any word if A and B are unitarily equivalent. W. Specht [S] proved that the converse is also true. That is, the set $\{tr(w(A, A^*)) : w(x, y) \text{ is any word in } x \text{ and} y\}$ completely determines A up to unitary equivalence, and thus is a *complete* set of unitary invariants. There is a similar generalization due to N.Wiegmann [W] that considers not only two n-by-n matrices A and B but two finite sets of n-by-n matrices $\{A_1, A_2, \ldots, A_t\}$ and $\{B_1, B_2, \ldots, B_t\}$. In this situation we must consider words w with noncommuting variables in $x_1, y_1, x_2, y_2, \ldots, x_t, y_t$. It states that there exists a unitary matrix U such that $U^*A_iU = B_i$ for $i = 1, 2, \ldots, t$ if and only if for every word $w(x_1, y_1, x_2, y_2, \ldots, x_t, y_t)$ we have $tr(w(A_1, A_1^*, A_2, A_2^*, \ldots, A_t, A_t^*)) =$ $tr(w(B_1, B_1^*, B_2, B_2^*, \ldots, B_t, B_t^*))$.

The result of Specht gives an *infinite* set of complete unitary invariants for an n-by-n matrix A since one can form infinitely many words in A and A^* . Later, C. Pearcy showed in [P1] that a finite set of words would suffice. More precisely, let $\omega(k)$ denote the set of words in the variables x and y in which the sum of the exponents does not exceed k. Then A is unitarily equivalent to B if $tr(w(A, A^*) = tr(w(B, B^*))$ for every word w in $\omega(2n^2)$. This set contains fewer than 4^{n^2} elements. Of course, this upper bound is still far from satisfactory. In the other direction, Bhattarcharya [B] proved that for matrices whose nonzero singular values have multiplicity one, a family of about $(2n)^n$ traces would suffice. She also showed that there exist $n^2 + 1$ complex-valued continuous functions on $\mathbb{M}_n(\mathbb{C})$ which form a complete set of unitary invariants for n-by-n matrices where $\mathbb{M}_n(\mathbb{C})$ denotes the algebra by all n-by-n copmlex matrices. It suggests that one would like to find a complete set of specific unitary

invariants with the size of the set being a polynomial in n. For small n, it is easy to see three traces of words suffice for 2-by-2 matrices. Pearcy [P2] showed that a set of nine words suffices for n = 3 and Sibirskiï [Si] improved this by finding a set of seven words which suffices and forms a minimal set.

In Chapter 2 of this paper, we survey cases for n = 2, 3 and show that

$$\operatorname{tr}(A), \operatorname{tr}(A^2), \operatorname{tr}(A^*A)$$

form a complete set of unitary invariants for any *n*-by-*n* matrix A with rank 1. This will cover the 2-by-2 case immediately. We also show that three words are fewest possible. That is, one cannot find a set with two traces of words that is complete. For any 3-by-3 matrix A, we prove that

$$\operatorname{tr}(A), \operatorname{tr}(A^2), \operatorname{tr}(A^3), \operatorname{tr}(A^*A), \operatorname{tr}(A^{*2}A), \operatorname{tr}(A^{*2}A^2), \operatorname{tr}(A^{*2}A^2A^*A)$$

form a complete set of unitary invariants. The proof here is quite different from the computational proof given in [P2] and gives a bit more. The set given above plus $tr(A^*A)^2$ is actually complete with respect to any *n*-by-*n* matrix *A* with rank 2. This will cover the case for 3-by-3 matrix readily.

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In Chapter 3, we give the result that for another special class of matrices (matrix whose eigenvectors are not orthogonal), a set with no more than $n^4 + 1$ words suffices to determine such matrices up to unitary equivalence. We denote the algebra generated by A and A^* over the complex numbers by $\operatorname{Alg}(A, A^*)$. This is the set of all polynomial expressions $p(A, A^*)$, where p(x, y) is any polynomial in the noncommuting variables x and y. The method lies in first proving $\{A^{*i}A^j : 0 \leq i, j < n\}$ spans $\operatorname{Alg}(A, A^*)$ for such an A. Then one modifies the method in [P1] somewhat to obtain that $\{\operatorname{tr}(A^{*i}A^jA^{*k}A^l) : 0 \leq i, j, k, l < n\} \cup \{\operatorname{tr}(A^n)\}$ forms a complete set of unitary invariants for A in this class.

2 Unitary equivalence for 2-by-2 and 3-by-3 Matrices

2.1 2-by-2 Matrices

First we state a basic lemma for general n-by-n matrices that will be applied repeatedly later.

Lemma 2.1 Let A and B be two n-by-n matrices. If $tr(A^i) = tr(B^i)$ for $1 \le i \le n$, then A and B have the same eigenvalues counting algebraic multiplicities and hence $tr(A^i) = tr(B^i)$ for all integers i.

Proof. Let a_1, \ldots, a_n be eigenvalues of A and b_1, \ldots, b_n be eigenvalues of B, each repeated according to its algebraic multiplicity. We define two classes of homogeneous polynomials $S_r, G_r : \mathbb{C}^n \to \mathbb{C}$ for $1 \leq r \leq n$ by

$$S_r(x_1, x_2, \dots, x_n) = \sum_{i=1}^n x_i^r$$

and

$$G_{r}(x_{1}, x_{2}, \dots, x_{n}) = \sum_{\substack{1 \leq i_{1} < i_{2} < \\ \cdots < i_{r} \leq n}} x_{i_{1}} x_{i_{2}} \dots x_{i_{r}}$$

I. Note that $S_{1} = G_{1}$.

We also define $G_0 = 1$. Note that $S_1 = G_1$.

The statement that $\operatorname{tr}(A^i) = \operatorname{tr}(B^i)$ for $1 \leq i \leq n$ is the same as $S_r(a_1, \ldots, a_n) = S_r(b_1, \ldots, b_n)$ for $1 \leq r \leq n$. We want to show that this implies $G_r(a_1, \ldots, a_n) = G_r(b_1, \ldots, b_n)$ for $1 \leq r \leq n$. In fact, we have the following so-called Newton's identities:

$$rG_r - S_1G_{r-1} + \dots + (-1)^r S_r G_0 = 0$$
(2.1)

for $1 \leq r \leq n$ (cf. [Pr, p.20]). Using $S_r(a_1, \ldots, a_n) = S_r(b_1, \ldots, b_n)$ for $1 \leq r \leq n$ together with (2.1), one concludes that $G_r(a_1, \ldots, a_n) = G_r(b_1, \ldots, b_n)$ for $1 \leq r \leq n$.

Now we have

$$\prod_{i=1}^{n} (x - a_i) = \sum_{i=0}^{n} (-1)^i G_i(a_1, ..., a_n) x^{n-i}$$
$$= \sum_{i=0}^{n} (-1)^i G_i(b_1, ..., b_n) x^{n-i} = \prod_{i=1}^{n} (x - b_i).$$
(2.2)

So $a_1, ..., a_n$ and $b_1, ..., b_n$ both represent the zeros of the same polynomial and hence must coincide.

To prove (2.1) we define another class of homogenous polynomials $K_r^{(m)} : \mathbb{C}^n \to \mathbb{C}$ for integers r and m such that r < n and $1 \le m \le n$ by

$$K_r^{(m)}(x_1, x_2, \dots, x_n) = \sum_{\substack{1 \le i_1 < i_2 < \dots < i_r \le n \\ i_1, i_2, \dots, i_r \neq m}} x_{i_1} x_{i_2} \dots x_{i_r}.$$

We also define $K_0^{(m)} = 1$ and $K_{-1}^{(m)} = 0$.

Notice that $G_r = x_m K_{r-1}^{(m)} + K_r^{(m)}$ for every m such that $1 \le m \le n$ and every r such that $0 \le r \le n$. So we have

$$S_i G_r = \sum_{j=1}^n x_j^i G_r = \sum_{j=1}^n x_j^i (x_j K_{r-1}^{(j)} + K_r^{(j)}).$$
(2.3)

Substituting (2.3) back into (2.1), we have

$$\begin{aligned} rG_r + \sum_{i=1}^r (-1)^i G_{r-i} S_i &= rG_r + \sum_{i=1}^r (-1)^i \sum_{j=1}^n x_j^i (x_j K_{r-i-1}^{(j)} + K_{r-i}^{(j)}) \\ &= rG_r + \sum_{j=1}^n \left[\sum_{i=1}^r (-1)^i x_j^{i+1} K_{r-i-1}^{(j)} + \sum_{i=1}^r (-1)^i x_j^i K_{r-i}^{(j)} \right] \\ &= rG_r - \sum_{j=1}^n x_j K_{r-1}^{(j)} + \sum_{j=1}^n \left[\sum_{i=1}^{r-1} (-1)^i x_j^{i+1} K_{r-i-1}^{(j)} + \sum_{i=2}^r (-1)^i x_j^i K_{r-i}^{(j)} \right] \\ &= rG_r - \sum_{j=1}^n x_j K_{r-1}^{(j)} + \sum_{j=1}^n \left[\sum_{i=1}^{r-1} (-1)^i x_j^{i+1} K_{r-i-1}^{(j)} + \sum_{i=1}^{r-1} (-1)^{i+1} x_j^{i+1} K_{r-i-1}^{(j)} \right] \\ &= rG_r - \sum_{j=1}^n x_j K_{r-1}^{(j)} = 0. \end{aligned}$$

Hence the assertion is proved.

Suppose we know beforehand that A and B have r common eigenvalues (counting algebraic multiplicities). Then it is clear that we only have to check $\operatorname{tr}(A^i) = \operatorname{tr}(B^i)$ for $1 \leq i \leq n - r$. We will use mostly Lemma 2.1 in the following form:

Corollary 2.2 Let A and B be two n-by-n matrices both of rank r. If $tr(A^i) = tr(B^i)$ for $1 \le i \le r$, then $tr(A^i) = tr(B^i)$ for all i.

Theorem 2.3 Let A and B be matrices both of rank 1. If tr(A) = tr(B) and $tr(A^*A) = tr(B^*B)$, then A is unitarily equivalent to B.

Proof. Since A and B are both of rank 1, $\operatorname{tr}(A) = \operatorname{tr}(B)$ guarantees that $\operatorname{tr}(A^i) = \operatorname{tr}(B^i)$ for any integer *i* and that A and B have equal characteristic polynomials of degree 2. So all it remains to check are words w with exponents all equal to 1. That is, we still need to check if $\operatorname{tr}(A^*A)^i = \operatorname{tr}(B^*B)^i$ for any integer *i*. Similarly, since A^*A and B^*B are both of rank 1, $\operatorname{tr}(A^*A) = \operatorname{tr}(B^*B)$ guarantees that $\operatorname{tr}(A^*A)^i = \operatorname{tr}(B^*B)^i$ for any integer *i*. Hence we conclude that for every word w we have $\operatorname{tr}(w(A, A^*)) = \operatorname{tr}(w(B, B^*))$. So by Specht's theorem A is unitarily equivalent to B.

Theorem 2.4 Let A and B be 2-by-2 matrices. If tr(A) = tr(B), $tr(A^2) = tr(B^2)$ and $tr(A^*A) = tr(B^*B)$, then A is unitarily equivalent to B.

Proof. Since $\operatorname{tr}(A) = \operatorname{tr}(B)$ and $\operatorname{tr}(A^2) = \operatorname{tr}(B^2)$ it follows that A and B have the same eigenvalues. Let λ be one of the common eigenvalues of A and B and let $A' = A - \lambda I_2$ and $B' = B - \lambda I_2$. One can easily check that $\operatorname{tr}(A) = \operatorname{tr}(B)$ and $\operatorname{tr}(A^*A) = \operatorname{tr}(B^*B)$ imply $\operatorname{tr}(A') = \operatorname{tr}(B')$ and $\operatorname{tr}(A'^*A') = \operatorname{tr}(B'^*B')$. Since both A' and B' are of rank 1, using Theorem 2.1 we obtain that A' is unitarily equivalent to B'. Hence A is unitarily equivalent to B as well.

and the

In Theorem 2.4 we use traces of three words to characterize 2-by-2 matrices up to unitary equivalence. Can we still make it better? The answer is no.

Theorem 2.5 Traces of two words do not suffice to determine a 2-by-2 matrix up to unitary equivalence.

Proof. Let A and B be two 2-by-2 matrices. Using the property that $\operatorname{tr}(PQ) = \operatorname{tr}(QP)$ for any matrices P and Q, one needs only to consider words w of the following two types: (1) $w(x,y) = y^{i_1}x^{j_1} \dots y^{i_n}x^{j_n}$ with positive integers $i_1, j_1, \dots, i_n, j_n$ and (2) $w(x,y) = y^n$ with positive integer n.

Take w_1 and w_2 to be any two words. First suppose that neither w_1 nor w_2 takes the form $(yx)^n$ for some integer n. This means that if w_1 is of the first type then there must exist some exponent of w_1 which is larger than 1. Now we may give a counterexample of

$$A = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} \text{ and } B = \begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}.$$

That $\operatorname{tr}(w_1(B, B^*)) = \operatorname{tr}(w_2(B, B^*)) = 0$ is trivial. Suppose that w_1 is of the first type. Then we have $\operatorname{tr}(w_1(A, A^*)) = 0$ because $A^2 = 0$. Suppose that w_1 is of the second type. Then we still have $\operatorname{tr}(w_1(A, A^*)) = 0$ as well. So we conclude that $\operatorname{tr}(w_1(A, A^*)) = 0$. Similarly, we have $\operatorname{tr}(w_2(A, A^*)) = 0$. So we have

$$\operatorname{tr}(w_1(A, A^*)) = \operatorname{tr}(w_2(A, A^*)) = \operatorname{tr}(w_1(B, B^*)) = \operatorname{tr}(w_2(B, B^*)) = 0$$

while A is NOT unitarily equivalent to B.

Now suppose that we have one of the words, say, w_1 equal to $(yx)^n$. If the sum of the exponents of w_2 is even, we may take

$$A = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \text{ and } B = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

as a counterexample. It is easily seen that

$$w_1(A, A^*) = w_1(B, B^*) = w_2(A, A^*) = w_2(B, B^*) = I_2.$$

Thus we again have

$$tr(w_1(A, A^*)) = tr(w_1(B, B^*))$$
 and $tr(w_2(A, A^*)) = tr(w_2(B, B^*))$

while A is clearly NOT unitarily equivalent to B.

Finally if we have $w_1 = (yx)^n$ for some integer n and the sum of the exponents of w_2 is odd, we may take

$$A = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \text{ and } B = \begin{pmatrix} i & 0 \\ 0 & -i \end{pmatrix}.$$

One can check that $\operatorname{tr}(w_1(A, A^*)) = \operatorname{tr}(w_1(B, B^*))$ and $\operatorname{tr}(w_2(A, A^*)) = \operatorname{tr}(w_2(B, B^*))$ while A is clearly NOT unitarily equivalent to B.

Since we have exhausted all possible cases, we conclude that traces of two words do not suffice to determine any 2-by-2 matrix up to unitary equivalence. \Box

2.2 3-by-3 Matrices

Let A and B be two n-by-n matrices of rank 2. Then tr(A)=tr(B) and $tr(A^2)=tr(B^2)$ guarantee that they have equal characteristic polynomials of degree 3 (unless n = 2, but this case is already solved). So we only have to check

$$\operatorname{tr}(w(A, A^*)) = \operatorname{tr}(w(B, B^*))$$

for words $w(x,y) = y^{i_1}x^{j_1}\cdots y^{i_n}x^{j_n}$ with $1 \leq i_1, j_1, \ldots, i_n, j_n \leq 2$. For the sake of convenience we denote P, Q, R, S words in W by

$$P(x,y) = yx, Q(x,y) = y^{2}x, R(x,y) = yx^{2}, S(x,y) = y^{2}x^{2}.$$
(2.4)

It is clear that every word we need to verify can be written as the product of P,Q,Rand S. Also for w a word in W, we will often denote $w(A, A^*)$ and $w(B, B^*)$ as w_A and w_B . Whenever we write $\operatorname{tr}(w_A)^{1,2} = \operatorname{tr}(w_B)^{1,2}$ it is to be understood that we mean $\operatorname{tr}(w_A)^i = \operatorname{tr}(w_B)^i$ for i = 1, 2.

Let $w^{(1)}, \ldots, w^{(8)}$ be words in W defined by

$$w^{(1)}(x,y) = x, \quad w^{(2)}(x,y) = x^2, \quad w^{(3)}(x,y) = x^3, \quad w^{(4)}(x,y) = P,$$

$$w^{(5)}(x,y) = Q, \quad w^{(6)}(x,y) = S, \quad w^{(7)}(x,y) = P^2, \quad w^{(8)}(x,y) = SP.$$
(2.5)

We are going to show that $\operatorname{tr}(w_A^{(i)}) = \operatorname{tr}(w_B^{(i)})$ for $1 \le i \le 8$ implies the unitary equivalence of A and B. This can be easily applied to arbitrary 3-by-3 matrices. There is a small difference in that $w^{(7)}$ is actually not needed in the 3-by-3 case.

Lemma 2.6 Let A and B be two n-by-n matrices both of rank 2. Suppose tr $(w_A^{(i)})$ =tr $(w_B^{(i)})$ for $1 \le i \le 8$ as defined in (2.5). Then for every word w of length 2 we have tr (w_A) =tr (w_B) .

Proof. Using $tr(A^*) = \overline{tr(A)}$ and tr(AB) = tr(BA) we only need to verify

$$\operatorname{tr}(QP)_{A} = \operatorname{tr}(QP)_{B}, \quad \operatorname{tr}(RQ)_{A} = \operatorname{tr}(RQ)_{B}, \quad \operatorname{tr}(QQ)_{A} = \operatorname{tr}(QQ)_{B},$$

$$\operatorname{tr}(SS)_{A} = \operatorname{tr}(SS)_{B}, \quad \operatorname{tr}(SQ)_{A} = \operatorname{tr}(SQ)_{B}.$$
(2.6)

For any complex numbers u and v, $uA + vA^*$ and $uB + vB^*$ are both matrices with rank at most 4 if A and B both have rank 2. So $tr(uA + vA^*)^i = tr(uB + vB^*)^i$ for $1 \le i \le 4$ and hence $tr(uA + vA^*)^i = 1896$

 $\operatorname{tr}(uB + vB^*)^i$ for any integer *i* and any complex numbers *u* and *v*. One can check that $\operatorname{tr}(w_A^{(i)}) = \operatorname{tr}(w_B^{(i)})$ for $1 \le i \le 7$ implies $\operatorname{tr}(uA + vA^*)^i =$

 $\operatorname{tr}(uB+vB^*)^i$ for $1 \leq i \leq 4$. Using the equality for i=5 and comparing the coefficients of u^2v^3 one gets

$$5(\operatorname{tr}(A^{*3}A^2) + \operatorname{tr}(QP)_A) = 5(\operatorname{tr}(B^{*3}B^2) + \operatorname{tr}(QP)_B).$$

However since $\operatorname{tr}(A^{*3}A^2)$ can be written as a linear combination of $\operatorname{tr}(S_A), \operatorname{tr}(Q_A)^*$ and $\operatorname{tr}(A^2)$ while $\operatorname{tr}(S_A) = \operatorname{tr}(S_B)$ and $\operatorname{tr}(Q_A) = \operatorname{tr}(Q_B)$, one gets $\operatorname{tr}(A^{*3}A^2) = \operatorname{tr}(B^{*3}B^2)$ and thus $\operatorname{tr}(QP)_A = \operatorname{tr}(QP)_B$ as well.

Next we apply $tr(uA + vA^*)^6 = tr(uB + vB^*)^6$ and comparing the coefficients of u^3v^3 to get

$$6(\operatorname{tr}(SP)_A + \operatorname{tr}(RQ)_A + \operatorname{tr}(A^{*3}A^3)) + 2 * \operatorname{tr}(P^3)_A = 6(\operatorname{tr}(SP)_B + \operatorname{tr}(RQ)_B + \operatorname{tr}(B^{*3}B^3)) + 2 * \operatorname{tr}(P^3)_B$$

Since $\operatorname{tr}(P^{1,2})_A = \operatorname{tr}(P^{1,2})_B$ while P_A and P_B are both matrices of rank 2, it implies that $\operatorname{tr}(P^i)_A = \operatorname{tr}(P^i)_B$ for any integer *i*. We also have $\operatorname{tr}(A^{*3}A^3) = \operatorname{tr}(B^{*3}B^3)$ which is

trivial and $tr(SP)_A = tr(SP)_B$ which is given. So we deduce that $tr(RQ)_A = tr(RQ)_B$.

Next we consider matrices $uA^*A + vA^*$ and $uB^*B + vB^*$ for any complex numbers u and v. Since $uA^*A + vA^* = A^*(uA + vI_n)$, $uA^*A + vA^*$ has rank 2 for any u and v. Similarly $uB^*B + vB^*$ has rank 2. So $\operatorname{tr}(uA^*A + vA^*)^{1,2} = \operatorname{tr}(uB^*B + vB)^{1,2}$ implies $\operatorname{tr}(uA^*A + vA^*)^i = \operatorname{tr}(uB^*B + vB^*)^i$ for any integer i. One could see that

$$\operatorname{tr}(A^{1,2}) = \operatorname{tr}(B^{1,2}), \operatorname{tr}(P^{1,2})_A = \operatorname{tr}(P^{1,2})_B, \operatorname{tr}(Q_A) = \operatorname{tr}(Q_B)$$

indeed imply $tr(uA^*A + vA^*)^{1,2} = tr(uB^*B + vB)^{1,2}$. Using the equality for i=4 and matching the coefficients of u^2v^2 one gets

$$4\mathrm{tr}(A^{*3}AA^*A) + 2\mathrm{tr}(Q^2)_A = 4\mathrm{tr}(B^{*3}BB^*B) + 2\mathrm{tr}(Q^2)_B.$$

However we have $\operatorname{tr}(A^{*3}AA^*A) = \operatorname{tr}(B^{*3}BB^*B)$ so we infer that $\operatorname{tr}(Q^2)_A = \operatorname{tr}(Q^2)_B$.

Now we consider $uA^{*2}A + vA$ and $uB^{*2}B + vB$ for any complex numbers u and v. For the same reason as before,

$$tr(uA^{*2}A + vA)^{1,2} = tr(uB^{*2}B + vB)^{1,2}$$

implies $\operatorname{tr}(uA^{*2}A + vA)^i = \operatorname{tr}(uB^{*2}B + vB)^i$ for any integer *i*. Using the equality for i=4 and comparing the coefficients of u^2v^2 one gets

$$4\mathrm{tr}(A^{*2}AA^{*2}A^{3}) + 2\mathrm{tr}(S^{2})_{A} = 4\mathrm{tr}(B^{*2}AB^{*2}B^{3}) + 2\mathrm{tr}(S^{2})_{B}.$$

Consider matrices $uQ_A + vA^2$ and $uQ_B + vB^2$ for any complex numbers u and v. One can check that indeed $\operatorname{tr}(uQ_A + vA^2)^{1,2} = \operatorname{tr}(uQ_B + vB^2)^{1,2}$, so we have

$$\operatorname{tr}(uQ_A + vA^2)^i = \operatorname{tr}(uQ_B + vB^2)^i$$

for every integer *i*. Using the equality for i=3 and matching the coefficients of uv^2 one gets $tr(A^{*2}AA^{*2}A^3)=tr(B^{*2}BB^{*2}B^3)$. Hence we conclude that $tr(S^2)_A=tr(S^2)_B$ as well.

Next we consider $uA^{*2}A^2 + vA^*A + tA^*$ and $uB^{*2}B^2 + vB^*B + tB^*$ for any complex number u, v and t. For the same reason as before,

$$\operatorname{tr}(uA^{*2}A^2 + vA^*A + tA^*)^{1,2} = \operatorname{tr}(uB^{*2}B^2 + vB^*B + tB^*)^{1,2}$$

implies tr $(uA^{*2}A^2 + vA^*A + tA^*)^i = tr(uB^{*2}B^2 + vB^*B + tB^*)^i$ for any integer *i*. One can check that

$$\operatorname{tr}(A) = \operatorname{tr}(B), \operatorname{tr}(P_A) = \operatorname{tr}(P_B), \operatorname{tr}(S_A) = \operatorname{tr}(S_B),$$

$$\operatorname{tr}(Q_A) = \operatorname{tr}(Q_B), \operatorname{tr}(SP)_A = \operatorname{tr}(SP)_B$$

indeed guarantee that $tr(uA^{*2}A^2 + vA^*A + tA^*)^{1,2} = tr(uB^{*2}B^2 + vB^*B + wB^*)^{1,2}$. Using the equality for i=3 and matching the coefficients of uvt one gets

$$3(\operatorname{tr}(SQ)_A + \operatorname{tr}(A^{*3}A^2A^*A)) = 3(\operatorname{tr}(SQ)_B + \operatorname{tr}(B^{*3}B^2B^*B)).$$

Since we have $\operatorname{tr}(A^{*3}A^2A^*A) = \operatorname{tr}(B^{*3}B^2B^*B)$, we deduce that $\operatorname{tr}(SQ)_A = \operatorname{tr}(SQ)_B$.

Lemma 2.7 Let A and B be two n-by-n matrices both of rank 2. Suppose $\operatorname{tr}(w_A^i) = \operatorname{tr}(w_B^i)$ for $1 \leq i \leq 8$. Let k be any integer. Assume that for every word w of length k we have $\operatorname{tr}(w_A^{1,2}) = \operatorname{tr}(w_B^{1,2})$. If for every word w' of length k + 1 we have $\operatorname{tr}(w_A') = \operatorname{tr}(w_B')$, then $\operatorname{tr}(w_A')^2 = \operatorname{tr}(w_B')^2$ as well.

Proof. Suppose a word w' of length n+1 can be written as, without loss of generality, say, Qw with w a word of length n. Since for any complex number u and v matrices $uQ_A + vw_A$ and $uQ_B + vw_B$ both have rank 2,

$$\operatorname{tr}(uQ_A + vw_A)^{1,2} = \operatorname{tr}(uQ_B + vw_B)^{1,2}$$

implies

$$\operatorname{tr}(uQ_A + vw_A)^i = \operatorname{tr}(uQ_B + vw_B)^i$$

for every integer i. One could check that

$$\operatorname{tr}(Q_A) = \operatorname{tr}(Q_B), \operatorname{tr}(Q^2)_A = \operatorname{tr}(Q^2)_B, \operatorname{tr}(w_A) = tr(w_B),$$

$$\operatorname{tr}(w_A)^2 = \operatorname{tr}(w_B)^2 \operatorname{tr}(Qw)_A = \operatorname{tr}(Qw)_B$$

indeed guarantee that $\operatorname{tr}(uQ_A + vw_A)^{1,2} = \operatorname{tr}(uQ_B + vw_B)^{1,2}$. Using the equality for i=4 and matching the coefficients of u^2v^2 one gets

$$\operatorname{tr}(QwQw)_A + \operatorname{tr}(Q^2w^2)_A = \operatorname{tr}(QwQw)_B + \operatorname{tr}(Q^2w^2)_B$$

Using the equality for i=3 and matching the coefficients of u^2v we get

$$\operatorname{tr}(Q^2 w)_A = \operatorname{tr}(Q^2 w)_B.$$

Now we consider $uQ_A^2 + w_A$ and $uQ_B^2 + w_B$. For the same reason as above,

$$tr(uQ_A^2 + w_A)^{1,2} = tr(uQ_B^2 + w_B)^{1,2}$$

implies

$$\operatorname{tr}(uQ_A^2 + w_A)^i = \operatorname{tr}(uQ_B^2 + w_B)^i$$

for every integer i. One could check that

$$tr(Q_A^2) = tr(Q_B^2), tr(Q_A^4) = tr(Q_B^4), tr(w_A) = tr(w_B), tr(w_A)^2 = tr(w_B)^2, tr(Q^2w)_A = tr(Q^2w)_B$$

imply $\operatorname{tr}(uQ_A^2 + w_A)^{1,2} = \operatorname{tr}(uQ_B^2 + w_B)^{1,2}$. Using the equality for i=3 and matching the coefficients of uv^2 one gets $\operatorname{tr}(Q^2w^2)_A = \operatorname{tr}(Q^2w^2)_B$. And thus we deduce $\operatorname{tr}(QwQw)_A = \operatorname{tr}(QwQw)_B$ as well. This is just $\operatorname{tr}(w'_A)^2 = \operatorname{tr}(w'_B)^2$.

We say that two words w and w' are cyclically equivalent if w' can be obtained from a cyclic permutation of w. For example, PQRS, QRSP, RSPQ and SPQR are all cyclically equivalent. Note that for two cyclically equivalent words w and w', $\operatorname{tr}(w_A)=\operatorname{tr}(w'_A)$ and $\operatorname{tr}(w_B)=\operatorname{tr}(w'_B)$. Thus when we check if $\operatorname{tr}(w_A)=\operatorname{tr}(w_B)$ we can always freely change w to any word w' that is cyclically equivalent to w.

Theorem 2.8 Let A and B be two n-by-n matrices both of rank 2. If $tr(w_A^{(i)}) = tr(w_B^{(i)})$ for $1 \le i \le 8$, then A is unitarily equivalent to B.

Proof. We proceed by induction on the length n of words. From Lemmas 2.6 and 2.7 we see that for words w of length 1 or 2, one has $tr(w_A)^{1,2} = tr(w_B)^{1,2}$. This proves our assertion for n = 1 and n = 2.

Now suppose for n = k, k+1, every word w of block n satisfies $\operatorname{tr}(w_A)^{1,2} = \operatorname{tr}(w_B)^{1,2}$. Let w' be any word of length k+2. If w' is not any of the forms $(SP)^n, (PS)^n, (QR)^n$ or $(RQ)^n$, then we have w cyclically equivalent to one of the following 12 cases:

(1)
$$w = PPK$$
 for some word K of length k

Consider for any complex numbers u and v matrices $uP_A + vK_A$ and $uP_B + vK_B$. Note that $uP_A + vK_A$ and $uP_B + vK_B$ are both of rank 2. Also,

$$\operatorname{tr}(uP_A + vK_A)^{1,2} = \operatorname{tr}(uP_B + vK_B)^{1,2}$$

implies

$$\operatorname{tr}(uP_A + vK_A)^i = \operatorname{tr}(uP_B + vK_B)^i$$

for every integer *i*. From $\operatorname{tr}(P_A) = \operatorname{tr}(P_B)$ we see that the coefficients of *u* of both sides are equal. Since *K* is of length *k*, by induction hypothesis we have $\operatorname{tr}(K_A) = \operatorname{tr}(K_B)$ and so the coefficients of *v* of both sides are equal. From $\operatorname{tr}(P_A)^2 = \operatorname{tr}(P_B)^2$ we see that the coefficients of u^2 of both sides are equal. Since *PK* is of length k + 1, by induction hypothesis we have $\operatorname{tr}(PK)_A = \operatorname{tr}(PK)_B$ and so the coefficients of *uv* of both sides are equal. Finally, since *K* is of length *k*, by the induction hypothesis we have $\operatorname{tr}(K_A)^2 = \operatorname{tr}(K_B)^2$ and so the coefficients of v^2 of both sides are equal. Hence we conclude that $\operatorname{tr}(uP_A + vK_A)^{1,2} = \operatorname{tr}(uP_B + vK_B)^{1,2}$. Using the equality for *i*=3 and matching the coefficients of uv^2 we get $\operatorname{tr}(PPK)_A = \operatorname{tr}(PPK)_B$.

(2) w = PRK for some word K of length k:

Consider for any complex numbers u and v the matrices $uP_A + vAK_A$ and $uP_B + vBK_B$. One could check indeed that $\operatorname{tr}(uP_A + vAK_A)^{1,2} = \operatorname{tr}(uP_B + vBK_B)^{1,2}$ and this implies $\operatorname{tr}(uP_A + vAK_A)^i = \operatorname{tr}(uP_B + vBK_B)^i$ for every integer i. Using the equality for i = 3 and matching the coefficients of u^2v one gets $\operatorname{tr}(PRK)_A = \operatorname{tr}(PRK)_B$.

(3) w = RPK for some word K of length k:

Matching the coefficients of uvt in $tr(uP_A + vR_A + tK_A)^3 = tr(uP_B + vR_B + tK_B)^3$ one gets $tr(PRK)_A + tr(RPK_A) = tr(PRK)_B + tr(RPK_B)$. Thus from (2) we deduce $tr(RPK_A) = tr(RPK_B)$.

(4) w = PQK for some word K of length k:

Note that $\operatorname{tr}(PQK)_A^* = \operatorname{tr}(K^*RP)_A = \operatorname{tr}(RPK^*)_A$. Applying (3) we obtain $\operatorname{tr}(PQK)_A = \operatorname{tr}(PQK)_B$.

(5) w = QPK for some word K of length k:

Matching the coefficients of uvt in $tr(uQ_A + vP_A + tK_A)^3 = tr(uQ_B + vP_B + tK_B)^3$, one gets $tr(QPK)_A + tr(PQK_A) = tr(QPK)_B + tr(PQK_B)$. Thus from (2) we deduce $tr(QPK_A) = tr(QPK_B)$.

(6) w = QQK for some word K of length k:

Consider for any complex numbers u and v the matrices $uQ_A + vK_A$ and $uQ_B + vK_B$. One could check indeed that $\operatorname{tr}(uQ_A + vK_A)^{1,2} = \operatorname{tr}(uQ_B + vK_B)^{1,2}$ and this implies $\operatorname{tr}(uQ_A + vK_A)^i = \operatorname{tr}(uQ_B + vK_B)^i$ for every integer i. Using the equality for i = 3 and matching the coefficients of u^2v one gets $\operatorname{tr}(QQK)_A = \operatorname{tr}(QQK)_B$.

(7) w = QSK for some word K of length k:

Consider for any complex numbers u and v the matrices $uQ_A + vAK_A$ and $uQ_B + vBK_B$. One could check indeed that $\operatorname{tr}(uQ_A + vAK_A)^{1,2} = \operatorname{tr}(uQ_B + vBK_B)^{1,2}$ and this implies $\operatorname{tr}(uQ_A + vAK_A)^i = \operatorname{tr}(uQ_B + vBK_B)^i$ for every integer i. Using the equality for i = 3 and matching the coefficients of u^2v one gets $\operatorname{tr}(QSK)_A = \operatorname{tr}(QSK)_B$.

(8) w = RRK for some word K of length k:

Consider for any complex numbers u and v the matrices $uR_A + vK_A$ and $uR_B + vK_B$. One could check indeed that $\operatorname{tr}(uR_A + vK_A)^{1,2} = \operatorname{tr}(uR_B + vK_B)^{1,2}$ and this implies $\operatorname{tr}(uR_A + vK_A)^i = \operatorname{tr}(uR_B + vK_B)^i$ for every integer i. Using the equality for i = 3 and matching the coefficients of $u^2 v$ one gets $tr(RRK)_A = tr(RRK)_B$.

(9) w = SQK for some word K of length k:

Matching the coefficients of uvt in $tr(uQ_A + vS_A + tK_A)^3 = tr(uQ_B + vS_B + tK_B)^3$ one gets $tr(QSK)_A + tr(SQK_A) = tr(QSK)_B + tr(SQK_B)$. Thus from (7) we deduce $tr(SQK_A) = tr(SQK_B)$.

(10) w = RSK for some word K of length k: Note that $\operatorname{tr}(RSK)_A^* = \operatorname{tr}(K^*SQ)_A = \operatorname{tr}(SQK^*)_A$. Applying (9) we obtain $\operatorname{tr}(RSK)_A = \operatorname{tr}(RSK)_B$.

(11) w = SRK for some word K of length k:

Matching the coefficients of uvt in $tr(uS_A + vR_A + tK_A)^3 = tr(uS_B + vR_B + tK_B)^3$ one gets $tr(RSK)_A + tr(SRK_A) = tr(RSK)_B + tr(SRK_B)$, Thus from (10) we deduce $tr(SRK_A) = tr(SRK_B)$.

(12) W = SSK for some word K of length k

Consider for any complex numbers u and v the matrices $uS_A + vK_A$ and $uS_B + vK_B$. One could check indeed that $\operatorname{tr}(uS_A + vK_A)^{1,2} = \operatorname{tr}(uS_B + vK_B)^{1,2}$ and this implies $\operatorname{tr}(uS_A + vK_A)^i = \operatorname{tr}(uS_B + vK_B)^i$ for every integer i. Using the equality for i = 3 and matching the coefficients of u^2v one gets $\operatorname{tr}(SSK)_A = \operatorname{tr}(SSK)_B$.

For $w' = (PS)^n$, we showed in Lemma 2.6 that $\operatorname{tr}(PS)_A = \operatorname{tr}(PS)_B$ and $\operatorname{tr}(PSPS)_A = \operatorname{tr}(PSPS)_B$. Since $(PS)_A$ and $(PS)_B$ are both of rank 2, this implies that $\operatorname{tr}(PS)_A^n = \operatorname{tr}(PS)_B^n$ for all integers n. Similarly, we have $\operatorname{tr}(QR)_A^n = \operatorname{tr}(QR)_B^n$ for all integers n. So we have proved that $\operatorname{tr}(w'_A) = \operatorname{tr}(w'_B)$ for every word w' of length k + 2. Finally, from Lemma 2.7 and the induction hypothesis that $\operatorname{tr}(w_A)^{1,2} = \operatorname{tr}(w_B)^{1,2}$ for every word w of length k + 1 we conclude that $\operatorname{tr}(w'_A)^2 = \operatorname{tr}(w'_B)^2$ for every w' of length k + 2. Thus we conclude by induction that $\operatorname{tr}(w_A) = \operatorname{tr}(w_B)$ for any w of length $n, n \geq 1$. So A is unitarily equivalent to B.

Now let us come back to the 3-by-3 cases. First we show that $w^{(7)}$ is redundant.

Lemma 2.9 Suppose A and B are two 3-by-3 matrices. If $\operatorname{tr}(w_A^{(i)}) = \operatorname{tr}(w_B^{(i)})$ for $1 \le i \le 6$, then $\operatorname{tr}(w_A^{(7)}) = \operatorname{tr}(w_B^{(7)})$

Proof. For any complex numbers u and v consider matrices $uA + vA^*$ and $uB + vB^*$. One could check that tr $(w_A^i) = tr(w_B^i)$ for $1 \le i \le 5$ implies

$$\operatorname{tr}(uA + vA^*)^{1,2,3} = \operatorname{tr}(uB + vB^*)^{1,2,3}.$$

Hence we have $\operatorname{tr}(uA + vA^*)^i = \operatorname{tr}(uB + vB^*)^i$ also for every integer *i*. Using the equality for i = 4 and matching the coefficients of u^2v^2 one gets

$$\operatorname{tr}(w_A^{(6)}) + \operatorname{tr}(w_A^{(7)}) = \operatorname{tr}(w_B^{(6)}) + \operatorname{tr}(w_B^{(7)}).$$

From tr $(w_A^{(6)}) = \operatorname{tr}(w_B^{(6)})$ we then deduce tr $(w_A^{(7)}) = \operatorname{tr}(w_B^{(7)})$

Lemma 2.10 Suppose A and B are two 3-by-3 matrices and tr $(w_A^{(i)}) = \operatorname{tr}(w_B^{(i)})$ for $1 \leq i \leq 6$ and i = 8. Then, for any complex number λ , $\operatorname{tr}(w_{(A-\lambda I_3)}^{(i)}) = \operatorname{tr}(w_{(B-\lambda I_3)}^{(i)})$ for $1 \leq i \leq 6$ and i = 8 as well.

Proof. The only case we really need to verify is $\operatorname{tr}(w_{(A-\lambda I_3)}^{(8)}) = \operatorname{tr}(w_{(B-\lambda I_3)}^{(8)})$. This in turn is equivalent to checking that if $\operatorname{tr}(A^{*2}AA^*A) = \operatorname{tr}(B^{*2}BB^*B)$. Using

$$\operatorname{tr}(uA + vA^*)^7 = \operatorname{tr}(uB + vB^*)^7$$

and matching the coefficients of u^3v^4 one gets indeed

$$\operatorname{tr}(A^{*2}AA^{*}A) = \operatorname{tr}(B^{*2}BB^{*}B).$$

Theorem 2.11 If A and B are two 3-by-3 matrices such that $\operatorname{tr}(w_A^{(i)}) = \operatorname{tr}(w_B^{(i)})$ for $1 \leq i \leq 6$ and i = 8, then A is unitarily equivalent to B.

Proof. First from $\operatorname{tr}(A)^{1,2,3} = \operatorname{tr}(B)^{1,2,3}$ we conclude that A and B have the same eigenvalues. Subtracting a common eigenvalue λ we get $A' = A - \lambda I_3$ and $B' = B - \lambda I_3$ both having rank 2. From Lemma 2.5 we have $\operatorname{tr}(w_{A'}^{(i)}) = \operatorname{tr}(w_{B'}^{(i)})$ for $1 \leq i \leq 6$ and i = 8. From Lemma 2.4 we see $\operatorname{tr}(w_{A'}^{(7)}) = \operatorname{tr}(w_B^{(7)})$ as well. From Theorem 2.1 we conclude that A' is unitarily equivalent to B'. Thus A is also unitarily equivalent to B.

3 Matrices with Eigenvectors Not Orthogonal

We prove in this chapter the most general result in this paper, namely, for any matrix whose eigenvectors are not orthogonal a set with $n^4 + 1$ words suffices to determine it up to unitary equivalence. For example,

$$A = \begin{pmatrix} a & b \\ 0 & c \end{pmatrix}, \ b \neq 0$$

is such a matrix.

Theorem 3.1 Let A be an n-by-n matrix such that none of its eigenvectors are orthogonal. Then $\{A^{*i}A^j : 0 \leq i, j < n\}$ is a linearly independent set and spans $Alg(A, A^*)$.

Proof. Suppose that

$$f(x) = (x - \lambda_1)^{p_1} (x - \lambda_2)^{p_2} \dots (x - \lambda_m)^{p_m}$$

is the characteristic polynomial of A with $\lambda_1, \lambda_2, ..., \lambda_m$ the distinct eigenvalues of A. Since none of the eigenvectors of A are orthogonal, we have in particular ker $(A - \lambda_i I_n) = 1$ for all i, on $1 \le i \le m$. This implies that f is also the minimal polynomial of A. Suppose on the contrary that $\{A^{*i}A^j : 0 \le i, j < n\}$ is linearly dependent. Then there exist polynomials $g_0, g_1, \ldots, g_{n-1}$ of degree < n, not all zero, such that

$$g_{n-1}(A^*)A^{n-1} + \dots + g_0(A^*) = 0.$$
 (3.1)

Let $h_{n-1} = \gcd(\overline{f}, g_{n-1})$ and write $g_{n-1} = qh_{n-1}$ with q also a polynomial. Then $q(A^*)$ is invertible with its inverse also a polynomial in A^* . Thus we can multiply both sides of (3.1) with $j(A^*)^{-1}$ and obtain

$$h_{n-1}(A^*)A^{n-1} + \dots + h_0(A^*) = 0$$
(3.2)

with $h_0, h_1, \ldots, h_{n-1}$ still polynomials of degree $\langle n$. Since h_{n-1} divides \overline{f} , there is some r such that $h_{n-1}(x)$ divides $(x - \overline{\lambda_1})^{p_1} \cdots (x - \overline{\lambda_r})^{p_r-1} \cdots (x - \overline{\lambda_m})^{p_m}$. We multiply both sides of (3.2) with $(A^* - \overline{\lambda_1})^{p_1} \cdots (A^* - \overline{\lambda_r})^{p_r-1} \cdots (A^* - \overline{\lambda_m})^{p_m}$ and obtain for some nonzero polynomial h the relation

$$(A^* - \overline{\lambda_1})^{p_1} \cdots (A^* - \overline{\lambda_r})^{p_r - 1} \cdots (A^* - \overline{\lambda_m})^{p_m} h(A) = 0.$$
(3.3)

Similarly, take $l=\gcd(f,h)$ and denote h = lu. Then u(A) is invertible with its inverse also a polynomial in A. So we multiply both sides of (3.3) with $u(A)^{-1}$ and obtain

$$(A^* - \overline{\lambda_1})^{p_1} \cdots (A^* - \overline{\lambda_r})^{p_r - 1} \cdots (A^* - \overline{\lambda_m})^{p_m} l(A) = 0.$$
(3.4)

Since *l* divides *f* and is of degree < n, there is some *s* such that *l* divides $(x - \lambda_1)^{p_1} \cdots (x - \lambda_s)^{p_s - 1} \cdots (x - \lambda_m)^{p_m}$. So we multiply both sides of (3.4) with $(A - \lambda_1)^{p_1} \cdots (A - \lambda_s)^{p_s - 1} \cdots (A - \lambda_m)^{p_m}$ and obtain

$$(A^* - \overline{\lambda_1})^{p_1} \cdots (A^* - \overline{\lambda_r})^{p_r - 1} \cdots (A^* - \overline{\lambda_m})^{p_m} (A - \lambda_1)^{p_1} \cdots (A - \lambda_s)^{p_s - 1} \cdots (A - \lambda_m)^{p_m} = 0. \quad (3.5)$$

Since f is the minimal polynomial of A, there exists some nonzero v in \mathbb{C}^n which is not in ker $(A - \lambda_r)^{p_r-1}$ but in ker $(A - \lambda_r)^{p_r}$. So $(A - \lambda_r)^{p_r-1}v$ is nonzero and belongs in ker $(A - \lambda_r)$. Let v_r in \mathbb{C}^n be an eigenvector of A corresponding to the eigenvalue λ_r . Then for some nonzero complex number a we have $(A - \lambda_r)^{p_r-1}v = av_r$. Similarly, for some nonzero w in \mathbb{C}^n and nonzero complex number b, we have $(A - \lambda_s)^{p_s-1}w = bv_s$, where v_s is an eigenvector of A corresponding to the eigenvalue λ_s . From (3.5), we must have

$$\left\langle (A - \lambda_1)^{p_1} ... (A - \lambda_r)^{p_r - 1} ... (A - \lambda_m)^{p_m} v, (A - \lambda_1)^{p_1} ... (A - \lambda_s)^{p_s - 1} ... (A - \lambda_m)^{p_m} w \right\rangle = 0.$$

This in turn implies

$$\left\langle \left(\prod_{1 \le i \le m; i \ne r} (A - \lambda_i)^{p_i}\right) v_r, \left(\prod_{1 \le i \le m; i \ne s} (A - \lambda_i)^{p_i}\right) v_s \right\rangle = 0$$

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and hence

So we find a pair of eigenvectors of A which are orthogonal, contradicting to our assumption on A. Thus the set $\{A^{*i}A^j: 0 \le i, j < n\}$ must be linearly independent. Since this set contains n^2 elements and $\text{Alg}(A, A^*)$ is of dimension at most n^2 , this implies that $\{A^{*i}A^j: 0 \le i, j < n\}$ spans $\text{Alg}(A, A^*)$.

Suppose now that A and B are two n-by-n matrices with none of the eigenvectors of A orthogonal. By Theorem 3.1, for any integers p and q we can write $A^p A^{*q}$ as a linear combination of $A^{*i}A^j, 0 \leq i, j < n$. The next theorem shows that if

$$\operatorname{tr}(A^n) = \operatorname{tr}(B^n)$$
 and $\operatorname{tr}(A^{*i}A^jA^{*k}A^l) = \operatorname{tr}(B^{*i}B^jB^{*k}B^l), \ 0 \le i, j, k, l < n_j$

then we can also write $B^p B^{*q}$ as a linear combination of $\{B^{*i}B^j : 0 \le i, j < n\}$ with the same coefficients.

Theorem 3.2 Let A and B be two n-by-n matrices and p and q be some integers. Suppose that there exists a set of complex numbers $\{a_{ij} : 0 \leq i, j < n\}$ such that $A^p A^{*q} = \sum_{i,j=0}^{n-1} a_{ij} A^{*i} A^j$. If

$$tr(A^n) = tr(B^n)$$
 and $tr(A^{*i}A^jA^{*k}A^l) = tr(B^{*i}B^jB^{*k}B^l), \ 0 \le i, j, k, l < n_j$

then we also have $B^p B^{*q} = \sum_{i,j=0}^{n-1} a_{ij} B^{*i} B^j$.

Proof. From $A^p A^{*q} = \sum_{i,j=0}^{n-1} a_{ij} A^{*i} A^j$, we have

$$\operatorname{tr}\left((A^{p}A^{*q} - \sum_{i,j=0}^{n-1} a_{ij}A^{*i}A^{j})^{*}(A^{p}A^{*q} - \sum_{i,j=0}^{n-1} a_{ij}A^{*i}A^{j})\right) = 0.$$

This is the same as

$$\operatorname{tr}(A^{q}A^{*p}A^{p}A^{*q}) - \operatorname{tr}(\sum_{i,j=0}^{n-1} a_{ij}A^{q}A^{*p}A^{*i}A^{j}) - \operatorname{tr}(\sum_{i,j=0}^{n-1} a_{ij}A^{*i}A^{j}A^{q}A^{*p}) - \operatorname{tr}(\sum_{i,j,k,l=0}^{n-1} a_{ij}a_{kl}A^{*j}A^{i}A^{*k}A^{l}) = 0. \quad (3.6)$$

Since $\operatorname{tr}(A^i) = \operatorname{tr}(B^i), 1 \leq i \leq n, A$ and B have equal characteristic polynomials and thus $\operatorname{tr}(A^{*i}A^jA^{*k}A^l) = \operatorname{tr}(B^{*i}B^jB^{*k}B^l), 0 \leq i, j, k, l < n$ implies $\operatorname{tr}(A^{*i}A^jA^{*k}A^l) = \operatorname{tr}(B^{*i}B^jB^{*k}B^l)$ for all nonnegative integers i, j, k and l. Substituting this back into (3.6), we obtain

$$\operatorname{tr}(B^{q}B^{*p}B^{p}B^{*q}) - \operatorname{tr}(\sum_{i,j=0}^{n-1} a_{ij}B^{q}B^{*p}B^{*i}B^{j}) - \operatorname{tr}(\sum_{i,j=0}^{n-1} a_{ij}B^{*i}B^{j}B^{q}B^{*p}) - \operatorname{tr}(\sum_{i,j,k,l=0}^{n-1} a_{ij}a_{kl}B^{*j}B^{i}B^{*k}B^{l}) = 0 \quad (3.7)$$

as well. So we have

$$\operatorname{tr}(B^{p}B^{*q} - \sum_{i,j=0}^{n-1} a_{ij}B^{*i}B^{j})^{*}(B^{p}B^{*q} - \sum_{i,j=0}^{n-1} a_{ij}B^{*i}B^{j}) = 0.$$

Thus $B^p B^{*q} - \sum_{i,j=0}^{n-1} a_{ij} B^{*i} B^j = 0.$

Now we are ready to prove the main result.

Theorem 3.3 Let A and B be two n-by-n matrices. Assume that no pair of the eigenvectors of A are orthogonal to each other. If $\operatorname{tr}(A^{*i}A^jA^{*k}A^l) = \operatorname{tr}(B^{*i}B^jB^{*k}B^l)$ for $0 \leq i, j, k, l < n$ and $\operatorname{tr}(A^n) = \operatorname{tr}(B^n)$, then A is unitarily equivalent to B.

Proof. We are going to show that for every word w(x, y), $tr(w(A, A^*)) = tr(w(B, B^*))$ and then apply Specht's theorem to conclude that A is unitarily equivalent to B. Using the property that tr(AB) = tr(BA), it suffices to consider only words w of the form $w(x, y) = y^{i_1} x^{j_1} \cdots y^{i_n} x^{j_n}$.

We proceed by induction on the length n. For n=1, this is already assumed. Suppose that the assertion is true for n = k. Consider the case n = k + 1. Then

$$\operatorname{tr}(w(A, A^*)) = \operatorname{tr}(A^{*i_1}A^{j_1} \cdots A^{*i_k}(A^{j_k}A^{*i_{k+1}})A^{j_{k+1}})$$
$$= \operatorname{tr}\left(\sum_{p,q=0}^{n-1} a_{pq}A^{*i_1}A^{j_1} \cdots A^{*(i_k+p)}A^q\right) = \sum_{p,q=0}^{n-1} a_{pq}\operatorname{tr}\left(A^{*i_1}A^{j_1} \cdots A^{*(i_k+p)}A^q\right). \quad (3.8)$$

Since $A^{*i_1}A^{j_1}\cdots A^{*(i_k+p)}A^q$ is of length k, by the induction hypothesis we have

$$\operatorname{tr}(A^{*i_1}A^{j_1}\cdots A^{*(i_k+p)}A^q) = \operatorname{tr}(B^{*i_1}B^{j_1}\cdots B^{*(i_k+p)}B^q).$$
(3.9)

Substituting this back into (3.8), we have

$$\operatorname{tr}(w(A, A^*)) = \sum_{p,q=0}^{n-1} a_{pq} \operatorname{tr} \left(A^{*i_1} A^{j_1} \cdots A^{*(i_k+p)} A^q \right)$$
$$= \sum_{p,q=0}^{n-1} a_{pq} \operatorname{tr} \left(B^{*i_1} B^{j_1} \cdots B^{*(i_k+p)} B^q \right)$$
$$= \operatorname{tr} \left(\sum_{p,q=0}^{n-1} a_{pq} B^{*i_1} B^{j_1} \cdots B^{*(i_k+p)} B^q \right)$$
$$= \operatorname{tr}(B^{*i_1} B^{j_1} \cdots B^{*i_k} (B^{j_k} B^{*i_{k+1}}) B^{j_{k+1}}) = \operatorname{tr}(w(B, B^*)).$$

Thus by the mathematical induction we conclude that $tr(w(A, A^*)) = tr(w(B, B^*))$ for every word w. So A is unitarily equivalent to B.

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