# 國立交通大學資訊科學與工程研究所

## 碩士論 文



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## 中華民國九十五年六月

#### 條件式容錯超立方體下的邊泛迴圈之研究 Edge-bipancyclicity of conditional faulty hypercubes

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國立交通大學 資訊科學與工程研究所 碩士論文

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A Thesis Submitted to Institute of Computer Science and Engineering College of Computer Science National Chiao Tung University in partial Fulfillment of the Requirements for the Degree of

Master

in

Institute of Computer Science and Engineering

June 2006

Hsinchu, Taiwan, Republic of China

中華民國九十五年六月

## 條件式容錯超立方體下的邊泛迴圈 之研究

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摘要

Xu et al.先前曾提過相關的論文研究,證明了當N大於四時,在n維度的超力 方體下若壞邊的數量小餘n-1時,通過任意指定的一個邊,仍可找到從 6 到 2<sup>n</sup>這 樣各種長度的迴圈,而限制條件在於並非所有壞邊皆集中在同一個點,意即每個 點仍然存在有兩個好邊。

在這篇論文中,我們在相似的條件下,若壞邊不集中在同一個點,則壞邊個 數能夠增加到 2n-5 個,且通過任意指定的一個邊,仍可找到長度從 6 到 2<sup>m的</sup>各種 的迴圈。此外我們仍證明了,當壞邊達到 2n-4 時,如此是無法被證明的,所以 我們的結論是最佳的結果。

關鍵字:迴圈,邊泛迴圈,條件式容錯超力方體,容錯

## Edge-bipancyclicity of conditional faulty hypercubes

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#### Abstract

Xu et al. showed that for any set of faulty edges F of an n-dimensional hypercube  $Q_n$  with  $|F| \leq n-1$ , each edge of  $Q_n$ -F lies on a cycle of every even length from 6 to  $2^n$ ,  $n \geq 4$ , provided not all edges in F are incident with the same vertex. In this paper, we find that under similar condition, the number of faulty edges can be much greater and the same result still holds. More precisely, we show that, for up to |F|=2n-5 faulty edges, each edge of the faulty hypercube  $Q_n$ -F lies on a cycle of every even length from 6 to  $2^n$  with each vertex having at least two healthy edges adjacent to it, for  $n \geq 3$ . Moreover, this result is optimal in the sense that the result can not be guaranteed, if there are 2n-4 faulty edges.

**Keywords:** cycles, Pancyclic, Conditional fault, Hypercube, Fault-tolerant

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## Chapter 1

## Introduction

The ring embedding problem, which deals with all the possible lengths of the cycles in a given graph, is investigated in a lot of interconnection networks [2, 3, 4]. If a graph contains cycles of all lengths, it is called pancyclic [7]. Bipancyclicity is essentially a restriction of the concept of pancyclicity to cycle of even lengths. A bipartite graph is *vertex-bipancyclic* [6] if every vertex lies on a cycle of every even length from 4 to |V(G)|. Similarly, a bipartite graph is *edge-bipancyclic* if every edge lies on a cycle of every even length from 4 to |V(G)|. A bipartite graph is *k-edge-fault-tolerant edge-bipancyclic* if G - F remains edge-bipancyclic for any set of faulty edges  $F \subset E(G)$  with  $|F| \leq k$ . A path P is a sequence of adjacent vertices, written as  $\langle v_0, v_1, ..., v_m \rangle$ . The *length* of a path P, denoted by l(P), is the number of edges in P. A hamiltonian cycle is a graph with a spanning cycle. In addition we call e a *healthy edge* when e is fault-free in a graph.

Chan and Lee [1] considered an injured *n*-dimensional hypercube where each vertex is incident with at least two healthy edges, and proved that it still contains a hamiltonian cycle even it has (2n - 5) edge faults. Tsai [8] proved that such injured hypercube  $Q_n$ contains a cycle of every even length from 4 to  $2^n$ , even if it has up to (2n - 5) edge faults. Recently, Xu et al. [9] showed that for any set of faulty edges F od  $Q^n$  with  $|F| \le n - 1$ , each edge of  $Q_n - F$  lies on a cycle of every even length from 6 to  $2^n$ ,  $n \ge 4$ , provided not all faulty edges are incident with the same vertex. We observe that not all faulty edges a re incident with the same vertex is equivalent to stating that each vertex has at least two healthy edges adjacent to it, if  $|F| \le n - 1$ . In this paper, we consider a set of faulty edges satisfies the condition that each vertex of  $Q_n - F$  is incident with at least two healthy edges. Such a set of faulty edges F is called a set of conditional faulty edges and  $Q_n - F$  is called a conditional faulty hypercube. We find that under this condition, the number of faulty edges can be much greater and the same result still holds. We show that, for up to |F| = 2n - 5 conditional faulty edges, each edge of a faulty hypercube  $Q_n - F$  lies on a cycle of every even length from 6 to  $2^n$  with each vertex having at least two healthy edges adjacent to it, for  $n \ge 3$ . We observe that, if |F| < 2n - 5, we may arbitrarily delete some more edges to make a faulty edge set  $F' \supseteq F$  and |F'| = 2n - 5. If our result holds for F', it holds for F. From now on, we shall assume |F| = 2n - 5.

The above result is optimal in the sense that result can not be guaranteed, if there are 2n - 4 conditional faulty edges. For example, take a cycle of length four in  $Q_n$ , let  $\langle u_1, u_2, u_3, u_4 \rangle$  be the consecutive vertices on this cycle. Suppose that all the (n-2) edges incident to vertex  $u_1$  (respectively vertex  $u_3$ ) are faulty except those two edges on the four cycle are healthy. There are 2(n-2) conditional faulty edges. (see Fig.1.1) Then there does not exist a hamiltonian cycle in this faulty  $Q_n$ , for  $n \geq 3$ .

We now give a formal definition of a hypercube. An *n*-dimensional hypercube is denoted by  $Q_n$  with the vertex set  $V(Q_n)$  and the edge set  $E(Q_n)$ . Each vertex u of  $Q_n$ can be distinctly labeled by a *n*-bit binary strings,  $u = u_{n-1}u_{n-2}...u_1u_0$ . There is an edge

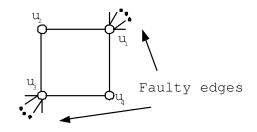


Figure 1.1: Illustration for the  $Q_n$  with (2n-4) edge fault.

between two vertices if and only if their binary labels differ in exactly one bit position. Let u and v be two adjacent vertices. If the binary labels of u and v differ in *i*th position, then the edge between them is said to be in dimension i and the edge (u, v) is called an *i*th dimension edge. Let i be a fixed position, we use  $Q_{n-1}^0$  to denote the subgraph of  $Q_n$  induced by  $\{u \in V(Q_n) | u_i = 0\}$  and  $Q_{n-1}^1$  to denote the subgraph of  $Q_n$  induced by  $\{u \in V(Q_n) | u_i = 1\}$ . We say that  $Q_n$  is decomposed into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by dimension i, and  $Q_{n-1}^0$  and  $Q_{n-1}^1$  are (n-1)-dimensional subcube of  $Q_n$  induced by the vertices with the *i*th bit position being  $\theta$  and 1 respectively.  $Q_{n-1}^0$  and  $Q_{n-1}^1$ , denoted by  $u^{(1)}$ , such that  $(u, u^{(1)}) \in E(Q_n)$ . Conversely, for each  $u \in V(Q_{n-1}^1)$ , there is one vertex in  $Q_{n-1}^0$ , denoted by  $u^{(0)}$ , such that  $(u, u^{(0)}) \in E(Q_n)$ . Let  $D_i$  be the set of all edges with one end in  $Q_{n-1}^0$  and  $Q_{n-1}^1$  and  $Q_{n-1}^1$ . These edges are called crossing edges in the *i*th dimension edges. Consequently,  $|D_i| = 2^{n-1}$  for all  $0 \le i \le n-1$ .

## Chapter 2

## Some Preliminaries

To prove our main theorem, we need some preliminary results.

Lemma 1 [5]  $Q_n$  is edge-bipancyclic, and is (n-2)-edge-fault-tolerant edge-bipancyclic, for  $n \ge 3$ .

**Lemma 2** [9] Each edge of  $Q_4 - F$  lies on a cycle of every even length from 6 to  $2^n = 16$ for any  $F \subset E(Q_4)$  with |F| = 3, provided not all the faulty edges in F are incident with the same vertex.

**Lemma 3** [9] Any two edges in  $Q_n$  are included in a hamiltonian cycle, for  $n \ge 2$ .

**Proof.** We prove the lemma by induction on  $n \ge 2$ . Obviously, the lemma is true for n = 2. Assume that the lemma is true for every k with  $2 \le k < n$ . Let e and e' be two edges in  $Q_n$  and express  $Q_n = L_k \odot R_k$  such that none of e and e' is k-dimensional. Without loss of generality, we may assume  $e \in L_k$ . Furthermore, we can suppose that e' is in  $L_k$ , otherwise consider  $e'_L$  instead of e'. By the induction hypothesis, there exists a

Hamiltonian cycle C containing e and e' in  $L_k$ . Let  $u_L v_L$  be an edge on C different from e and e'. The corresponding C' is a Hamiltonian cycle in  $R_k$  containing  $u_R v_R$ ,  $e_R$  and  $e'_R$ . Let  $P = C - u_L v_L$  and  $P' = C' - u_R v_R$ . Then  $P + u_L v_L + P' + u_R v_R$  is a Hamiltonian cycle in  $Q_n$  containing e and e'.

The above lemma can be improved; In addition to the hamiltonian cycle, there are cycles of smaller lengthes passing through these two edges. We have the following lemma.

**Lemma 4** Any two edges in  $Q_n$  are included in a cycle of length  $2^n$ ,  $2^n - 2$  and  $2^n - 4$  respectively, for  $n \ge 3$ .

**Proof.** Let  $e_1$ ,  $e_2$  be two arbitrary edges in  $Q_n$ . Since  $n \ge 3$ , we can decompose  $Q_n$  into two subcubes  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by some dimension i, such that  $e_1$  and  $e_2$  are not crossing edges. We then consider the following two cases:

**Case 1:** Both edges  $e_1$  and  $e_2$  are in the same subcube  $Q_{n-1}^i$ . Without loss of generality, we assume that  $\{e_1, e_2\} \subseteq E(Q_{n-1}^0)$ . By Lemma 3, there exists a hamiltonian cycle  $C_0$ going through  $e_1$  and  $e_2$  in  $Q_{n-1}^0$ , and the length of  $C_0$  is  $2^{n-1}$ . Since  $n \geq 3$ , there is a third edge (u, v) other than  $e_1$  and  $e_2$  on cycle  $C_0$ . We write  $C_0$  as  $\langle u, P_0, v, u \rangle$ . It follows from the definition of the hypercubes,  $(u^{(1)}, v^{(1)})$  is an edge in  $Q_{n-1}^1$ . By Lemma 1, in  $Q_{n-1}^1$ , there exists a cycle  $C_1$  of every even length  $4 \leq l(C_1) \leq 2^{n-1}$  going through  $(u^{(1)}, v^{(1)})$ . We write  $C_1$  as  $\langle u^{(1)}, P_1, v^{(1)}, u^{(1)} \rangle$ . Thus,  $\langle u, P_0, v, v^{(1)}, P_1, u^{(1)}, u \rangle$  can form a cycle of length  $2^n$ ,  $2^n - 2$  or  $2^n - 4$  respectively passing through edges  $e_1$  and  $e_2$  in  $Q_n$ , if we adjust the length of  $P_1$  properly.

**Case 2**:  $e_1$  and  $e_2$  are in different subcubes. Without loss of generality, we assume that  $e_1 \in E(Q_{n-1}^0), e_2 \in E(Q_{n-1}^1)$ . By Lemma 1, there exists a cycle  $C_0$  of every even length,  $4 \leq l(C_0) \leq 2^{n-1}$  going through  $e_1$  in  $Q_{n-1}^0$ . Since  $n \geq 3$ , we can choose an edge (u, v) on cycle  $C_0$  and by definition,  $(u^{(1)}, v^{(1)})$  is an edge in  $Q_{n-1}^1$  such that  $(u, v) \neq e_1$ , and  $(u^{(1)}, v^{(1)}) \neq e_2$ . We write  $C_0$  as  $\langle u, P_0, v, u \rangle$ . By Lemma 3, there exists a cycle  $C_1$  of length  $2^{n-1}$  going through  $e_2$  and  $(u^{(1)}, v^{(1)})$ . We write  $C_1$  as  $\langle u^{(1)}, P_1, v^{(1)}, u^{(1)} \rangle$ . Thus, we can adjust the length of  $P_0$  properly such that it produces a cycle  $\langle u, P_0, v, v^{(1)}, P_1, u^{(1)}, u \rangle$  of length  $2^n$ ,  $2^n - 2$  or  $2^n - 4$  respectively going through edges  $e_1$ ,  $e_2$  in  $Q_n$ .

This proves the lemma.

**Lemma 5** Let  $Q_n$  be an n-dimensional hypercube,  $n \ge 2$ , and let  $e_1$  and  $e_2$  be two edges in the same dimension *i*. Then there exists another dimension  $j \ne i$  such that decomposing  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by dimension *j*, we have (1) neither  $e_1$  nor  $e_2$  is crossing edges, (2) not  $e_1$  and  $e_2$  are in the same subcube.

**Proof.** Let  $e_1 = (a, b)$  and  $e_2 = (s, t)$  be two edges in the same dimension *i*. Let  $a = a_n \dots a_i \dots a_1$  and  $s = s_n \dots s_i \dots s_1$ . Then  $b = b_n \dots \overline{b_i} \dots b_1$  and  $t = t_n \dots \overline{t_i} \dots t_1$ . Since  $e_1 \neq e_2$  and  $n \geq 2$ , there exists another dimension  $j \neq i$ , such that  $a_j \neq s_j$ . We decompose  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by dimension j. Then,  $e_1$  and  $e_2$  are not crossing edges and in the different subcubes.

In some special cases, Lemma 4 still holds even if there is one faulty edge. We need the following lemma later.

**Lemma 6** Let  $Q_n$  be an n-dimensional hypercube and  $e_0$  be a faulty edge in the *i*th dimension. Any two healthy edges  $e_1$  and  $e_2$  in the *i*th dimension are included in a cycle of length  $2^n$ ,  $2^n - 2$  and  $2^n - 4$  respectively, for  $n \ge 4$ . **Proof.** Since  $e_0$ ,  $e_1$  and  $e_2$  are all in the *i*th demension, and  $n \ge 4$ . By Lemma 5, we can choose a dimension *j* different from *i* such that  $e_1$  and  $e_2$  are in different subcubes. Without loss of generality, we may assume that  $e_1$  is in  $Q_{n-1}^0$  and  $e_2$  is in  $Q_{n-1}^1$  and  $e_0$  is in  $Q_{n-1}^0$ . By Lemma 1, in  $Q_{n-1}^0 - \{e_0\}$ , there exists a cycle  $C_0$  of every length  $4 \le l(C_0) \le 2^{n-1}$  going through  $e_1$ . Since  $n \ge 4$ , we can choose an edge (u, v) on cycle  $C_0$  such that  $(u, v) \ne e_1$ , and  $(u^{(1)}, v^{(1)}) \ne e_2$ . We write  $C_0$  as  $\langle u, P_0, v, u \rangle$ . By Lemma 3, in  $Q_{n-1}^1$ , there exists a hamiltonian cycle  $C_1$  going through  $e_2$  and  $(u^{(1)}, v^{(1)})$ . We write  $C_1$  as  $\langle u^{(1)}, P_1, v^{(1)}, u^{(1)} \rangle$ . Thus,  $\langle u, P_0, v, v^{(1)}, P_1, u^{(1)}, u \rangle$  is a cycle of length  $2^n$ ,  $2^n - 2$  or  $2^n - 4$  respectively, if we adjust the length of  $P_0$  properly.

Let F be a set of faulty edges of  $Q_n$ . Suppose that we decompose  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by dimension j, and let  $F_L = F \cap E(Q_{n-1}^0)$ ,  $F_R = F \cap E(Q_{n-1}^1)$ . Suppose that F is a set of conditional faulty edges of  $Q_n$ . If we arbitrarily decompose  $Q_n$  into  $Q_{n-1}^0$ and  $Q_{n-1}^1$  by a dimension,  $F_L$  and  $F_R$  may not be conditional faulty edges in  $Q_{n-1}^0$  and  $Q_{n-1}^1$  respectively. However, we will show that it is always possible to find some suitable dimension such that decomposing by this dimension, both  $F_L$  and  $F_R$  are conditional faulty sets in  $Q_{n-1}^0$  and  $Q_{n-1}^1$  respectively.

**Lemma 7** Consider the n-dimensional hypercube  $Q_n$ , for  $n \ge 4$ . Let F be a set of conditional faulty edges with |F| = 2n - 5. There are at most two vertices in  $Q_n$  incident with (n-2) faulty edges.

**Proof.** If there are three vertices in  $Q_n$  incident with (n-2) faulty edges, the number of faulty edge F is at least 3n - 8. However, (3n - 8) > (2n - 5) for all  $n \ge 4$  which is a contradiction.

**Lemma 8** Consider the n-dimensional hypercube  $Q_n$ ,  $n \ge 4$ . Let F be a set of conditional faulty edges with |F| = 2n - 5. If there are two vertices x and y both incident with n-2faulty edges, then x and y are adjacent in  $Q_n$  and the edge (x,y) is a faulty edge. Suppose that (x,y) is in dimension j. Then decomposing  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by dimension j, both  $F_L$  and  $F_R$  are sets of conditional faulty edges in  $Q_{n-1}^0$  and  $Q_{n-1}^1$  respectively. Moreover,  $|F_L| \le 2n - 6$  and  $|F_R| \le 2n - 6$ .

**Proof.** If there are two vertices x and y in  $Q_n$  incident with (n-2) faulty edges, then these two vertices are connected by a faulty edge. Otherwise, |F| = 2(n-2) = 2n-4 > 2n-5which is a contradiction. Suppose the edge (x, y) is in dimension j, we decompose  $Q_n$ into two subcubes. It is clearly that each vertex in  $Q_{n-1}^0$  and  $Q_{n-1}^1$  is still incident with at least two healthy edges, and  $F_L$  and  $F_R$  are both conditional faulty edges in  $Q_{n-1}^0$  and  $Q_{n-1}^1$  respectively. Then,  $|F_L| = |F_R| = n-3 \le 2n-6$ , for  $n \ge 4$ .

**Lemma 9** Consider an n-dimensional hypercube  $Q_n$ , for  $n \ge 4$ . Let F be a set of conditional faulty edges with |F| = 2n - 5. Suppose that there exists exactly one vertex xhaving (n-2) faulty edges indicident with it. Since  $n - 2 \ge 2$ , let  $e_1$  and  $e_2$  be two faulty edges incident with x, and let  $e_1$  and  $e_2$  be jth and kth dimension edges respectively. Then decomposing  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by either one of these two dimensions j and k,  $F_L$ and  $F_R$  are still sets of conditonal faulty edges in  $Q_{n-1}^0$  and  $Q_{n-1}^1$  respectively. Moreover,  $|F_L| \le 2n - 6$  and  $|F_R| \le 2n - 6$ .

**Proof.** If there exists only one vertex x having (n-2) faulty edges incident with it, there are at least two faulty edges  $e_1$  and  $e_2$  incident with it, since  $n \ge 4$ . Obviously, these two faulty edges are in different dimensions. Without loss of generality, we may assume

that  $e_1$  is in dimension j and  $e_2$  is in dimension k, for  $j \neq k$ . We can decompose  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by either jth or kth dimension, and either  $e_1$  or  $e_2$  is a crossing edge. Therefore, each vertex in these two subcubes is incident with at least two healthy edges and  $|F_L| \leq 2n - 6$  and  $|F_R| \leq 2n - 6$ .

**Lemma 10** Let  $Q_n$  be an n-dimensional hypercube, F be a set of faulty edges with  $|F| \ge 2$ , and e be a healthy edge,  $n \ge 2$ . Then there exists a dimension j, decomposing  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by this dimension, such that e is not a crossing edge and not all the faulty edges are in the same subcube.

**Proof.** Suppose that e = (u, v) is in dimension *i*. If there is a faulty edge *f* not in dimension *i*, say in dimension *j*. We decompose  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by dimension *j*. Then *f* is a crossing edge but *e* is not, and all the faulty edges are not in the same subcube. Otherwise, all the faulty edges are in the same dimension *i* as *e* is in. We now choose any two faulty edges  $f_1$  and  $f_2$  in *F*. By Lemma 5,  $Q_n$  can be decomposed into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by some dimension  $j \neq i$  such that edges  $f_1$  and  $f_2$  are not in the same subcube, and  $e_0$  is not a crossing edge.

## Chapter 3

## Main theorem

We now prove our main result.

**Theorem 1** Let  $Q_n$  be an n-dimensional hypercube, and F be a set of conditional faulty edges with  $|F| \le 2n - 5$ . Then each edge of the conditional faulty hypercube  $Q_n - F$  lies on a cycle of every even length from 6 to  $2^n$ , for  $n \ge 3$ .

**Proof.** We prove this lemma by induction on n. For n = 3, since 2n - 5 = n - 2, by Lemma 1, the result is true. For n = 4, 2n - 5 = n - 1, by Lemma 2, the result holds. Assume the lemma holds for n - 1, for some  $n \ge 5$ , we shall show that it is true for n.

As we mentioned before, we may assume |F| = 2n - 5. Let e = (u, v) be an edge in  $Q_n - F$ . We shall find a cycle of every even length from 6 to  $2^n$  passing through e in  $Q_n - F$ . Assume that e is an *i*th dimension edge,  $e \in D_i$ , for some  $i \in \{1, 2, ..., n\}$ . The proof is divided into three major cases:

**Case 1**: There are two vertices x and y in  $Q_n$  incident with (n-2) faulty edges. By Lemma 8, (x, y) is an edge in  $Q_n$  and is a faulty edge. We denote this edge by  $e_f$ . Suppose

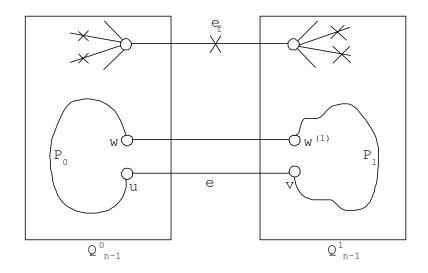


Figure 3.1: Illustration for of Theorem Case 1.1.

that  $e_f$  is a *j*th dimension edge. We decompose  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by dimension *j*. We then consider two further cases:

**1.1**  $e_f = (x, y)$  and e = (u, v) are in the same dimension. Thus, j = i and  $e_f \in D_i$ . (Fig. 3.1) In this case, e is an edge crossing  $Q_{n-1}^0$  and  $Q_{n-1}^1$ . Without loss of generality, assume that  $u \in V(Q_{n-1}^0)$  and  $v \in V(Q_{n-1}^1)$ . Since  $n \ge 5$ , u has a neighbor vertex  $w \in Q_{n-1}^0$ , by the definition of hypercube,  $w^{(1)}$  is a neighbor of v such that the edge  $(w, w^{(1)})$  is a healthy edge and  $(w, w^{(1)})$  is a crossing edge between  $Q_{n-1}^0$  and  $Q_{n-1}^1$ . By lemma 1, there exists a cycle  $C_0$  in  $Q_{n-1}^0 - F_L$  passing through (u, w) of every even length  $4 \le l(C_0) \le 2^{n-1}$  and a cycle  $C_1$  in  $Q_{n-1}^1 - F_R$  going through  $(v, w^{(1)})$  of every even length  $4 \le l(C_1) \le 2^{n-1}$ . We write  $C_0$  as  $\langle u, P_0, w, u \rangle$ , and  $C_1$  as  $\langle v, P_1, w^{(1)}, v \rangle$ . Thus,  $\langle u, P_0, w, w^{(1)}, v, u \rangle$  is a cycle of length 6 with  $l(P_0) = 3$ . For every even l,  $8 \le l \le 2^n$ , we may choose  $C_0$  and  $C_1$  such that  $l(C_0) = l(C_1) = \frac{l}{2}$ . Thus,  $\langle u, P_0, w, w^{(1)}, P_1, v, u \rangle$  can form a cycle of length l through e in  $Q_n - F$ , if we adjust the length of  $P_0$  and  $P_1$  properly.

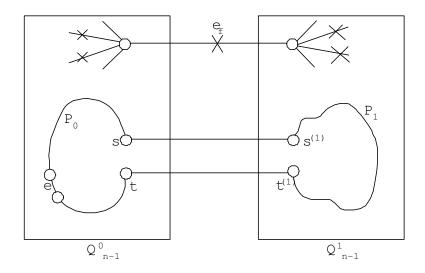


Figure 3.2: Illustration for of Theorem Case 1.2.

1.2  $e_f$  and e are in different dimension. Thus,  $j \neq i$  and  $e_f \notin D_i$ . (Fig 3.2) In this case, e is in  $Q_{n-1}^0$  or  $Q_{n-1}^1$ . Without loss of generality, we may assume that  $e \in E(Q_{n-1}^0)$ . By Lemma 1, there exists a cycle C in  $Q_{n-1}^0 - F_L$  going through the edge e of every even length l,  $6 \leq l \leq 2^{n-1}$ . Let  $C_0$  be a cycle of length  $2^{n-1} - 2$  or  $2^{n-1}$  passing through e in  $Q_{n-1}^0 - F_L$ . Since  $n \geq 5$ , there exists an edge (s,t) on  $C_0$  such that neither s nor t is adjacent to  $e_f$  and  $(s,t) \neq e$ . We write  $C_0$  as  $\langle s, P_0, t, s \rangle$ . By definition,  $(s^{(1)}, t^{(1)})$  is an edge in  $Q_{n-1}^1 - F_R$  through  $(s^{(1)}, t^{(1)})$  of every even length  $4 \leq l(C_1) \leq 2^{n-1}$ . We write  $C_1$  as  $\langle s^{(1)}, P_1, t^{(1)}, s^{(1)} \rangle$ . For every even  $l, 2^{n-1} + 2 \leq l \leq 2^n$ .  $\langle s, P_0, t, t^{(1)}, P_1, s^{(1)}, s \rangle$  is a cycle of length l going through e in  $Q_n - F$  if we adjust the length of  $P_0$  and  $P_1$  properly.

**Case 2**: There is exactly one vertex in  $Q_n$  incident with (n-2) faulty edges. Let x be the vertex having (n-2) faulty edges incident with it. Let  $f_1$  and  $f_2$  be two faulty edges incident with x, so  $f_1$  and  $f_2$  are in different dimensions j and k. By Lemma

9, decomposing  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by either *j*th or *k*th dimension, both  $F_L = F \cap E(Q_{n-1}^0)$  and  $F_R = F \cap E(Q_{n-1}^1)$  are sets of conditional faulty edges in  $Q_{n-1}^0$  and  $Q_{n-1}^1$  respectively. Between dimension *j* and *k*, we choose one to decompose  $Q_n$  into  $Q_{n-1}^0$  and  $Q_{n-1}^1$ , say dimension *j*, such that the required edge *e* is not a crossing edge. Therefore, there is a faulty edge crossing  $Q_{n-1}^0$  and  $Q_{n-1}^1$ , we denote this edge by  $e_f$ , and  $e_f \in F \cap D_j$  is incident with *x*. Without loss of generality, we may assume that  $x \in V(Q_{n-1}^0)$ .

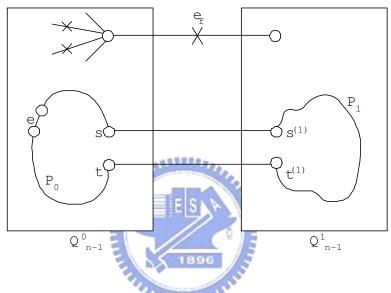
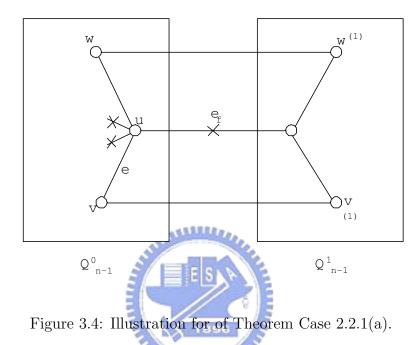


Figure 3.3: Illustration for of Theorem Case 2.1.

**2.1**: Suppose  $|F_L| \leq 2n - 7$  and  $|F_R| \leq 2n - 7$ . (Fig. 3.3) Without loss of generality, we further assume that  $e \in E(Q_{n-1}^0)$ . By induction hypothesis, there exists a cycle Cin  $Q_{n-1}^0 - F_L$  of every even length  $6 \leq l(C) \leq 2^{n-1}$  passing through e. Let  $C_0$  be a cycle of length  $2^{n-1} - 4 \leq l(C_0) \leq 2^{n-1}$  through e in  $Q_{n-1}^0 - F_L$ . Since  $|C_0 - e| \geq 2^{n-1} - 4 - 1 > 2(2n - 5) = 2|F \cap D_j|$ , for all  $n \geq 5$ . There exists an edge (s, t) on  $C_0$ such that (s, t) is not e, and both  $(s, s^{(1)})$  and  $(t, t^{(1)})$  are healthy edges. We write  $C_0$  as  $\langle s, P_0, t, s \rangle$ . By induction hypothesis, there exists a cycle  $C_1$  in  $Q_{n-1}^1 - F_R$  of every even length  $6 \leq l(C_1) \leq 2^{n-1}$  passing through  $(s^{(1)}, t^{(1)})$ . We write  $C_1$  as  $\langle s^{(1)}, P_1, t^{(1)}, s^{(1)} \rangle$ . For every even l,  $2^{n-1} + 2 \le l \le 2^n$ ,  $\langle s, P_0, t, t^{(1)}, P_1, s^{(1)}, s \rangle$  can form a cycle of length l going through e in  $Q_n - F$ , if we adjust the length of  $P_0$  and  $P_1$  properly.

**2.2**:  $|F_L| = 2n - 6$  or  $|F_R| = 2n - 6$ , say the former case. In this case,  $|F \cap D_j| = 1$ and  $|F \cap E(Q_{n-1}^1)| = |F_R| = 0$ .



**2.2.1**: *e* is in subcube  $Q_{n-1}^0$ . To find a cycle of length 6 passing through e = (u, v), we discuss the case that whether *e* is incident with *x* or not. If *e* is incident with *x*, without loss of generality, we assume that u = x. (Fig. 3.4) Thus,  $(v, v^{(1)})$  is a healthy edge. Since  $F_L$  is a set of conditional faulty edges in  $Q_{n-1}^0$ , vertex u = x has two healthy edges incident with it. Let *w* be a neighbor of *u* in  $Q_{n-1}^0$  such that (w, u) and  $(w, w^{(1)})$  are healthy edges and  $w \neq v$ . Thus,  $\langle u, v, v^{(1)}, u^{(1)}, w^{(1)}, w, u \rangle$  is a cycle of length 6 in  $Q_n - F$ . Otherwise, *e* is not incident with *x*, then  $(u, u^{(1)})$  and  $(v, v^{(1)})$  are healthy edges. (Fig. 3.5) By Lemma 1, there exists a cycle  $C_1 = \langle u^{(1)}, P_1, v^{(1)}, u^{(1)} \rangle$  of length four in  $Q_{n-1}^1$  through the edge  $(u^{(1)}, v^{(1)})$ . Thus,  $\langle u, u^{(1)}, P_1, v^{(1)}, v, u \rangle$  is a cycle of length 6 in  $Q_n - F$ , where  $l(P_1) = 3$ .

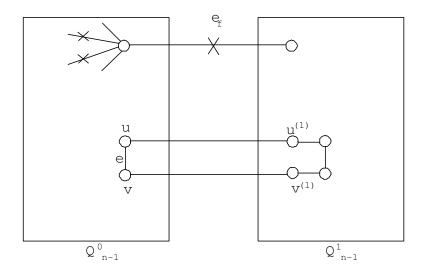


Figure 3.5: Illustration for of Theorem Case 2.2.1(b).

Let  $e_1$  be a faulty edge in  $Q_{n-1}^0$  that is not adjacent to  $e_f$ . Though  $e_1$  is a faulty edge, we treat it as a healthy edge temporarily, then the total number of faulty edge in  $Q_{n-1}^0$  is 2n-7. By induction hypothesis, there exists a cycle  $C_0$  of every length  $6 \leq l(C_0) \leq 2^{n-1}$ going through e in  $Q_{n-1}^0 - \{F_L - \{e_1\}\}$ . If  $C_0$  passes  $e_1$ , we choose  $e_1$ , or else, we choose any one edge on  $C_0$  which is not adjacent to  $e_f$ . Let the chosen edge be denoted by (s, t). We write cycle  $C_0$  as  $\langle s, P_0, t, s \rangle$ . Since  $|F \cap D_j| = 1$  and  $|F_R| = 0$ ,  $(s, s^{(1)})$ ,  $(t, t^{(1)})$  and  $(s^{(1)}, t^{(1)})$  are all healthy edges. Thus,  $\langle s, P_0, t, t^{(1)}, s^{(1)}, s \rangle$  is a cycle of length 8 in  $Q_n - F$ if  $l(P_0) = 6$ . Suppose that  $10 \leq l \leq 2^n$  and l is even. By Lemma 1, in  $Q_{n-1}^1$ , there exists a cycle  $C_3$  of length  $4 \leq l(C_3) \leq 2^{n-1}$  passing through  $(s^{(1)}, t^{(1)})$ . We write  $C_3$ as  $\langle s^{(1)}, P_3, t^{(1)}, s^{(1)} \rangle$ . Thus,  $\langle s, P_0, t, t^{(1)}, P_3, s^{(1)}, s, t \rangle$  is a cycle of length l through e in  $Q_n - F$ , if we adjust the length of  $P_0$  and  $P_3$  properly.

**2.2.2**: *e* is in subcube  $Q_{n-1}^1$ . (Fig. 3.6) By Lemma 1, there exists a cycle *C* of every even length  $4 \le l \le 2^{n-1}$  passing through *e* in  $Q_{n-1}^1$ . Suppose that  $2^{n-1} + 2 \le l \le 2^n$  and

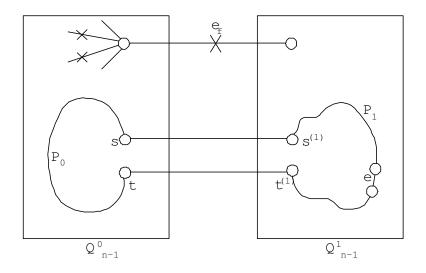


Figure 3.6: Illustration for of Theorem Case 2.2.2.

l is even. Since  $F_L$  is a set of conditional faulty edges, there are at most (n-3) faulty edges adjacent to  $e_f$  in  $Q_{n-1}^0$ . For  $n \ge 5$ ,  $n-3 \ge 2$ , we can choose a faulty edge  $e_2 = (s,t)$ in  $Q_{n-1}^0$  such that  $e_2$  is not adjacent to  $e_f$  and  $(s^{(1)}, t^{(1)})$  is not e. Treating the edge  $e_2$  as a healthy edge, by induction hypothesis, there exists a cycle  $C_0$  of length  $6 \le l(C_0) \le 2^{n-1}$ going through  $e_2$  in  $Q_{n-1}^0 - F_L$ . We write  $C_0$  as  $\langle s, P_0, t, s \rangle$ , and observe that  $(s, s^{(1)})$  and  $(t, t^{(1)})$  are healthy edges. By Lemma 4, there exists a cycle  $C_1$  of every length  $2^{n-1} - 4$ ,  $2^{n-1} - 2$ , or  $2^{n-1}$  through  $(s^{(1)}, t^{(1)})$  and e in  $Q_{n-1}^1$ . We write  $C_1$  as  $\langle s^{(1)}, P_1, t^{(1)}, s^{(1)} \rangle$ . Thus,  $\langle s, P_0, t, t^{(1)}, P_1, s^{(1)}, s \rangle$  is a cycle of even length l through e in  $Q_n - F$ , if we adjust the length of  $P_0$  and  $P_1$  properly.

**Case 3**: Every vertex in  $Q_n$  is incident with at most (n-3) faulty edges. In this case, suppose that e = (u, v) is in dimension *i*. By Lemma 10,  $Q_n$  can be decomposed into  $Q_{n-1}^0$  and  $Q_{n-1}^1$  by a dimension *j* different from *i* such that *e* is not a crossing edge and not all the faulty edges are in the same subcube. Then  $|F_L| \leq 2n - 6$  and  $|F_R| \leq 2n - 6$ . Next, we consider two further cases:

**3.1**: At least one faulty edge is a jth dimension edge. Thus,  $|F \cap D_j| \neq 0$ .

We then consider two cases: (a) $|F_L| \le 2n-7$  and  $|F_R| \le 2n-7$ , and (b) $|F_L| = 2n-6$ or  $|F_R| = 2n-6$ . The proof of this subcase is exactly the same as that of case 2.

**3.2**: None of the faulty edges is a *j*th dimension edge. Thus,  $|F \cap D_j| = 0$ .

**3.2.1**:  $|F_L| \leq 2n - 7$  and  $|F_R| \leq 2n - 7$ . Without loss of generality, we may assume that  $e \in E(Q_{n-1}^0)$ . By induction hypothesis, there exists a cycle C of even length  $6 \leq l(C) \leq 2^{n-1}$  in  $Q_{n-1}^0 - F_L$  passing through e. Let  $C_0$  be a cycle of even length  $2^{n-1} - 4 \leq l(C_0) \leq 2^{n-1}$  going through e in  $Q_{n-1}^0 - F_L$ . There exists an edge (s,t) other than e in  $C_0$ . Since  $|F \cap D_j| = 0$ ,  $(s, s^{(1)})$  and  $(t, t^{(1)})$  are healthy edges. We write  $C_0$  as  $\langle s, P_0, t, s \rangle$ . By induction hypothesis, there exists a cycle  $C_1$  of every even length  $6 \leq l(C_1) \leq 2^{n-1}$  in  $Q_{n-1}^1 - \{F_R - (s^{(1)}, t^{(1)})\}$  through  $(s^{(1)}, t^{(1)})$ . We write  $C_1$  as  $\langle s^{(1)}, P_1, t^{(1)}, s^{(1)} \rangle$ . For even  $l, 2^{n-1} + 2 \leq l \leq 2^n$ ,  $\langle s, P_0, t, t^{(1)}, P_1, s^{(1)}, s \rangle$  can form a cycle of length l going through e in  $Q_n - F$ , if we adjust the length of  $P_0$  and  $P_1$  properly.

**3.2.2**: Suppose  $|F_L| = 2n - 6$  or  $|F_R| = 2n - 6$ , say the former case. In this case,  $|F_R| = 1$ . We then consider two cases: (a) e is in subcube  $Q_{n-1}^0$ , and (b) e is in subcube  $Q_{n-1}^1$ .

(a) e = (u, v) is in subcube  $Q_{n-1}^0$ . Since  $|F \cap D_j| = 0$ , both  $(u, u^{(1)})$  and  $(v, v^{(1)})$ are healthy edges. Let l be an even number with  $6 \leq l \leq 2^{n-1}$ . By Lemma 1, there exists a cycle  $C_1$  of every even length from 4 to  $2^{n-1}$  passing through  $(u^{(1)}, v^{(1)})$  in  $Q_{n-1}^1 - \{F_R - (u^{(1)}, v^{(1)})\}$ . We write  $C_1$  as  $\langle u^{(1)}, P_1, v^{(1)}, u^{(1)} \rangle$ . No matter  $(u^{(1)}, v^{(1)})$  is healthy or not,  $\langle u, u^{(1)}, P_1, v^{(1)}, v, u \rangle$  forms a cycle of length l through e in  $Q_n - F$ . Suppose that  $2^{n-1} + 2 \leq l \leq 2^n$ . Let  $e_1$  be a faulty edge in  $Q_{n-1}^0$ . We may treat  $e_1$  as a healthy edges temporarily. By induction hypothesis, there exists a cycle  $C_0$  of length  $6 \leq l(C_0) \leq 2^{n-1}$ going through e in  $Q_{n-1}^0 - \{F_L - \{e_1\}\}$ . If  $C_0$  passes the edge  $e_1$ , we choose  $e_1$  to be deleted. Otherwise, we choose another edge other than e on cycle  $C_0$ . Let the chosen edge be denoted by (s, t). We write the cycle  $C_0$  as  $\langle s, P_0, t, s \rangle$ . Treating  $(s^{(1)}, t^{(1)})$  as a healthy edge, by Lemma 1, there exists a cycle  $C_3$  of every even length from 4 to  $2^{n-1}$ passing through  $(s^{(1)}, t^{(1)})$  in  $Q_{n-1}^1 - \{F_R - (s^{(1)}, t^{(1)})\}$ . We write  $C_3$  as  $\langle s^{(1)}, P_3, t^{(1)}, s^{(1)} \rangle$ , then  $\langle s, P_0, t, t^{(1)}, P_3, s^{(1)}, s \rangle$  is the cycle of length l through e in  $Q_n - F$ .

(b): e is in subcube  $Q_{n-1}^1$ . Let  $e_1$  be the only faulty edge in  $Q_{n-1}^1$ . By Lemma 1, there exists a cycle C of every even length from 6 to  $2^{n-1}$  through e in  $Q_{n-1}^1 - \{e_1\}$ . Suppose that  $2^{n-1} + 2 \leq l \leq 2^n$ , and l is even. Let  $e_0 = (s, t)$  be a faulty edge in  $Q_{n-1}^0$  such that  $(s^{(1)}, t^{(1)}) \neq e$ . By induction hypothesis, there exists a cycle  $C_0$  of length  $6 \leq l(C_0) \leq 2^{n-1}$  in  $Q_{n-1}^0 - \{F_L - \{e_0\}\}$  going through e. We write  $C_0$  as  $\langle s, P_0, t, s \rangle$ . If  $(s^{(1)}, t^{(1)}) = e_1$ , treat  $e_1$  as a healthy edge temporarily, by Lemma 4, there exists a cycle  $C_1$  of length  $2^{n-1} - 4$ ,  $2^{n-1} - 2$ , or  $2^{n-1}$  respectively going through both  $(s^{(1)}, t^{(1)})$  and e in  $Q_{n-1}^1$ . We write  $C_1$  as  $\langle s^{(1)}, P_1, t^{(1)}, s^{(1)} \rangle$ . Thus,  $\langle s, P_0, t, t^{(1)}, P_1, s^{(1)}, s \rangle$  can form a cycle of length l through e in  $Q_n - F$ , if we adjust the length of  $P_0$  and  $P_1$  properly. Otherwise, if  $(s^{(1)}, t^{(1)}) \neq e_1$ , by Lemma 6, there exists a cycle  $C_3$  of length  $2^{n-1} - 2$ , or  $2^{n-1} - 4$ , respectively, going through  $e^{1} - \{e_1\}$ . We write  $C_3$  as  $\langle s^{(1)}, P_3, t^{(1)}, s^{(1)} \rangle$ . Thus,  $\langle s, P_0, t, t^{(1)}, P_1, - \{e_1\}$ . We write  $C_3$  as  $\langle s^{(1)}, P_3, t^{(1)}, s^{(1)} \rangle$ . Thus,  $\langle s, P_0, t, t^{(1)}, P_3, s^{(1)}, s \rangle$  can form a cycle of length e in  $Q_n - F$ , if we adjust the length of  $P_0$  and  $P_1 - \{e_1\}$ . We write  $C_3$  as  $\langle s^{(1)}, P_3, t^{(1)}, s^{(1)} \rangle$ . Thus,  $\langle s, P_0, t, t^{(1)}, P_3, s^{(1)}, s \rangle$  can form a cycle of length l through e in  $Q_n - F$ , if we adjust the length  $Q_{n-1}^{1} - \{e_1\}$ . We write  $C_3$  as  $\langle s^{(1)}, P_3, t^{(1)}, s^{(1)} \rangle$ . Thus,  $\langle s, P_0, t, t^{(1)}, P_3, s^{(1)}, s \rangle$  can form a cycle of length l through e in  $Q_n - F$ , if we adjust the length of  $P_0$  and  $P_3$  properly.

This completes the proof.

## Chapter 4

## Conclusions

Since every component in the network may have different reliability, it is important to consider properties of a network with some conditional faults. We consider the *n*-dimensional hypercube with some faulty edges such that each vertex is incident to at least two nonfaulty edges. We use induction to prove that the *n*-dimensional hypercube  $Q_n$ ,  $n \ge 3$ , is (2n-5)-edge fault-tolerant conditional edge bipancyclic.

There exists an n-dimensional hypercube with (2n-4) edge faults, in which each vertex incident to at least two nonfaulty edges, such that for any pair of vertices does not exist path joining them. for example, let u = 00...0 and v = 100...1. One can consider the (2n-4) faulty edge in  $Q_n : e_i(u)$  and  $e_i(v)$  for all  $1 \le i \le n-2$  (see Fig. 1.1). obviously, vertices u and v each have exactly two nonfaulty edges incident to them. Hence the four edges  $e_0(u), e_n - 1(u), e_0(v)$ , and  $e_n - 1(v)$  from a 4-cycle by themselves. Therefore, in this  $Q_n, n \ge 3$ , it is impossible to make a faulty a hamiltonian cycle joining any pair of vertices. On the other hand, it is also impossible to make a hamiltonian cycle for  $n \ge 3$ . There, our result are optimal.In our future work, there is a direction to be studied:

If hypercube  $Q_n$  is faulty free, there exist n cycles with the same length being L for

each integer  $4 \leq L \leq 2^n$  . We will discuss that whether they are mutually independent or not.



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