

多媒體工程研究所

碩士論文

應用於 H. 264 可調式視訊串流之動態前向糾錯—失 真 最 佳 化 演 算 法 Dynamic FEC-Distortion Optimization for H.264 Scalable Video Streaming

研究生:溫維中

指導教授:蕭旭峰 教授

中華民國九十七年一月

在穩定性起伏大的通道中傳輸多媒體串流時,無論於應用層或是連接層,前向糾錯 編碼已經被視為一種適合用來提供服務質量(Quality of Service)的方案。在這篇論文中, 我們提出了兩種前向糾錯-失真最佳化的演算法,將頻寬適當地分配於視訊資料以及保 護資料,以求得到最佳的視訊品質。前向糾錯-失真最佳化演算法中的最佳化評量基準 建構於不平等錯誤保護,並且考慮了H.264/MPEG-4 AVC 可調式視訊編碼標準中所使用 的運動補償和跨層預測所帶來的誤差影響。同時,我們的演算法也適用於同視訊層的不 同幀會有不等視訊品質貢獻的情況。另外,我們也實做了輕量級的誤差隱蔽功能,用來 輔助演算法達到較佳的視訊品質。

111111

m

Abstract

Forward error correction codes have been shown to be a feasible solution either in application layer or in link layer to fulfill the need of Quality of Service for multimedia streaming over the fluctuant channels. In this paper, we propose the Dynamic FEC-distortion optimization algorithm to efficiently utilize the bandwidth for better video quality. The optimization criterions are based on the unequal error protection by taking account of the error drifting problems from both temporal motion compensation and inter-layer prediction of H.264/MPEG-4 AVC scalable video coding. Also, it can adapt to the content-dependent quality contribution of each video frame in a video layer. Lightweight error-concealment is also incorporated with the proposed algorithms for better H.264 SVC decoding performance.

Acknowledgement

能夠完成這篇論文,首先必須要感謝老師的指導。在無數次的挫折中,老師總能抽 絲剝繭地帶領我慢慢找出問題的癥結,點醒應該要注意的方向。老師也教導了我許多求 學上的正確態度,讓我獲益斐淺,最後終於完成這篇論文。同時也要感謝實驗室的同學 們,在我研究和生活遇到瓶頸的時候,總能適時地給予支持。最後,謹將這篇論文獻給

我最摯愛的父母以及家



摘要	I
ABSTRACT	II
ACKNOWLEDGEMENT	III
CONTENTS	IV
LIST OF TABLES	VI
LIST OF FIGURES	VII
I. INTRODUCTION	1 -
II. SCALABLE VIDEO CODING	3 -
2.1 H.264/MPEG-4 AVC Scalable Extension	3 -
2.2 MOTION PREDICTION	4 -
2.3 ERROR CONCEALMENT TOOL	5 -
III. FEC-DISTORTION OPTIMIZATION ALGORITHMS	- 8 -
3.1 FORWARD ERROR CORRECTION CODES	
3.1.1 Reed-Solomon code	9 -
3.1.2 How to generate a systematic Reed-Solomon erasure code	9 -
3.1.3 Encoding and decoding of systematic Reed-Solomon codes	10 -
3.2 Berkeley Algorithm	10 -
3.2.1 Component	10 -
3.2.2 Distortion	11 -
3 2 3 Algorithm	- 12 -
3.2.4 Problem when applying to H 264 SVC	13
3.2.4 From the approximation of 1.204 SV C	1J - 16
2.2.1.4L suid as	10 -
5.5.1 Algorithm	10 -
3.3.2 How to determine the a factor	18 -
3.3.3 Packetization	22 -
3.4 DYNAMIC FEC-DISTORTION OPTIMIZATION ALGORITHM	
3.4.1 Classifying Pictures to Clusters	25 -
3.4.2 Protecting Scheme	25 -
3.4.3 Algorithm	26 -
3.4.4 How to determine β factor	29 -
IV. SIMULATION RESULTS	32 -
4.1 SIMULATIONS FOR FLAT FEC-DISTORTION OPTIMIZATION	32 -

Contents

4.1.1 Environment	- 32 -
4.1.2 Simulation Results	33 -
4.2 SIMULATIONS FOR DYNAMIC FEC-DISTORTION OPTIMIZATION	36 -
4.2.1 Environment	36 -
4.2.2 Simulation Results	- 40 -
V. CONCLUSION	46 -
REFERENCE	47 -



List of Tables

Table 3-1	An example of picture sizes of a GOP	15
Table 3-2	Encoding parameters and PSNRs	17
Table 3-3	Experiment data	21
Table 3-4	Protection results	23
Table 3-5	Encoding parameters and PSNRs	30
Table 4-1	Encoding parameters and PSNRs	32
Table 4-2	Protection results	33
Table 4-3	Encoding parameters and PSNR	37
Table 4-4	Encoding parameters and PSNR	39
Table 4-5	Protection results	40



List of Figures

Fig 2-1	Basic coder structure for the scalable extension of H.264/MPEG-4 AVC
Fig 2-2	Hierarchical-B motion prediction
Fig 2-3a	Case 1: Error concealment when a non-base layer is lost
Fig 2-3b	Case 2: Error concealment when a base layer is lost
Fig 3-1	The Vandermonde matrix used in our FEC code
Fig 3-2	Component
Fig 3-3	Rate-Distortion curve
Fig 3-4	Packet lengths
Fig 3-5	Partially received layer may contribute PSNR
Fig 3-6	PacketLossRate-Alpha curve
Fig 3-7a	Alpha-PSNR curves of video "mobile"
Fig 3-7b	Alpha-PSNR curves of video "crew"
Fig 3-8	An example of picture classification
Fig 3-9	An example of DFDO protecting scheme
Fig 3-10	A binary tree for 3-layer video
Fig 3-11a	Beta-PSNR curve for video sequence "ice"
Fig 3-11b	Beta-PSNR curve for video sequence "city"
Fig 4-1	Given bandwidth
Fig 4-2	Simulation results different packet loss rates
Fig 4-3	Simulation results different packet loss rates
Fig 4-4a	Given bandwidth
Fig 4-4b	Given bandwidth
Fig 4-5a	Given bandwidth
Fig 4-5b	Given bandwidth
Fig 4-6a	Simulation result of FFDO
Fig 4-6b	Simulation result of DFDO
Fig 4-6c	The difference of simulation results
Fig 4-6d	The difference of simulation results
Fig 4-7a	Simulation result @ 15 fps
Fig 4-7b	Simulation result @ 15 fps
Fig 4-8	Simulation result with different packet loss rate
Fig 4-9a~f	Video sequences

I. Introduction

Personal, home, or handheld entertainment systems, such as DVB-H [1] and IPTV which is under construction to be a standard by ITU-T, have been an emerging research and industrial emphasis due to the great progress of the network communications and joint multimedia/channel coding technologies. It is rather challenging to fulfill the needs for Quality of Service and Quality of Experience requirements in the mobile environments of such entertainment systems that might suffer from dynamic channel fluctuation.

Besides Automatic Repeat reQuest (ARQ) which possibly suffers from the intolerable end-to-end packet delay and exacerbated jitter, forward error correction codes have been shown to be a feasible solution. In DVB-H, Multi-Protocol Encapsulated Forward Error Correction (MPE-FEC) is used by interleaving the information packet and the protection packets from Reed-Solomon code to deal with the burst error. The error protection strength in MPE-FEC is not really content-dependent. Besides Reed-Solomon code, rateless erasure codes (also known as fountain code [2]), such as raptor code [3], provide virtually infinite protection symbols and the modified version of such code has been recently adopted in 3GPP [4]. However, unlike Reed-Solomon error erasure code which shows maximum distance separable property, fountain codes generally have less coding efficiency.

In [5], Tan et al. proposed layered FEC for sub-band coded scalable video multicast using equation-based rate control while adaptive FEC is adopted to recover the lost packets so that the distortion function can be minimized with the optimized subscription of video and FEC layers, under an assumption that different frames in a video layer shall have the same distortion.

In [6], an adaptive FEC scheme as part of the reliable layered multimedia streaming over either unicast or multicast was proposed. The main objective of the FEC scheme is to maximize the streaming throughput while maintaining an upper bound of the error rate for each scalable video layer that FEC fails to decode. However, the upper bounds are pre-set without further explanation.

The impact of packet loss and FEC overhead on scalable bit-plane coded video in best-effort networks is analyzed in [7] and a similar optimization algorithm was proposed to allocate the bandwidth resource to FEC and video data.

In this paper, we propose FEC-Distortion optimization algorithms that take account of the error drifting problems from both temporal motion compensation and inter-layer prediction of H.264/MPEG-4 AVC scalable video coding, as well as the content-dependent visual quality contribution of each video frame in a video layer to achieve better quality of service with the same resource. In case of occasional packet error that is not recoverable by the FEC scheme, lightweight error-concealment is also incorporated with the proposed algorithms for better quality of reconstructed video.

The rest of this paper is organized as follows. In Chapter II we introduce the scalable video coding and the error concealment tool we added in the reference software. In chapter III we describe predecessor's work on the same topic as well as our two FEC optimization algorithms, Flat FEC-Distortion Optimization and Dynamic FEC-Distortion Optimization, to be used with H.264 scalable video coding, followed by Chapter IV, the simulation results and discussions. The concluding remarks are presented in Chapter V.

II. Scalable Video Coding

Traditional video coding techniques, such as waveform-based and content-dependent methods [14], aim to optimize the coding efficiency at a fix bit rate, and this presents a difficulty when multiple users try to access the same video through different communication links. [14] For example, a video encoded at 1.5 mbps can be downloaded and played at real time via a high-speed link, such as ADSL. But for users with only a modem connection at 56 kbps, it is not possible to receive sufficient bits for real-time play. [14] In such cases, the same video needs to be encoded at different bit rates to fulfill a range of users' abilities. Usually, a scalable video codec (SVC) provides scalability on SNR, spatial, temporal, or any combination of them, in a single encoded video stream. High-speed link users can choose to download the entire stream to have the best video experience, and for users with low-speed or low computation power environment, they can download partial of the same stream.

2.1 H.264/MPEG-4 AVC Scalable Extension

H.264/MPEG-4 AVC scalable extension is a standard developed jointly by ITU-T and ISO. These two groups created the Joint Video Team (JVT) to develop the H.264/MPEG-4 standard which has features such as hierarchical prediction structure, usage and extension of the network abstraction layer (NAL) unit concept of H.264/MPEG-4 AVC, and fine granular quality scalabilities, etc. [13] JVT also provides a reference software named JSVM, which not

only can encode/decode videos according to the H.264 SVC standard but also provides a toolset to down-sample/up-sample contents and calculate PSNR.

Figure 2-1 shows the basic structure of H.264 SVC, where the base layer is AVC compatible and the other layers provide scalabilities such as SNR, spatial, and temporal. In the decoder side, in order to decode a picture of higher layer, all the data of lower layers of the same picture are required. The more the video layers received, the better the video quality will



Fig 2-1: Basic coder structure for the scalable extension of H.264/MPEG-4 AVC [13]

2.2 Motion Prediction

In video coding, motion compensation is a technique used to improve coding efficiency. Since a video is composed by a series of continuous pictures and neighborhood pictures are usually very similar, we can use the reconstruction of the former picture as a prediction of later one. By encoding the difference between the prediction picture and the actual one, there usually will be a large range of 0s in the residual picture, and the number of bits can be further reduced by run-length coding and entropy coding. A picture is called I-picture or intra-coded if it needs no reference picture; a picture is called P-picture or inter-coded if it needs reference picture. B-picture is a special case of P-picture. A B-picture can use up to two reconstructed pictures as its prediction.

Currently, the prediction method supported by the reference software JSVM is only hierarchical-B. Figure 2-2 illustrates an example. The first picture of a video sequence is intra-coded as IDR (instantaneous decoder refresh) picture. A key picture and all the others pictures located between two key pictures (the IDR picture is also a key picture) form a group of pictures (GOP). The key pictures are either intra-coded or inter-coded by using previous key pictures as reference. In general, key pictures represent the minimal temporal resolution that can be decoded. [13]



A → B: A is reference by B

Fig 2-2: Hierarchical-B motion prediction

2.3 Error Concealment Tool

In multimedia communications, error resilient and error concealment are both tools used

to moderate the quality degradation when a video stream can't be fully received. Error resilient means this function is provided by encoder, such as including some extra headers in stream to help the detection of lost information, or simple duplication of encoded frames. In contrast, error concealment is implemented at the decoder side only.

When a scalable layer of a picture is lost, the decoder in the reference software JSVM may either crashes or skips the entire picture. However, even if the decoder doesn't crash, the decoded video still suffers from the error propagation and the visual quality dramatically decreases. In our implementation of error concealment, we fix this problem by frame duplication and interpolation. Thus, further reference to this lost frame can be found and the quality degradation caused by reference error is moderated. Figure 2-3a and 2-3b illustrate two examples of how the error concealment tool works when picture loss occurs. When a non-base layer is lost, decoder uses the reconstructed frame of its lower layer as its replacement. On the other hand, if the base layer of a picture is lost, decoder replaces the lost frame with the reconstructed base layer of nearest picture along the

time axis, and then resamples the replacement to proper resolution.



Fig 2-3a: Case 1: Error concealment when a non-base layer is lost



Fig 2-3b: Case 2: Error concealment when a base layer is lost



III. FEC-Distortion Optimization Algorithms

In this chapter, we will briefly introduce forward error correction code and then describe the algorithm proposed by Tan et al. in [5] as well as ours. When applied to a scalable video stream, these algorithms choose a subscription – which records how many video layers or clusters of video data to be protected and the corresponding protection strength so that the minimal distortion for each GOP can be approached. We also describe the modified version of [5] to apply Tan's algorithm to the scalable extension of H.264/AVC codec, which is denoted as Flat FEC-Distortion Optimization algorithm described in Section 3.3.

3.1 Forward Error Correction Codes

Forward error correction (FEC) code is a technique used to protect data through transmission. The advantage of using FEC code over Automatic Repeat reQuest (ARQ) on the multimedia streaming is that it can avoid the retransmission of lost data which will increase the total delay and is not feasible in real-time play situations such as live shows or video conferences.

Different types of FEC codes have different properties and error-correction abilities. An error detection/correction code recovers information by two steps. First, it detects whether if error occurs and the location of errors. And then it tries to correct the corruption. On the other hand, error erasure codes know the location of errors in advance, and it can recover the source

information as long as the number of received encoded information exceeds a predefined threshold. We use error erasure codes because in the case of transmitting video streams, receivers know which packets are lost. All the algorithms described in later sections use one of the most famous systematic erasure codes, Reed-Solomon code [10], to protect video streams. The term "systematic" means that the information data is part of the encoded data.

3.1.1 Reed-Solomon code

Reed-Solomon code was introduced by I.S. Reed and G. Solomonand. It is widely used in data storage, data transmission, mail encoding, and satellite transmission. The minimal data unit in Reed-Solomon codes is symbol, which may consist of one or more bits. Before encoding, information will be divided into some message blocks which consist of a predetermined number of symbols. In a Reed-Solomon erasure code with parameter pair (N, K), a message block consists of K source symbols and is protected by generated N-K protecting symbols. At receiver side, it can recover information as long as there are more than or equal to K received encoded symbols.

3.1.2 How to generate a systematic Reed-Solomon erasure code

We use the FEC code provided in [10] as our systematic Reed-Solomon erasure code. It uses Vandermonde matrix [16] to construct the generating matrix G, which is used to encode the source information as well as to decode the protected message blocks. Building generating matrices by Vandermonde matrix is faster than any other systematic MDS [12,15] erasure codes due to matrix inversion and matrix-vector multiplication [16,17].

The (i,j) element of generating matrix of Vandermonde in [10] can be formed as shown in figure 3-1, where the x_i 's are elements of $GF(p^r)$, p is a prime and r is the number of bits in a symbol. For more detail information about Reed-Solomon erasure code and Vandermonde matrix, please refer to [12,15,16,17].

$$g_{ij} = x_i^{j-1}$$

Fig 3-1: The Vandermonde matrix used in our FEC code 3.1.3 Encoding and decoding of systematic Reed-Solomon codes

Before encoding, source information needs to be divided into a number of message blocks, which consist of a predefined number of symbols that can be recognized by Reed-Solomon code. A message block is encoded by multiplying it with the generating matrix, and a protected message block can be decoded by multiplying it with the inverse of the generating matrix. 11111

3.2 Berkeley Algorithm

3.2.1 Component

The video codec used in [5] is 3D-wavelet sub-band coding, where each sub-band of a picture will be divided into N coefficient blocks of equal sizes and a component is formed by grouping coefficient blocks from different sub-bands. As illustrated in figure 3-1, there are seven sub-bands and each is divided into nine equal sizes coefficient blocks. Then, a component is formed by grouping a coefficient block of different spatial location from each sub-band.

A GOP (group of pictures) can be divided into M components, and each component consists of many video layers, where each video layer has different playing rate. Thus, spatial and temporal scalabilities are constituted.

Components are independently decodable because they are independently variable length coded. Besides, since a component is formed by different spatial locations of coefficient blocks, the distortion of losing any component can be assumed to be equal.



Fig 3-2 [8]: A component is formed by grouping a coefficient block of different spatial location from each sub-band.

3.2.2 Distortion

The estimation of distortion comes from a pre-measured Rate-Distortion curve, as illustrated in figure 3-3.



Fig 3-3 [8]: Rate-Distortion curve.

3.2.3 Algorithm

By giving available bandwidth and packet lost rate, the proposed algorithm of Tan et al. chooses a subscription that has the minimal distortion to determine the number of video layers to be protected and the protection strength for each layer. Equation (1) describes the main idea. In order to find the best subscription S^* , the function D(s,p) is used to estimate distortion of current GOP when applying subscription s and packet lost rate p. R(s) is the bandwidth used by applying subscription s and must be less than or equal to the available bandwidth B. M is the set of all possible subscriptions under B. The final version of distortion function used in (1) is (4), which is derived from (2) and (3).

Since components of the same layer of a GOP are independently decodable and have similar importance, the distortion of a GOP is about to be the sum of distortions of all components of all subscribed layers. Thus, equation (2) is formed. In (2), M is the number of components, L is the number of subscribed layers, $p_i^{(c)}$ means the decodable probability of the only first i layers of component c, and $D_i^{(c)}$ is the corresponding distortion which is measured by rate-distortion curve mentioned above.

Again, since all components are similarly important, Tan et al. assumes that distortions of all components are the same, thus equation (3) can be derived from (2). And by defining $D_i = MD_i^{(1)}$, equation (4) can be derived from (3), and is the final version of distortion function used in (1).

In equation (6), $q_i^{(c)}$ is the probability that a lost packet of component c of layer i is unrecoverable. The decodable probability of the only first i layers, $p_i^{(c)}$, is illustrated as (5). By the definition of Reed-Solomon code, when using parameter pair (N, K), there must be more than or equal to K out of N symbols received to recover all source symbols.

$$S^* = \underset{s \in M, R(s) \leq B}{\operatorname{arg min}} D(s, p) \tag{1}$$

)

$$D(s,p) \approx \sum_{c=1}^{M} D^{(c)} = \sum_{c=1}^{M} \sum_{i=0}^{L} p_i^{(c)} \cdot D_i^{(c)}$$
(2)

$$D(s,p) \approx M \sum_{i=0}^{L} p_i^{(1)} \cdot D_i^{(1)}$$
(3)

$$D(s,p) \approx \sum_{i=0}^{L} p_i \cdot D_i$$
(4)

$$p_{i}^{(c)} = \begin{cases} q_{i}^{(c)} \prod_{k=0}^{i-1} \left(1 - q_{k}^{(c)} \right) & \text{if } 0 \le i < L \\ \prod_{k=0}^{L-1} \left(1 - q_{k}^{(c)} \right) & \text{if } i = L \end{cases}$$
(5)

$$q_{i}^{(c)} = p \cdot \left[\sum_{w=0}^{M-1} \binom{M + k_{i} - 1}{w} (1 - p)^{w} p^{M + k_{i} - 1 - w} \right]$$
(6)

3.2.4 Problem when applying to H.264 SVC

When applying the proposed algorithm in [5] to the scalable extension of H.264/AVC, it

faces a problem of packet length. In [5], each packet contains one component and the components of the same layer are supposed to be of similar sizes. But in the case of the scalable extension of H.264/AVC, the NAL sizes may vary in a large range, and this causes problem.

As shown in figure 3-4, where each row is a packet, since NALs may have different sizes, the packet lengths also differ. The problem occurs when the largest and smallest NAL sizes dramatically differ. Because the length of protecting packets must be the same as the size of the largest NAL, the utilization of bandwidth is very inefficient if the packet lengths dramatically vary.

Fig 3-4: The length of protecting packets must be the same as the largest data packet Table 3-1 is an example of NAL sizes which are extracted from an actually encoded video stream and the GOP size is 16. The red and bold numbers are the largest NAL sizes in corresponding video layers, and the blue and bold are smallest ones. We can observe that the largest NAL size is about 26 ~ 30 times larger than the smallest one.

Picture ID	Size (Byte) in 0 th layer	Size (Byte) in 1 st layer	Size (Byte) in 2 nd layer
1	92	150	125
2	186	255	260
3	100	165	112
4	602	618	411
5	98	194	172
6	222	366	291
7	107 E S	199	151
8	896	819	570
9	103	176	142
10	250	364	245
11	129	200	166
12	532	608	427
13	112	177	171
14	226	312	245
15	101	151	152
16	3034	4582	2962

Table 3-1: An example of picture sizes of a GOP

3.3 Flat FEC-Distortion Optimization Algorithm

In [5], Tan et al. proposed a layered FEC algorithm for sub-band coded scalable video multicast using equation-based rate control such that packet loss is one of the parameters to regulate the sending rate while adaptive FEC is adopted to recover the lost packets so that the distortion can be minimized with optimized subscription S* as described in (1) and (4), under an assumption that different frames in a video layer shall have the same distortion measure.

Instead of the sub-band scalable video coding with layered structure on both video and FEC data in [5], our proposed FEC optimization algorithms are based on the H.264/MPEG-4 AVC scalable extension, which is an amendment to the H.264/MPEG-4 AVC standard and it is scheduled to be finalized in 2007. The base layer of a Scalable Video Coding (SVC) bit-stream is usually coded in compliance with H.264 while new scalable tools are added for supporting spatial, SNR, and temporal scalability [9].

When applying to the scalable extension of H.264/AVC, the algorithm in [5] does not take the effect of reference pictures into account. That means, when a layer of a picture is lost and can't be recovered by FEC, the distortion of other pictures which reference the lost one would increase due to reference error, but in [5] this is not represented. In addition, [5] has the problem of inefficient bandwidth using as described in Section 3.2.4.

3.3.1 Algorithm

Our flat FEC-distortion optimization algorithm (FFDO) inherits the idea of equation (1)

and (2) from [5]: selecting a subscription that not only fits the available bandwidth but also has the maximal PSNR to determine the number of video layers to be protected and the protecting strength for each video layer. In contrast with [5], FFDO estimates the PSNR of a GOP of layers instead of PSNR of individual pictures. A factor α is added to the distortion measuring function to reflect the quality degradation caused by partially received video layer.

The PSNR of a GOP is the summation of PSNR of all subscribed layers. Equation (8) illustrates the contribution measuring function, where p_i is the FEC decodable probability of only the first i layers, D_i is the corresponding PSNR and L is the number of subscribed layers. The leaky factor α in (9) is a real value between 0 and 1, and is used to represent the PSNR contributed by the partially received higher layers. $S^* = \underset{s \in M, R(s) \leq B}{\operatorname{max}} \operatorname{PSNR}(s, p)$ (7) $PSNR(s, p) = \sum_{i=0}^{L} p_i \cdot D_i$ (8) $D_i = PSNR_{i-1} + \alpha \times (PSNR_i - PSNR_{i-1}), \quad 0 \leq \alpha < 1$ (9)

Use fig 3-5 to explain the leaky factor α , where the arrows indicate the reference relationships. In our implementation of error concealment tool, when a layer of picture is unavailable, decoder will take the reconstructed picture of lower layer to substitute the lost one and the further higher layers. In fig 3-5, since the lost NAL belongs to a picture which will not be referenced by any other pictures, the PSNR contribution of this video layer should be multiplied by a leaky factor α to reflect the quality degradation caused by the lost picture.



Fig 3-5: Partially received layer may contribute PSNR

Though the definitions of p_i in FFDO and [5] are the same, the definitions of q_i in both

algorithms are different. In FFDO, q_i is the FEC undecodable probability of layer i, as illustrated in (11), where the protecting parameter (N, K) of layer i is (M + k_i , M).

$$p_{i} = \begin{cases} q_{i} \prod_{k=0}^{i-1} (1 - q_{k}) & \text{if } 0 \leq i < L \\ \prod_{k=0}^{L-1} (1 - q_{k}) & \text{if } i = L \end{cases}$$
(10)
$$q_{i} = \sum_{w=0}^{M-1} \binom{M+k_{i}}{w} (1 - p)^{w} p^{M+k_{i}-w}$$
(11)

Sequence Parameter Set Network Abstraction Layer (NAL) units and Picture Parameter Set NAL units [11] have essential header information in order to decode the video properly and they are assigned strongest error correction code (n=256).

3.3.2 How to determine the α factor

100

The α factor in (9) embraces the video quality degradation caused by both partially received layer data and reference errors, and is too complex and time-consuming to estimate the accurate value of α from all picture layer lost cases in a GOP. Hence, we select α by experiments. We do experiments on two testing sequences, mobile and crew, with same environment parameters to see the variation tendency of best α when the packet lost rate varies from 0.12 to 0.48. The encoding parameters and the PSNR are listed in table 3-2, where the PSNRs are calculated by the source video and the up-sampled reconstructed layers. For each given packet loss rate, α is set from 0 to 0.9 to generate the protected video streams. Each protected stream passes 100 lossy channels which have the same specified packet loss rate but different random seeds to generate lost pattern. The experiments data are listed in table 3-3 and the

figure 3-6 shows the best	α	for both	videos with	different p	acker	loss rate.
---------------------------	---	----------	-------------	-------------	-------	------------

Table 3-2: Encoding parameters and PSINKs						
Layer ID	Resolution	Bit-rate (kbit/sec)	PSNR ^{mobile}	PSNR ^{crew}		
0	QCIF	200	21.5101	29.5959		
1	QCIF	400	21.9109	30.3443		
2	CIF	600	28.0416	31.5937		
3	CIF	800	29.6475	32.1536		
4	4CIF	1000	30.7404	33.4996		

Table 3-2: Encoding parameters and PSNRs

From figure 3-6 we can observe that the α factor becomes smaller as the packet loss rate increases. This is reasonable because the partially received layers lost more picture data as the increasing of packet loss rate. From figure 3-7, we observe that there is a large range of α with which the associated PSNRs don't differ much, for two videos with any packet loss rate. So we choose the α to be the median α values shown in figure 3-6 to be used in later simulations.



Fig 3-7a: Alpha-PSNR curves of video "mobile" under different packet loss rate



Fig 3-7b: Alpha-PSNR curves of video "crew" under different packet loss rate

Table 3-3: Experiment data to find the best α factor, where red are the largest PSNRs, green are the smallest ones, and PSNR(i, i-1) is the PSNR difference between protected layers

	α	# protected layers	PSNR(i, i-1) ^{mobile}	PSNR(i, i-1) ^{crew}	PSNR ^{mobile}	PSNR ^{crew}
	0	3~4	1.6059	0.5599	27.1221	31.3382
2	0.1	3~4	1.6059	0.5599	27.1701	31.3964
§ 0.1	0.2	3~4	1.6059	0.5599	27.1701	31.3964
Rate	0.3	3~4	1.6059	0.5599	27.1690	31.4436
[SSC	0.4	3~4	1.6059	0.5599	27.1655	31.4099
it Lo	0.5	3~4	1.6059	0.5599	27.1268	31.4101
acke	0.6	3~4	1.6059	0.5599	27.1268	31.3894
P.	0.7	3~4	1.6059	0.5599	27.0701	31.3808
	0.8	3~4	1.6059	0.5599	26.7671	31.3799
	0.9	3~4	1.6059	0.5599	26.0669	31.3280

	α	# protected layers	PSNR(i, i-1) ^{mobile}	PSNR(i, i-1) ^{crew}	PSNR ^{mobile}	PSNR ^{crew}
	0	2~3	6.1307	1.2494	23.7793	30.7573
4	0.1	2~3	6.1307	1.2494	23.7793	30.7577
0.2	0.2	2~3	6.1307	1.2494	23.7793	30.7802
Rate	0.3	2~3	6.1307	1.2494	23.7793	30.7802
l ssc	0.4	2~3	6.1307	1.2494	23.7100	30.7748
st Lo	0.5	2~3	6.1307	1.2494	23.7100	30.7736
acke	0.6	2~3	6.1307	1.2494	22.9704	30.7736
P	0.7	2~3	6.1307	1.2494	22.9704	30.6994
	0.8	2~3	6.1307	1.2494	22.8145	30.6041
	0.9	2~3	6.1307	1.2494	22.6387	30.3538
	α	# protected layers	PSNR(i, i-1) ^{mobile}	PSNR(i, i-1) ^{crew}	PSNR ^{mobile}	PSNR ^{crew}
	0	2~3	6.1307	1.2494	23.1585	30.2710
36	0.1	2~3	6.1307	1.2494	23.1615	30.3509
e 0.3	0.2	2~3	6.1307	1.2494	23.1615	30.3453
Rate	0.3	2~3	6.1307	1.2494	23.1330	30.3445
SSC	0.4	2~3	6.1307	1.2494	23.1265	30.3479
st Lo	0.5	2~3	6.1307	1.2494	23.1376	30.3315
acke	0.6	2~3	6.1307	1.2494	23.1139	30.3272
Р	0.7	2~3	6.1307	1.2494	23.0831	30.3096
	0.8	2~3	6.1307	1.2494	23.0235	30.3023
	0.9	2~3	6.1307	1.2494	22.8426	30.2908
	α	# protected layers	PSNR(i, i-1) ^{mobile}	PSNR(i, i-1) ^{crew}	PSNR ^{mobile}	PSNR ^{crew}
	0	1~2	0.4008	0.7484	21.7022	29.8771
48	0.1	1~2	0.4008	0.7484	21.6898	29.8818
e 0.∠	0.2	1~2	0.4008	0.7484	21.6542	29.8577
Rate	0.3	1~2	0.4008	0.7484	21.6217	29.8512
oss	0.4	1~2	0.4008	0.7484	21.6417	29.8297
et L	0.5	1~2	0.4008	0.7484	21.5526	29.7990
acki	0.6	1~2	0.4008	0.7484	21.5098	29.7108
Ą	0.7	1~2	0.4008	0.7484	21.4703	29.4410
	0.8	1~2	0.4008	0.7484	21.4504	29.4487
	0.9	1~2	0.4008	0.7484	21.3699	29.3134

3.3.3 Packetization

As mentioned in Section 3.3, when applying [5] to the extension of H.264/AVC, the bandwidth cannot be utilized efficiently due to the dramatically differing packet sizes. In our

design of FFDO, since we estimate distortion of a GOP from the point of view of layers instead of individual pictures, we have no need to enforce a packet to contain exactly one picture layer data. In our implementation, every packet has the same length and it is possible to have many NAL units in one packet or an NAL can be split into more than one packet. The packet length of a FEC coding session which protects a layer of a GOP is the ceiling of video data size/K, where K is from the protecting parameter (N, K) and is also the number of video data packets.

Table 3-4 lists an example of protections determined by [5] and FFDO under the same bandwidth and the same video bitstream. From the table and Fig. 3.4, it is easy to find out that FFDO uses bandwidth much more efficiently than the proposed FEC selection algorithm in [5], i.e., more bandwidth resource can be allocated to protection packets.

Table 3-4: Protection results determined by [5] and FFDO under same bandwidth ("0" denotes only the video data is sent and "x" denotes neither video nor protection data are sent.)

COP	Layer:	FFDO		Algor	ithm in [5]
(Rondwdith)	video size	Pkt Length	# of Protecting	Pkt Length	# of Protecting
(Dalluwului)	(Byte)	(Byte)	Packets	(Byte)	Packets
0	0: 12548	793	17	7948	4
(75526)	1: 13419	847	14	5081	0
(75550)	2: 12570	794	14	8636	0
1 (49508)	0: 12635	797	18	7932	1
	1: 13599	858	10	5108	0
	2: 12650	Х	Х	8511	0
2	0: 12026	759	17	7889	4
ک (7/7/۹)	1: 13288	839	15	5194	0
(74740)	2: 12449	787	14	8508	0

3 (48698)	0: 12613	796	18	8182	1
	1: 13042	824	10	5268	0
	2: 12419	Х	Х	8691	0
4 (79258)	0: 12806	808	17	8164	8
	1: 14154	893	14	Х	Х
	2: 13306	840	14	X	X

3.4 Dynamic FEC-Distortion Optimization Algorithm

FFDO is based on the assumption that different frames in the same video layer exhibit constant distortion. However, this is usually not the case for the real H.264 SVC videos. The distortion (or PSNR) depends on the content of each video frame as well as the quantization parameter used in each block. Due to the error propagation effect resulting from not only the prediction coding across the video layers but also the temporal motion compensation coding in each individual video layer, the distortion caused by different frame of a video layer can also vary. As a result, the global optimal bit allocation of H.264 SVC and FEC shall be calculated over all the possible bit allocation and packet loss combinations.

The main objective of our Dynamic FEC-Distortion Optimization (DFDO) algorithm is to increase the decodable probability of important pictures when a layer of a GOP cannot be completely received. After FFDO determines the best parameter configuration that indicates the number of video layers to be protected and the protecting strengths, DFDO is applied to each subscribed layer to reallocate the protecting packets. DFDO first classifies pictures into a number of clusters which have different importance, and then reallocates the protecting packets to clusters according to their importance in order to minimize the distortion of whole layer (or maximize the PSNR of whole layer).

3.4.1 Classifying Pictures to Clusters

Figure 3-8 is an example of picture classification, where pictures are classified by their temporal level. As mentioned in Section 2.2, when hierarchical B prediction is applied, each video layer can be decomposed to a number of temporal levels and every picture is predicted by the reconstructed pictures of other two pictures of lower temporal levels. Therefore we can say that the pictures belong to lower temporal levels are more important than those of higher



In the implementation phase, a cluster may have unused spaces after all pictures of cluster are filled inside. When this happens, we append the pictures of the next cluster to the current one until there are no free spaces remained Pictures Clusters



3.4.2 Protecting Scheme

Figure 3-9 is an example of DFDO protecting scheme, blue parts are the clusters of video data and green parts are the FEC data used to protect source clusters enclosed within brackets. K_i is the number of packets in cluster i, and M_i denotes the number of protection packets in

cluster $(0,1,\ldots,i)$. The summation of M_i must be the same as the number of packets allocated

to this layer by Flat FEC-Distortion Optimization algorithm.



Fig 3-9: An example of DFDO protecting scheme

3.4.3 Algorithm

The main idea of Dynamic FEC-Distortion Optimization algorithm is to find the best allocating patterns of protecting symbols such that the distortion of the whole video layer is minimized (or the PSNR of the whole video layer is maximized), and can be illustrated as equation (12). An allocating pattern is a vector where each element determines the number of protecting symbols for corresponding cluster. M^{*} is the set of all possible patterns. PSNR^{*}(s,p) is the corresponding PSNR when applying allocating pattern s to a video layer with average packet loss rate p. The PSNR of whole video layer can be estimated as the summation of PSNR contributed by each cluster and is illustrated as equation (13), where p^*_i means the decodable probability of the only first i clusters and D^*_i is the corresponding PSNR contribution.

$$S'' = \underset{s \in M^*}{\operatorname{arg\,max}} \operatorname{PSNR}^*(s, p) \tag{12}$$

$$PSNR^{*}(s,p) = \sum_{i=0}^{C} p_{i}^{*} \cdot D_{i}^{*}$$
(13)

Figure 3-10 is a binary tree for 3-layer video. Each path which starts from root and ends at leaf forms a decoding path. A Y-edge denotes that the corresponding source cluster is decodable, and an N-edge means that the corresponding source cluster is undecodable. For example, Path 0Y1N2Y denotes the event that cluster 0 is decodable followed by cluster 1, which is undecodable, and further followed by decodable cluster 2. Since cluster (0,1,...,i) protects source clusters 0 to i at the same time, all source clusters can be decoded at the end of the path 0Y1N2Y.

The decodable probability of the only first i clusters, p_{i}^{*} , can be derived by such binary trees. In figure 3-10, p_{3}^{*} is the sum of the probability of decoding paths 0Y1Y2Y, 0Y1N2Y, 0N1Y2Y, and 0N1N2Y. And for p_{2}^{*} , its value is the sum of the probability of decoding paths 0Y1Y2N and 0N1Y2N. And so on.



Fig 3-10: A binary tree for 3-layer video

We summarize p_{1i}^* in equation forms, as illustrated in (14) to (17). Equation (15) corresponds to the decodable probability of paths which start from ID-Y and there are pvK packets received from source clusters 0 to ID-1. Equation (16) corresponds to the undecodable probability of paths which start from ID-N and there are pvK packets received from source clusters 0 to ID-1. Equation (17) calculates the minimal number of packets required to decode source clusters 0 to ID. C is the number of source clusters in layer, p is the average lost rate, K_i is the number of source symbols in cluster i and M_i is the number of protecting symbols of protecting cluster (0,1,...,i).

$$p_i^* = Y(0, i, 0) + N(0, i, 0) \tag{14}$$

$$Y(ID,t, pvK) = \begin{cases} 0, if \ ID \ge t \\ \sum_{i=0}^{K_{w}} \sum_{j=K(ID)-pvK-i}^{M_{w}} p^{K_{w}+M_{w}-i-j}(1-p)^{i+j} \times \\ \sum_{i=0}^{K_{w}} \sum_{j=K(ID)-pvK-i}^{M_{w}} \left[Y(ID+1,t,K(ID)) + \\ N(ID+1,t,K(ID)) \right] , otherwise \end{cases}$$
(15)

$$N(ID,t,pvK) = \begin{cases} 0, if \ ID = t-1 \\ \sum_{i=0}^{K_{iD}-1-pvK \ K(ID)-1-pvK-i} C_{i}^{K_{iD}} C_{j}^{M_{iD}} p^{K_{iD}+M_{iD}-i-j} (1-p)^{i+j} \times \\ \sum_{i=0}^{K_{iD}-1-pvK \ K(ID)-1-pvK-i} \left[Y(ID+1,t,pvK+i) + \\ N(ID+1,t,pvK+i) \right] , otherwise$$
(16)

$$K(ID) = \sum_{i=0}^{D} K_i$$
(17)

The distortion measuring function D_{i}^{*} , which is corresponded to p_{i}^{*} , can be estimated as equation (18), where PSNR_j is the summation of PSNR contributed by pictures in cluster j. The β is a leaky factor to reflect the quality degradation caused by reference to an unavailable and error concealed picture. The exponential form of β is due to the hierarchical B prediction.

$$D_{i}^{*} = \sum_{j=0}^{i-1} PSNR_{j} + \sum_{j=i}^{C-1} PSNR_{j} \times \beta^{j-i+1}, \quad 0 \le \beta < 1$$
(18)

3.4.4 How to determine β factor

The value of β is chosen experimentally. We protect the video sequences "ice" and "city" with α 0.15 and different β_s and select the best one which shows the highest PSNR. Each protected stream passes 100 lossy channels which have the same packet loss rate 0.285714 but different random seeds to generate lost pattern. The encoding parameters and PSNR for each video layer are listed in table 3-5. Figures 3-11a and 3-11b show the simulation results of sequence "ice" and "city", respectively. From figures 3-11a and 3-11b, we can observe that DFDO has stable and best performance when β is between 0.6 to 0.95. Thus, we choose 0.75 in further simulations.

Layer ID	Resolution	Bit-rate (kbit/sec)	PSNR ^{ice} (dB)	PSNR ^{city} (dB)
0	QCIF	200	30.1402	25.0721
1	QCIF	400	30.5005	25.2716
2	CIF	600	32.2488	26.7457
3	CIF	800	32.6895	27.2123
4	4CIF	1000	36.3041	30.9114

Table 3-5: Encoding parameters and PSNRs



Fig. 3-11a: Beta-PSNR curve for video sequence "ice"



IV. Simulation Results

4.1 Simulations for Flat FEC-Distortion Optimization

4.1.1 Environment

We prepare two 300-frame video sequences, mobile and crew, which are encoded with the same parameters listed in table 4-1. For each video sequence, we simulate with four packet loss rates which are 0.12, 0.24, 0.36 and 0.48. The packet lost patterns are generated with independent and identical distribution. The bandwidth is given as shown in figure 4-1, where for odd-ID GOPs we give more bandwidth and for even ones we give less. The α factor used in FFDO is 0.15, as described in Section 3.3.2. Every simulation value is averaged from 100 times of simulations with different loss patterns. The protection result for both videos with four packet loss rates are listed in table 4-2, and the graphs of simulation results are illustrated in figure 4-2 and 4-3.

Layer ID	Resolution	Bit-rate (kbit/sec)	PSNR ^{mobile}	PSNR ^{crew}
0	QCIF	200	21.5101	29.5959
1	QCIF	400	21.9109	30.3443
2	CIF	600	28.0416	31.5937
3	CIF	800	29.6475	32.1536
4	4CIF	1000	30.7404	33.4996

Table 4-1: Encoding parameters and PSNRs



Fig 4-1: Given bandwidth

4.1.2 Simulation Results

Table 4-2: Protection results	The number is the amount of	protecting symbols)
10010 . 20110000000000000000000000000000		

		mobile					crew			
		Packet loss rate			Packet loss rate					
GOP	Layer	0.12	0.24	0.36	0.48	0.12	0.24	0.36	0.48	
	0	8	16	18	28	6	14	23	29	
	1 //	7	15	14	17	4	10	17	19	
1	2	7	14	13	- 0	4	10	17	10	
	3	6	Х	Х	Х	3	5	Х	Х	
	4	Х	Х	Х	Х	2	Х	Х	Х	
	0	5	.17	21	46	8	13	24	38	
2	1	4	11	7	X	5	6	14	0	
	2	3	X	X	X	5	0	Х	Х	
	0	10	16	18	28	11	21	23	30	
3	1	7	15	14	-18	9	18	19	17	
5	2	7	15	14	0	8	17	15	11	
	3	5	X	X	X	7	Х	Х	Х	
	0	5	17	21	45	7	18	21	47	
4	1	4	11	7	Х	4	13	10	Х	
	2	3	Х	Х	Х	3	Х	Х	Х	
	0	8	16	18	29	10	19	21	32	
5	1	7	15	14	17	8	16	16	21	
5	2	7	14	13	0	7	16	14	0	
	3	6	Х	Х	Х	6	Х	Х	Х	
	0	6	16	22	46	7	17	20	45	
6	1	3	11	6	Х	3	12	9	Х	
	2	3	Х	Х	Х	2	Х	Х	Х	

	0	8	16	18	28	9	16	20	35
7	1	7	15	14	17	6	13	12	24
7	2	7	14	13	0	6	12	9	Х
	3	6	Х	Х	Х	4	Х	Х	Х
	0	5	16	21	46	6	18	21	46
8	1	4	11	6	Х	3	13	10	Х
	2	3	Х	Х	Х	2	Х	Х	Х
	0	9	17	18	28	9	18	20	38
0	1	7	15	14	17	6	14	13	28
9	2	7	14	14	0	6	12	11	Х
	3	6	Х	Х	Х	5	Х	Х	Х
	0	5	16	21	44	14	15	21	38
10	1	3	10	5	X	12	11	4	Х
	2	3	X	Х	X	Х	Х	Х	Х
	0	8	16	17	28	8	16	20	34
11	1	7	14	14	15	6	13	12	22
11	2	7	14	13	0	6	12	9	Х
	3	6	Х	X	X	-5	X	Х	Х
	0	5	16	21	41	7	14	23	38
12	1	4	11	6	Х	3	9	0	Х
	2	3	X	Х	X	1	X	Х	Х
	0	8	16	17	28	8	16	22	31
13	1	7	14	13	14	6	12	14	18
15	2	7	13	13	0	6	11	0	Х
	3	6	Х	Х	Х	4	Х	Х	Х
	0	5	16	20	41	13	15	23	38
14	1	3	10	6	Х	11	9	0	Х
	2	3	Х	Х	Х	Х	Х	Х	Х
	0	8	16	17	28	8	16	20	32
1.5	1	7	14	13	15	6	13	12	20
13	2	7	13	13	0	6	12	8	X
	3	6	X	X	X	5	X	X	Х
	0	5	16	22	41	7	16	20	41
16	1	4	11	4	X	3	10	6	Х
	2	3	Х	Х	Х	1	Х	Х	Х



Fig. 4-2: Simulation results of sequence mobile with different packet loss rates



Fig 4-3: Simulation results of sequence crew with different packet loss rates

4.2 Simulations for Dynamic FEC-Distortion Optimization

4.2.1 Environment

We first encode two sequences, "football" and "soccer", both have frame rate 30 fps, to compare the protection performance between FFDO and DFDO. The encoding parameters are listed in table 4-3. The packet loss rate is 0.285714; the α factor used in FFDO is 0.15, and the β factor used in DFDO is 0.75. The bandwidth distributions for two videos are given in figure 4-4.





Fig 4-4b: Given bandwidth for soccer @ 30 fps

Table 4-3: Encoding parameters and PSNR for sequences football and soccer at 30fps

Layer ID	Resolution	Bit-rate (kbit/sec)	PSNR ^{football} (dB)	PSNR ^{soccer} (dB)
0	QCIF	200	27.5387	28.3235
1	QCIF	400	28.9634	28.9213
2	CIF	600	31.0809	30.0142
3	CIF	800	32.1219	30.4611
4	4CIF	1000	32.3232	32.1135

Figures 4-6a and 4-6b are the simulation results of football with FFDO and DFDO, respectively. For the convenience of comparison, we also present the graph of the PSNR difference between DFDO and FFDO, as shown in 4-6c, where the positive area indicates that DFDO outperforms FFDO by 0.4 dB in average. In fig 4-6d, we compare both algorithms with sequence soccer, and it shows that DFDO outperforms FFDO by 0.7 dB, averagely.

From figure 4-6c, we can observe that DFDO outperforms FFDO even more when a GOP has fewer video layers protected. This is because the percentage of the number of protected layers is higher in GOPs which have fewer video layers protected, and can be observed from the protection result as shown in table 4-5.

Further, in order to examine the performance of DFDO when the video content has wide and global motion, we down-sample the sequences football and soccer from 30fps to 15fps. The encoding parameters are listed in table 4-4, such as packet loss rate and both leaky factors used in simulation remain the same. The graph of given bandwidth is shown in figure 4-5. The simulation results for both videos are shown in figure 4-7a and 4-7b, and DFDO outperforms FFDO by 0.5~0.7 dB. In the case of football, it is almost twice as compared with the 30fps



Fig 4-5b: Given bandwidth for soccer @ 15 fps

Layer ID	Resolution	Bit-rate (kbit/sec)	PSNR ^{football} (dB)	PSNR ^{soccer} (dB)
0	QCIF	200	28.4286	28.6826
1	QCIF	400	29.9313	29.1768
2	CIF	600	32.7111	30.3753
3	CIF	800	34.2528	30.8541
4	4CIF	1000	34.4149	32.9559

Table 4-4: Encoding parameters and PSNR for sequences football and soccer at 30fps

We also compare both algorithms under the case of misestimate of packet loss rate.

Figure 4-8 shows the simulation results on the 30fps and 15fps football sequences, where the x axis means the actual packet loss rate in channel and the y axis means how much DFDO outperforms FFDO. The misestimated packet loss rate is 0.285714. And we can observe that, although the packet loss rate is not estimated accurately, our DFDO algorithm still outperforms FFDO about 0.25 to 0.75 dB.

11115

m

4.2.2 Simulation Results

COD	Lover	FEDO			DFDO					
UOP	Layer	FFDO	M_0	M_1	M_2	M ₃	M_4			
	0	18	0	0	0	0	18			
	1	16	0	0	0	0	16			
1	2	15	0	0	0	0	15			
	3	13	0	0	0	0	13			
	4	9	0	0	0	0	9			
	0	15	0	0	0	0	15			
2	1	11	0	0	0	0	11			
2	2	10	0	0	0	0	10			
	3	6	0	0	0	6	0			
	0	14	0	0	0	0	14			
3	51/1	8	0	0	0	7	1			
	2	4	0	0	4	0	0			
4	0	14	0	0	0	0	14			
4	1	4	0	0	4	0	0			
5	0	4	0	0	4	0	0			
6	0	. 14	0	0	0	0	14			
0	≥ 1	4		0	4	0	0			
	0	15	0	0	0	0	15			
7	1	8	0	0	0	8	0			
	2	6	0	0	0	6	0			
	0	16	0	0	0	0	16			
8	1	12	0	0	0	0	12			
0	2	11	0	0	0	0	11			
	3	4	0	0	4	0	0			
	0	18	0	0	0	0	18			
9	1	15	0	0	0	0	15			
	2	14	0	0	0	0	14			
	3	11	0	0	0	0	11			
	4	6	0	0	0	6	0			

 Table 4-5: Protection results (The number is the amount of protecting symbols)

	0	15	0	0	0	0	15
10	1	12	0	0	0	0	12
10	2	10	0	0	0	0	10
	3	5	0	0	0	5	0
	0	14	0	0	0	0	14
11	1	8	0	0	0	7	1
	2	5	0	0	0	5	0
12	0	14	0	0	0	0	14
12	1	4	0	0	4	0	0
13	0	4	0	0	4	0	0
14	0	14	0	0	0	0	14
14	1	4	0	0	4	0	0
	0	13	0	-0	0	0	13
15	1	9	0	0	0	0	9
	2	4	0	0	4	0	0
	0	15	-0	0	- 0	0	15
16	1/		0	0	0	0	11
10	2	- 11	0	0	0	0	11
	3	5	0	0	0	5	0
	0	18	0	18	X	Х	Х
17	-1	16	0	16	X	Х	Х
	2	15	0	15	X	Х	Х
	3	12	0	12	X	X	Х
	4	5	5	0	X	Χ	Χ
A REAL PROPERTY OF							





Fig 4-6b: Simulation result of DFDO with football, average PSNR is 28.8179 dB



Fig 4-6d: The different of simulation results between DFDO and FFDO, soccer @ 30 fps



Fig 4-7b: Simulation result between DFDO and FFDO, soccer @15 fps



Fig 4-9f: soccer

V. Conclusion

In multimedia transmission, forward error correction code is a useful technique to avoid the retransmission of lost packet, which may lead to a large delay and is not feasible to real-time play. The erasure code protects data by adding some redundancies to the source information. And for a video stream transmitted over a lossy channel which has limited bandwidth, it is important to distribute bandwidth among video data and protection data efficiently in order to have good visual experience.

In this paper, we first modify [5] to be flat FEC-distortion optimization (FFDO) algorithm, which not only can adapt to the scalable video streams encoded with H.264/MPEG-4 AVC scalable extension standard but also can take account inter-layer prediction. Then we propose the dynamic FEC-distortion optimization (DFDO) algorithm, which further improves FFDO so that the pictures within the same video layer of a GOP are protected according to their importance. Thus, when a video layer of GOP can not be completely received, DFDO has more chance to recover important pictures than FFDO does.

The simulation results show that the average PSNR of DFDO outperforms FFDO about 0.4 dB. If we make the video content moves wider by down-sampling from 30fps to 15 fps, DFDO outperforms FFDO about 0.8 dB. We also navigate the performance of DFDO under the case of misestimate of packet loss rate, and the simulation results show that DFDO still outperforms FFDO. Thus we can conclude that protecting pictures within the same video layer according to their importance can improve the visual quality.

Reference

- ETSI, "Digital Video Broadcasting (DVB): Transmission systems for handheld terminals," ETSI standard, EN 302 304 V1.1.1, 2004.
- [2] J. Byers, M. Luby, and M. Mitzenmacher, "A Digital Fountain Approach to Asynchronous Reliable Multicast," IEEE Journal on Selected Areas in Communications, 20(8), pp. 1528-1540, October 2002.
- [3] M. Luby et al., "Raptor Codes for Reliable Download Delivery in Wireless Broadcast Systems," IEEE CCNC, Las Vegas, NV, Jan. 2006.
- [4] 3GPP TS 26.346 V6.4.0, "Technical Specification Group Services and System Aspects; Multimedia Broadcast/Multicast Service (MBMS); Protocols and Codecs," Mar. 2006.
- [5] W.-T. Tan, A. Zakhor, "Video multicast using layered FEC and scalable compression," IEEE Transactions on Circuits and Systems for Video Technology, March 2001.
- [6] H.-F. Hsiao, A. Chindapol, J. A. Ritcey, and J.-N. Hwang, "Adaptive FEC Scheme for Layered Multimedia Streaming over Wired/Wireless Channels," Workshop on Multimedia Signal Processing, IEEE, pp. 1-4, Oct. 2005.
- [7] S.-R. Kang and D. Loguinov, "Modeling Best-Effort and FEC Streaming of Scalable Video in Lossy Network Channels," IEEE/ACM Trans. on Networking, Feb. 2007.
- [8] W.-T, A. Zakhor, "Real-Time Internet Video Using Error Resilient Scalable Compression and TCP-Friendly Transport Protocol," IEEE Transactions On Multimedia, Vol. 1, No. 2, JUNE 1999
- [9] H. Schwarz, D. Marpe, and T. Wiegand, "Overview of the Scalable H.264/MPEG4-AVC Extension," International Conference on Image Processing, IEEE, Oct., 2006.
- [10] L. Rizzo, "Effective Erasure Codes for Reliable Computer Communication Protocols", Computer Communication Review, 27(2):24-36, April 1997.

- [11] ITU-T Rec. & ISO/IEC 14496-10 AVC, "Advanced Video Coding for Generic Audiovisual Services," version 3, 2005.
- [12] San Ling and Chaoling Xing, "Coding Theory, A First Course", Cambridge University Press
- [13] http://ip.hhi.de/imagecom_G1/savce/index.htm

m

- [14] Yao Wang, Jorn Ostermann, and Ya-Qin Zhang, "Video Processing And Communications", Prentice Hall, pp. 349
- [15] S. R. Hankerson, D. G. Hoffman, D. A. Leonard, C. C. Lindner, K. T. Phelps, C. A. Rodger, J. R. Wall, "Coding Theory And Cryptography, The Essentials, 2nd edition", Marcel Dekker, Inc.
- [16] Jerome Lecan and Jerome Fimes, "Systematic MDS Erasure Codes Based on Vandermonde Matrices", IEEE Communications Letters, Vol. 8, No. 9, Sep. 2004
- [17] I.Gohberg and V.Olshevsky, "Fast algorithms with preprocessing for matrix-vector multiplication problems", Journal of Complexity, 1994

11111